

Computation of Minimal Realizations of Periodic Systems

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Abstract

We propose balancing related numerically reliable methods to compute minimal realizations of linear periodic systems with time-varying dimensions. The first method belongs to the family of *square-root* methods with guaranteed enhanced computational accuracy and can be used to compute balanced minimal order realizations. An alternative *balancing-free square-root* method has the advantage of a potentially better numerical accuracy in case of poorly scaled original systems. The key numerical computation in both methods is the solution of nonnegative periodic Lyapunov equations directly for the Cholesky factors of the solutions. For this purpose, a numerically reliable computational algorithm is proposed to solve nonnegative periodic Lyapunov equations with time-varying dimensions.

1 Introduction

In the last few years there has been a constantly increasing interest to develop numerical algorithms for the analysis and design of linear periodic discrete-time control systems with constant state-, input- and output-vector dimensions [1, 6, 8, 11]. Areas where significant theoretical results have been achieved for periodic systems with time-varying state dimensions are the solution of the minimal realization problem [4, 3] and robust pole assignment [7].

In this paper we develop balancing related numerical methods to compute minimal realizations of linear periodic systems with time-varying dimensions. The first method can be used to compute balanced minimal realizations and belongs to the family of *square-root* methods with guaranteed enhanced computational accuracy. The matrices of the minimal realization are computed using appropriate truncation matrices determined exclusively using the Cholesky factors of the periodic reachability and observability Gramians.

An alternative *balancing-free square-root* method for

minimal realization has the advantage of a potentially better numerical accuracy in case of poorly scaled original nonminimal systems. By avoiding the use of possibly ill-conditioned balancing based truncation formulas and by using instead well-conditioned transformations, the accuracy of computations can be often significantly improved. The second method extends the balancing-free square-root approach for standard systems [10] to periodic systems.

The key computation in the proposed computational approaches is the solution of nonnegative periodic Lyapunov equations directly for the Cholesky factors of the Gramians. For this purpose, a numerically reliable computational algorithm is proposed to solve nonnegative periodic Lyapunov equations with time-varying dimensions. The proposed algorithm is an extension of a method proposed by the author for constant dimensions [11].

2 Preliminaries

Consider the linear discrete-time K -periodic system

$$\begin{aligned}x_{k+1} &= A_k x_k + B_k u_k \\ y_k &= C_k x_k\end{aligned}\quad (1)$$

where the matrices $A_k \in \mathbb{R}^{n_{k+1} \times n_k}$, $B_k \in \mathbb{R}^{n_{k+1} \times m_k}$, $C_k \in \mathbb{R}^{p_k \times n_k}$ and the integers n_k , m_k , p_k are periodic with period $K \geq 1$. The transition matrix of the system (1) is defined by the $n_j \times n_i$ matrix $\Phi_A(j, i) = A_{j-1} A_{j-2} \cdots A_i$, where $\Phi_A(i, i) := I_{n_i}$. The state transition matrix over one period $\Phi_A(j + K, j) \in \mathbb{R}^{n_j \times n_j}$ is called the *monodromy matrix* of system (1) at time j and its eigenvalues are called *characteristic multipliers* at time j . Note that $\Lambda(\Phi_A(j + K, j))$ has always at least $n_j - \underline{n}$ zero elements, where $\underline{n} := \min_k \{n_k\}$. The rest of \underline{n} eigenvalues are independent of time j and form the *core characteristic multipliers*. The periodic system (1) is *asymptotically stable* if all characteristic multipliers belong to the open unit disk. For the definitions of reachability, observability and minimality of periodic systems we use the corresponding notions from

[4] for general time-varying systems.

Definition 1. The periodic system (1) is *reachable at time k* if

$$\text{rank } G_k = n_k, \quad (2)$$

where G_k is the infinite columns matrix

$$G_k = [B_{k-1} \ A_{k-1}B_{k-2} \ \cdots \ \Phi_A(k, i+1)B_i \ \cdots]. \quad (3)$$

The periodic system (1) is *completely reachable* if (2) holds for all k .

Definition 2. The periodic system (1) is *observable at time k* if

$$\text{rank } F_k = n_k, \quad (4)$$

where F_k is the infinite rows matrix

$$F_k = \begin{bmatrix} C_k \\ C_{k+1}A_k \\ \vdots \\ C_i\Phi_A(i, k) \\ \vdots \end{bmatrix}. \quad (5)$$

The periodic system (1) is *completely observable* if (4) holds for all k .

