

Parallel solution of large sparse eigenproblems using a Block-Jacobi-Davidson method

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Background

Aim: Calculate a set of outer eigenpairs (λ_i, v_i) of a large sparse matrix:

$$Av_i = \lambda_i v_i$$

Project: Equipping Sparse Solvers for Exascale (ESSEX) of the DFG SPPEXA programme

Outline

Block-Jacobi-Davidson algorithm

Performance analysis

Implementation

Results

Conclusion

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Jacobi-Davidson QR method

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Sketch of the algorithm

- 1: **while** not converged **do** ▷ Outer iteration
- 2: Project problem to small subspace
- 3: Solve small eigenvalue problem
- 4: Calculate approximation and residual
- 5: Approximately solve correction equation ▷ Inner iteration
- 6: Orthogonalize result
- 7: Enlarge subspace
- 8: **end while**

Outer iteration

Subspace iteration

Deflation

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Subspace iteration

- ▶ Galerkin projection onto small subspace $\mathcal{W} \subset \mathbb{R}^n$:

$$\begin{aligned} & Av - \lambda v \perp \mathcal{W}, & v \in \mathcal{W} \\ \Leftrightarrow & (W^T AW)s - \tilde{\lambda}s = 0, & \text{for orth. basis } W \end{aligned}$$

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- ▶ Project out already known eigenvector space Q :

$$\bar{A} := (I - QQ^T)A(I - QQ^T)$$

Jacobi-Davidson correction equation

Basic approach

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- Solve approximately:

$$(I - \tilde{Q}\tilde{Q}^T)(A - \tilde{\lambda}I)(I - \tilde{Q}\tilde{Q}^T)w_{k+1} = -r$$

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- ▶ Provides new direction w_{k+1} for the subspace iteration

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Performance analysis

- Required linear algebra operations

- spMMVM single node

- Block vector operations

- spMMVM inter-node

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Sparse matrix-multiple-vector multiplication (spMMVM)

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- ▶ different types of operations: ($V, W \in \mathbb{R}^{n \times n_b}$)

	local	all-reduction
BLAS 1	$V + W$	$\ v_i\ $
BLAS 3	VM	$V^T W$

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- ▶ message aggregation for all-reductions
- ▶ better code balance for BLAS 3 operations (memory bounded)

spMMVM single node

Setup

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- ▶ CRS format
- ▶ 6-core Intel Westmere CPU

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$$B_{CRS,NT} = \frac{6}{n_b} + \frac{8}{n_{nzs}} + \frac{\kappa}{2} \left[\frac{\text{bytes}}{\text{flops}} \right]$$

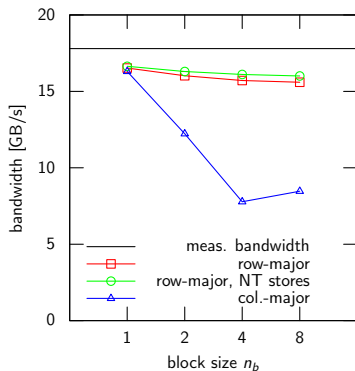
spMMVM single node (2)

Bandwidth

Performance

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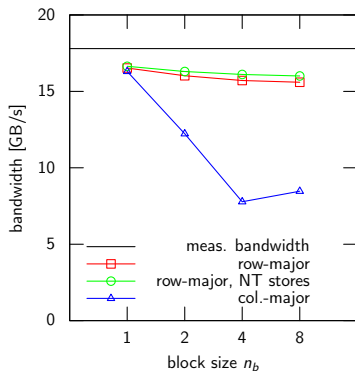
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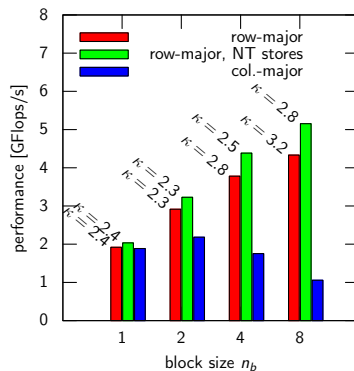
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Block vector operations

Jacobi-Davidson Operator

- ▶ spMMVM + 2 × GEMM:

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with $Q \in \mathbb{R}^{n \times 10}$, $i = 1, \dots, n_b$