Definition 3. The periodic system (1) is *minimal* if it is completely reachable and completely observable.

For an asymptotically stable periodic system, the $n_k \times n_k$ reachability Gramian at time k is defined as

$$P_k := \sum_{i=-\infty}^{k-1} \Phi_A(k, i+1)B_iB_i^T\Phi_A(k, i+1)^T = G_kG_k^T \geq 0,$$

where G_k is defined in (3). Similarly, the $n_k \times n_k$ observability Gramian at time k is defined as

$$Q_k = \sum_{i=k}^{\infty} \Phi_A(i, k)^TC_i^TC_i\Phi_A(i, k) = F_k^TF_k \geq 0.$$

Note that both Gramians are K -periodic matrices. Using the definitions of reachability and observability we have the following results.

Proposition 1 *The periodic system (1) is reachable at time k iff $P_k > 0$ and is completely reachable iff $P_k > 0$ for $k = 1, \dots, K-1$.*

Proposition 2 *The periodic system (1) is observable at time k iff $Q_k > 0$ and is completely observable iff $Q_k > 0$ for $k = 1, \dots, K-1$.*

Notation. For a K -periodic matrix X_k we use alternatively the *script* notation

$$\mathcal{X} := \text{diag}(X_0, X_1, \dots, X_{K-1}),$$

which associates the block-diagonal matrix \mathcal{X} to the cyclic matrix sequence X_k , $k = 0, \dots, K-1$. This notation is consistent with the standard matrix operations as for instance addition, multiplication, inversion as well as with several standard matrix decompositions (Cholesky, SVD). We denote with $\sigma\mathcal{X}$ the K -cyclic shift

$$\sigma\mathcal{X} = \text{diag}(X_1, \dots, X_{K-1}, X_0)$$

of the cyclic sequence X_k , $k = 0, \dots, K-1$. By using the script notation, the periodic system (1) will be alternatively denoted by the quadruple $(\mathcal{A}, \mathcal{B}, \mathcal{C})$.

3 Square-root Minimal Realization

For an asymptotically stable periodic system the two Gramians are nonnegative definite and satisfy *nonnegative (or positive) discrete periodic Lyapunov equations* (PDPLEs) as follows: the reachability Gramian \mathcal{P} satisfies the *forward-time* PDPLE

$$\sigma\mathcal{P} = \mathcal{A}\mathcal{P}\mathcal{A}^T + \mathcal{B}\mathcal{B}^T, \quad (6)$$

while the observability Gramian \mathcal{Q} satisfies the *reverse-time* PDPLE

$$\mathcal{Q} = \mathcal{A}^T\sigma\mathcal{Q}\mathcal{A} + \mathcal{C}^T\mathcal{C}. \quad (7)$$

Let $T_k \in \mathbb{R}^{n_k \times n_k}$ be a K -periodic invertible matrix. Two periodic systems $(\mathcal{A}, \mathcal{B}, \mathcal{C})$ and $(\tilde{\mathcal{A}}, \tilde{\mathcal{B}}, \tilde{\mathcal{C}})$ related by the transformation

$$(\tilde{\mathcal{A}}, \tilde{\mathcal{B}}, \tilde{\mathcal{C}}) = (\sigma T^{-1}\mathcal{A}T, \sigma T^{-1}\mathcal{B}, \mathcal{C}T) \quad (8)$$

are called *Lyapunov-similar* and (8) is called a *Lyapunov similarity transformation*. The Gramians $\tilde{\mathcal{P}}$ and $\tilde{\mathcal{Q}}$ of the transformed system $(\tilde{\mathcal{A}}, \tilde{\mathcal{B}}, \tilde{\mathcal{C}})$ satisfy

$$\tilde{\mathcal{P}} = T^{-1}\mathcal{P}T^{-T}, \quad \tilde{\mathcal{Q}} = T^T\mathcal{Q}T.$$

For a completely reachable and completely observable (i.e., minimal) periodic system, T can be determined such that the transformed Gramians are equal and diagonal and thus the transformed periodic system is *balanced* [11]. The diagonal elements of the balanced Gramians are called the *Hankel singular values* and are the positive square-roots of the eigenvalues of the product $\mathcal{P}\mathcal{Q}$. The maximum of them defines the *Hankel norm* of the given periodic system [1].

Let $\mathcal{P} = \mathcal{S}^T\mathcal{S}$ and $\mathcal{Q} = \mathcal{R}^T\mathcal{R}$ be in Cholesky factorized forms. For a minimal system, in analogy with the standard case [9], we can use the singular value decomposition

$$\mathcal{R}\mathcal{S}^T = \mathcal{U}\Sigma\mathcal{V}^T, \quad (9)$$

to compute the balancing transformation matrix \mathcal{T} and its inverse \mathcal{T}^{-1} as

$$\mathcal{T} = \mathcal{S}^T\mathcal{V}\Sigma^{-1/2}, \quad \mathcal{T}^{-1} = \Sigma^{-1/2}\mathcal{U}^T\mathcal{R}.$$

For an asymptotically stable non-minimal periodic system $(\mathcal{A}, \mathcal{B}, \mathcal{C})$, the balancing transformations are not defined since the system is not completely reachable and/or not completely observable. Thus, from Proposition 1 or 2 follows that Σ in (9) is not invertible. We will show how it is possible to determine a realization of the given periodic system which is minimal, that is, completely reachable and completely observable, and moreover balanced. The main result (Theorem 1) of this section can be seen as an extension to the periodic case of a similar result in [10] for standard systems.

Let us write the singular value decomposition (9) at each time instant k in the partitioned form

$$R_k S_k^T = [U_{k,1} \ U_{k,2}] \begin{bmatrix} \tilde{\Sigma}_k & 0 \\ 0 & 0 \end{bmatrix} [V_{k,1} \ V_{k,2}]^T, \quad (10)$$

where $\tilde{\Sigma}_k \in \mathbb{R}^{r_k \times r_k}$, $U_{k,1} \in \mathbb{R}^{n_k \times r_k}$, $V_{k,1} \in \mathbb{R}^{n_k \times r_k}$ and $\tilde{\Sigma}_k > 0$. From the above decomposition define, with $\tilde{\Sigma} = \text{diag}(\tilde{\Sigma}_0, \dots, \tilde{\Sigma}_{K-1})$, the *truncation* matrices

$$\mathcal{L} = \tilde{\Sigma}^{-\frac{1}{2}} \mathcal{U}_1^T \mathcal{R}, \quad \mathcal{T} = \mathcal{S}^T \mathcal{V}_1 \tilde{\Sigma}^{-\frac{1}{2}},$$

which are used to determine the reduced system matrices

$$\hat{\mathcal{A}} = \sigma \mathcal{L} \mathcal{A} \mathcal{T}, \quad \hat{\mathcal{B}} = \sigma \mathcal{L} \mathcal{B}, \quad \hat{\mathcal{C}} = \mathcal{C} \mathcal{T}. \quad (11)$$

The following is the main theoretical result of the paper (for the proof see the Appendix).

Theorem 1 *The periodic system $(\hat{\mathcal{A}}, \hat{\mathcal{B}}, \hat{\mathcal{C}})$ defined in (11) is a balanced minimal realization of the system $(\mathcal{A}, \mathcal{B}, \mathcal{C})$.*

The computation of the minimal realization relies exclusively on square-root information (the Cholesky factors of Gramians) and this leads to a guaranteed enhancement of the overall numerical accuracy of computations. The key computation in determining \mathcal{L} and \mathcal{T} is the solution of the two PDPLEs (6) and (7) with time-varying dimensions directly for the Cholesky factors of the Gramians. A numerically reliable procedure for this purpose is discussed Section 5.

4 Balancing-free Square-root Minimal Realization

To obtain a minimal realization from a non-minimal one we do not actually need to obtain a balanced minimal realization since this could involve ill-conditioned \mathcal{L} and \mathcal{T} matrices, if the original system is poorly balanced. To avoid potential accuracy losses, an alternative is to use a *balancing-free* approach to compute the two truncation matrices. A *square-root balancing-free* approach can be easily devised analogously as

in case of standard systems [10]. Consider the QR-decompositions

$$\mathcal{S}^T \mathcal{V}_1 = \tilde{\mathcal{T}} \mathcal{X}, \quad \mathcal{R}^T \mathcal{U}_1 = \tilde{\mathcal{Z}} \mathcal{Y}, \quad (12)$$

where \mathcal{X} and \mathcal{Y} are nonsingular matrices and $\tilde{\mathcal{T}}$ and $\tilde{\mathcal{Z}}$ are matrices with orthonormal columns. With the already computed $\tilde{\mathcal{T}}$ we define the corresponding $\tilde{\mathcal{L}}$ as

$$\tilde{\mathcal{L}} = (\tilde{\mathcal{Z}}^T \tilde{\mathcal{T}})^{-1} \tilde{\mathcal{Z}}^T. \quad (13)$$

We have the following result analogous to Theorem 1.