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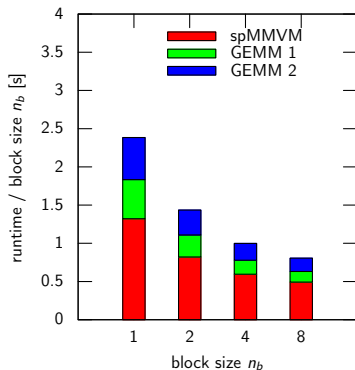
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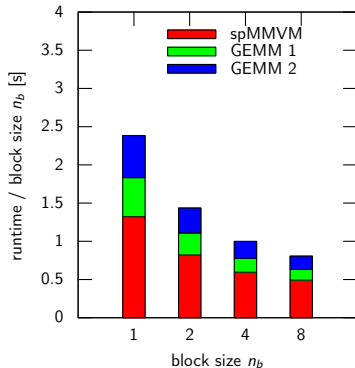
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- ▶ 1.6 times faster for $n_b = 2$
- ▶ 2.5 times faster for $n_b = 4$



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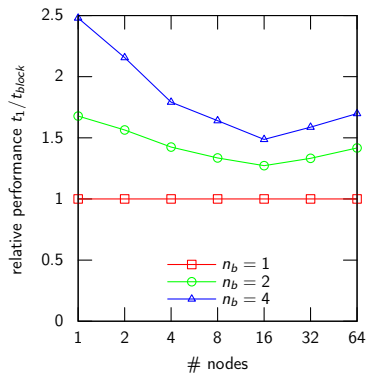
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phist (Pipelined Hybrid-parallel Linear Solver Toolkit)

GHOST (General Hybrid Optimized Sparse Toolkit)

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phist (Pipelined Hybrid-parallel Linear Solver Toolkit)

- ▶ General C-interface to linear algebra libraries:
 - ▶ Trilinos (C++, <http://trilinos.sandia.gov>)
 - ▶ GHOST
 - ▶ **self-written** (row-major storage, Fortran + C99, NT-stores)

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- ▶ Developed at the RRZE in Erlangen (ESSEX project partner)
- ▶ Hybrid-parallel MPI+OpenMP+CUDA

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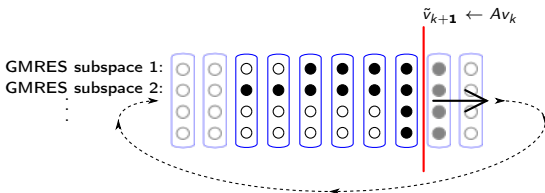
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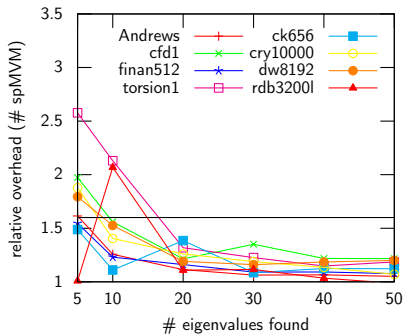
Numerical behavior

Performance of the complete algorithm

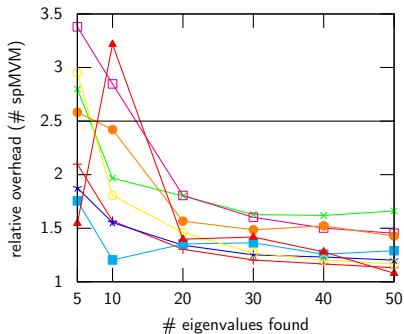
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Block size 2

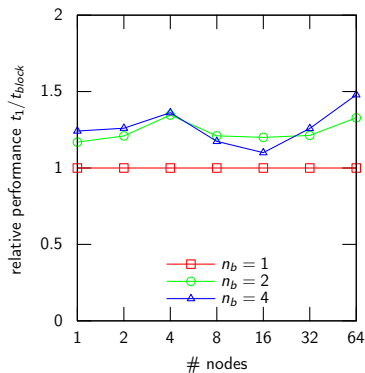


Block size 4

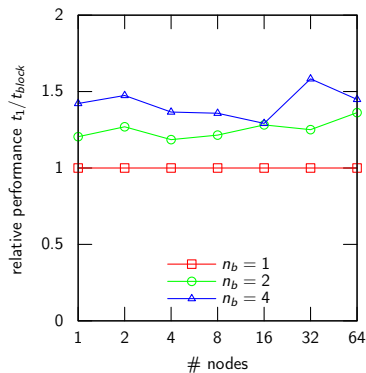


Performance of the complete algorithm

10 eigenvalues



20 eigenvalues



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 - ▶ Great impact of memory layout (row- instead of col.-major)
 - ▶ Improves code balance (node-level)
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⇒ Improved performance by factor 1.2–1.6

Conclusion (2)

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- ▶ Block Pipelined GMRES also interesting for other applications (FEAST in ESSEX)