Theorem 2 *The periodic system*

$$(\tilde{\mathcal{A}}, \tilde{\mathcal{B}}, \tilde{\mathcal{C}}) := (\sigma \tilde{\mathcal{L}} \mathcal{A} \tilde{\mathcal{T}}, \sigma \tilde{\mathcal{L}} \mathcal{B}, \mathcal{C} \tilde{\mathcal{T}})$$

with $\tilde{\mathcal{L}}$ and $\tilde{\mathcal{T}}$ defined in (12) and (13) is a minimal realization of the system $(\mathcal{A}, \mathcal{B}, \mathcal{C})$.

The resulted system $(\tilde{\mathcal{A}}, \tilde{\mathcal{B}}, \tilde{\mathcal{C}})$ is clearly not balanced. However, it can be shown that it is Lyapunov-similar to the balanced minimal realization (11). Since the balancing-free square-root approach always constructs well-conditioned truncation matrices, the use of this method for poorly scaled systems leads certainly to an improvement of the overall computational accuracy.

5 Solution of PDPLEs

The main computational problem to compute a minimal realization of an asymptotically stable periodic system is the solution of a PDPLE of the form

$$\mathcal{U}^T \mathcal{U} = \mathcal{A}^T \sigma \mathcal{U}^T \sigma \mathcal{U} \mathcal{A} + \mathcal{R}^T \mathcal{R} \quad (14)$$

directly for the Cholesky factor \mathcal{U} , where $U_k \in \mathbb{R}^{n_k \times n_k}$, $A_k \in \mathbb{R}^{n_{k+1} \times n_k}$, $R_k \in \mathbb{R}^{n_k \times n_k}$, and the dimension n_k are periodic with period $K \geq 1$. To solve PDPLEs with constant dimensions, numerically reliable algorithms have been recently proposed in [11], representing extensions of a method for standard systems [5]. In this section we describe an extension of the method of [11] to solve PDPLEs with time-varying dimensions.

A straightforward embedding of the problem with time-varying dimensions into a larger order problem with constant dimension allows to solve the PDPLE (14) by using algorithms for constant dimension [11]. Let $\bar{n} = \max_k \{n_k\}$ and consider the extended $\bar{n} \times \bar{n}$ matrices

$$A_{k,e} = \begin{bmatrix} A_k & 0 \\ 0 & 0 \end{bmatrix}, \quad R_{k,e} = \begin{bmatrix} R_k & 0 \\ 0 & 0 \end{bmatrix}, \quad U_{k,e} = \begin{bmatrix} U_k & 0 \\ 0 & 0 \end{bmatrix}, \quad (15)$$

where the zero matrices have appropriate dimensions. Then it is easy to see that \mathcal{U}_e is the solution of the PDPLE

$$\mathcal{U}_e^T \mathcal{U}_e = \mathcal{A}_e^T \sigma \mathcal{U}_e^T \sigma \mathcal{U}_e \mathcal{A}_e + \mathcal{R}_e^T \mathcal{R}_e \quad (16)$$

with constant dimensions. The main drawback of this approach is that working with extended matrices with many zero elements leads to an unnecessary loss of computational efficiency. Alternatively, an efficient approach can be devised which fully exploits the underlying problem structure.

In the approach which we propose, the key role plays an extension of the *periodic Schur decomposition* (PSD) of a cyclic product of square matrices and of the corresponding algorithms for its computation [2, 6].

Proposition 3 *Let $A_k \in \mathbb{R}^{n_{k+1} \times n_k}$, $k = 0, 1, \dots, K-1$, with $n_K = n_0$ be arbitrary matrices and let $\underline{n} = \min_k \{n_k\}$. Then there exist orthogonal matrices $Z_k \in \mathbb{R}^{n_k \times n_k}$ such that the matrices $\tilde{A}_k = Z_{k+1}^T A_k Z_k$ for $k = 0, \dots, K-1$ are block upper triangular*

$$\tilde{A}_k = \begin{bmatrix} \tilde{A}_{k,11} & \tilde{A}_{k,12} \\ 0 & \tilde{A}_{k,22} \end{bmatrix}, \quad (17)$$

where $\tilde{A}_{k,11} \in \mathbb{R}^{\underline{n} \times \underline{n}}$, $\tilde{A}_{k,22} \in \mathbb{R}^{(n_{k+1}-\underline{n}) \times (n_k-\underline{n})}$ for $k = 0, 1, \dots, K-1$. Moreover, $\tilde{A}_{K-1,11}$ is in a real Schur form (RSF), $\tilde{A}_{k,11}$ for $k = 0, \dots, K-2$ are upper triangular and $\tilde{A}_{k,22}$ for $k = 0, \dots, K-1$ are upper trapezoidal.

Proof: Let p be such that $n_p = \underline{n}$ and without loss of generality we assume $p = 0$. Let $U_1 \in \mathbb{R}^{n_1 \times n_1}$ be an orthogonal matrix such that

$$U_1^T A_0 = \begin{bmatrix} A_{0,11} \\ O \end{bmatrix},$$

where $A_{0,11} \in \mathbb{R}^{\underline{n} \times \underline{n}}$. An upper triangular $A_{0,11}$ results for example from the QR decomposition of A_0 . For $k = 1, \dots, K-1$, we successively determine the orthogonal matrices $U_{k+1} \in \mathbb{R}^{n_{k+1} \times n_{k+1}}$ such that each partitioned matrix

$$U_{k+1}^T A_k U_k = \begin{bmatrix} A_{k,11} & A_{k,12} \\ O & A_{k,22} \end{bmatrix}$$

is upper trapezoidal, $A_{k,11} \in \mathbb{R}^{\underline{n} \times \underline{n}}$ and $A_{k,22} \in \mathbb{R}^{(n_{k+1}-\underline{n}) \times (n_k-\underline{n})}$. At this moment, we redefine $A_{0,11}$ from

$$U_1^T A_0 U_{K-1} = \begin{bmatrix} A_{0,11} \\ O \end{bmatrix}.$$

Notice that, excepting the leading diagonal matrices, the transformed matrices $U_{k+1}^T A_k U_k$ have the requested form in (17). To finish, we use the results of [2] to determine the orthogonal matrices $Q_k \in \mathbb{R}^{\underline{n} \times \underline{n}}$ which brings the product $A_{K-1,11} \cdots A_{0,11}$ in RSF using the associated PSD. The overall transformation matrices Z_k , achieving the desired form in (17), are given by $Z_k = U_k \cdot \text{diag}(Q_k, I)$. \square

Note that by using the extended PSD instead of the PSD of the product of extended matrices, a notable

reduction of computational costs could arise if the difference $\bar{n} - \underline{n}$ is significant. It can be easily verified that each $\Phi_{\tilde{A}}(k+K, k) = Z_k^T \Phi_{\tilde{A}}(k+K, k) Z_k$ is in a particular RSF

$$\Phi_{\tilde{A}}(k+K, k) = \begin{bmatrix} \Phi_{\tilde{A}_{11}}(k+K, k) & \times \\ 0 & 0 \end{bmatrix}.$$

and each $\Lambda(\Phi_{\tilde{A}_{11}}(k+K, k))$ defines the set of \underline{n} core characteristic values (independent of k).

Let \mathcal{Z} be an orthogonal Lyapunov transformation to compute the extended PSD of the monodromy matrix $\Phi_A(K, 0)$ in the Proposition 3 and define $\tilde{\mathcal{A}} = \sigma \mathcal{Z}^T \mathcal{A} \mathcal{Z}$ and the upper triangular $\tilde{\mathcal{R}}$ such that $\tilde{\mathcal{R}}^T \tilde{\mathcal{R}} = \mathcal{Z}^T \mathcal{R}^T \mathcal{R} \mathcal{Z}$. The equation (14) becomes after premultiplication with $\tilde{\mathcal{Z}}^T$ and postmultiplication with \mathcal{Z}

$$\tilde{\mathcal{U}}^T \tilde{\mathcal{U}} = \tilde{\mathcal{A}}^T \sigma \tilde{\mathcal{U}}^T \sigma \tilde{\mathcal{U}} \tilde{\mathcal{A}} + \tilde{\mathcal{R}}^T \tilde{\mathcal{R}} \quad (18)$$

where $\tilde{\mathcal{U}} = \mathcal{U} \mathcal{Z}$. After solving this reduced equation for $\tilde{\mathcal{U}}$, the solution of (14) results as $\mathcal{U} = \tilde{\mathcal{U}} \mathcal{Z}^T$. In order to simplify the notations, we assume in what follows that the equation (14) is already in the reduced form (18) and \mathcal{R} is upper triangular.

Let us partition A_k , U_k and R_k analogously as

$$A_k = \begin{bmatrix} A_{k,11} & A_{k,12} \\ 0 & A_{k,22} \end{bmatrix}, \quad U_k = \begin{bmatrix} U_{k,11} & U_{k,12} \\ 0 & U_{k,22} \end{bmatrix},$$

$$R_k = \begin{bmatrix} R_{k,11} & R_{k,12} \\ 0 & R_{k,22} \end{bmatrix}$$

where the upper left blocks are $p_1 \times p_1$ matrices ($p_1 = 1$ or 2). Assuming U_{11} non-singular we can derive recursive formulas which are analogous with those in the case of constant dimensions. By rewriting (14) for the partitioned matrices, we obtain

$$\begin{aligned} \mathcal{U}_{11}^T \mathcal{U}_{11} &= \mathcal{A}_{11}^T \sigma \mathcal{U}_{11}^T \sigma \mathcal{U}_{11} \mathcal{A}_{11} + \mathcal{R}_{11}^T \mathcal{R}_{11} \\ \mathcal{U}_{12} &= \mathcal{M}_1^T \sigma \mathcal{U}_{12} \mathcal{A}_{22} + \mathcal{M}_1^T \sigma \mathcal{U}_{11} \mathcal{A}_{12} + \mathcal{M}_2^T \mathcal{R}_{12} \\ \mathcal{U}_{22}^T \mathcal{U}_{22} &= \mathcal{A}_{22}^T \sigma \mathcal{U}_{22}^T \sigma \mathcal{U}_{22} \mathcal{A}_{22} + \mathcal{R}_{22}^T \mathcal{R}_{22} + \mathcal{Y}^T \mathcal{Y} \end{aligned} \quad (19)$$

where $\mathcal{M}_1 = \sigma \mathcal{U}_{11} \mathcal{A}_{11} \mathcal{U}_{11}^{-1}$, $\mathcal{M}_2 = \mathcal{R}_{11} \mathcal{U}_{11}^{-1}$, and \mathcal{Y} is defined by the K -periodic matrix

$$Y_k = N_k \begin{bmatrix} R_{k,12} \\ U_{k+1,11} A_{k,12} + U_{k+1,12} A_{k,22} \end{bmatrix}$$

with each N_k satisfying

$$P_k = I_{2p_1} - \begin{bmatrix} M_{k,2} \\ M_{k,1} \end{bmatrix} \begin{bmatrix} M_{k,2} \\ M_{k,1} \end{bmatrix}^T = N_k^T N_k.$$

It can be shown that $P_k = P_k^2 \geq 0$ and $\text{rank } P_k = p_1$. N_k can be computed as $N_k = Q_{k,2}^T$ from the QR-decomposition

$$\begin{bmatrix} M_{k,2} \\ M_{k,1} \end{bmatrix} = \begin{bmatrix} Q_{k,1} & Q_{k,2} \end{bmatrix} \begin{bmatrix} V_k \\ 0 \end{bmatrix}.$$

Thus by solving successively the first and second equation from (19) for \mathcal{U}_{11} and \mathcal{U}_{12} , respectively, it remains to solve the third equation, where each $U_{k,22}$ has lower order $n_k - p_1$. After updating the Cholesky factorization the resulting PDPLE for \mathcal{U}_{22} has the same reduced form as the starting equation. The only difference between this procedure and that one in [11] for constant dimension is that after several steps some dimensions become zero. In such cases, the computations can continue without any breaks because, we can freely assume that the missing blocks in all matrices are zero matrices. In rest, the algorithmic details are almost the same as in case of the procedure for constant dimension, although the efficient implementation for time-varying dimensions certainly requires an increased bookkeeping effort. For more details refer to [11].

Remark. To compute the Cholesky factors of the reachability and observability Gramians a single computation of the extended PSD of the monodromy matrix $\Phi_A(K, 0)$ is sufficient. In this way, the cost to solve the two PDPLEs (6) and (7) is almost the same as the cost of solving a single PDPLE.

6 Conclusion

We proposed a numerically sound approach to compute minimal realizations of linear periodic systems with time-varying dimensions. The proposed approach relies on algorithms using exclusively square-root information in form of Cholesky factors of the Gramians and therefore they have guaranteed enhanced computational accuracy. A square-root balancing-free variant has been derived, which we believe to be a completely satisfactory numerical approach to solve periodic minimal realization problems. The key computation in the proposed approach is the numerical solution of PDPLEs directly for the Cholesky factors of the solutions. A numerically reliable computational algorithm has been proposed to solve PDPLEs with varying dimension. This algorithm extends similar algorithms to solve PDPLEs with constant dimension.

The proposed minimal realization methods for periodic systems are not restricted to asymptotically stable periodic systems. For an unstable system, a simple scaling can be used to enforce the stability of the starting representation. For instance, it is possible to replace only A_0 by αA_0 , where $0 < \alpha < 1$ is chosen such that $\alpha \Phi_A(K, 0)$ has eigenvalues in the open unit disc. For the modified system, we can apply either the square-root or balancing-free square-root approach to determine a minimal system. Finally, the computed \hat{A}_0 or \tilde{A}_0 needs to be rescaled to \hat{A}_0/α or \tilde{A}_0/α , respectively.

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A Proof of Theorem 1

Proof. The proof is in two steps. First we prove that $(\hat{\mathcal{A}}, \hat{\mathcal{B}}, \hat{\mathcal{C}})$ is a minimal balanced realization and then we prove that it realizes the same input-output operator as $(\mathcal{A}, \mathcal{B}, \mathcal{C})$.

Let substitute $\mathcal{P} = \mathcal{S}^T \mathcal{S}$ in the equation (6) and pre-multiply it by $\sigma \mathcal{L}$ and postmultiply it by $\sigma \mathcal{L}^T$ to obtain

$$\sigma(\mathcal{L} \mathcal{S}^T \mathcal{S} \mathcal{L}^T) = \sigma \mathcal{L} \mathcal{A} \mathcal{S}^T \mathcal{S} \mathcal{A}^T \sigma \mathcal{L}^T + \sigma \mathcal{L} \mathcal{B} \mathcal{B}^T \sigma \mathcal{L}^T. \quad (20)$$

It can be readily verified that

$$\sigma(\mathcal{L}\mathcal{S}^T\mathcal{S}\mathcal{L}^T) = \sigma\tilde{\Sigma}. \quad (21)$$

From the equation (9) we also have

$$\mathcal{R}\mathcal{S}^T\mathcal{V}_2 = 0. \quad (22)$$

Now, we replace $\mathcal{Q} = \mathcal{R}^T\mathcal{R}$ in the equation (7) and premultiply this equation by $\mathcal{V}_2^T\mathcal{S}$ and postmultiply it by $\mathcal{S}^T\mathcal{V}_2$, to obtain, using also (22),

$$\mathcal{V}_2^T\mathcal{S}\mathcal{A}^T\sigma\mathcal{R}^T\sigma\mathcal{R}\mathcal{A}\mathcal{S}^T\mathcal{V}_2 + \mathcal{V}_2^T\mathcal{S}\mathcal{C}^T\mathcal{C}\mathcal{S}^T\mathcal{V}_2 = 0.$$

It follows immediately that $\mathcal{C}\mathcal{S}^T\mathcal{V}_2 = 0$ and also

$$\sigma\mathcal{R}\mathcal{A}\mathcal{S}^T\mathcal{V}_2 = 0. \quad (23)$$

Using the identity $\mathcal{V}_1\mathcal{V}_1^T + \mathcal{V}_2\mathcal{V}_2^T = I$ and (23), as well as the expressions for \mathcal{L} and \mathcal{T} , the first term in the right-hand side of (20) becomes successively

$$\begin{aligned} \sigma\mathcal{L}\mathcal{A}\mathcal{S}^T\mathcal{S}\mathcal{A}^T\sigma\mathcal{L}^T &= \sigma\mathcal{L}\mathcal{A}\mathcal{S}^T\mathcal{V}_1\mathcal{V}_1^T\mathcal{S}\mathcal{A}^T\sigma\mathcal{L}^T + \sigma\tilde{\Sigma}^{-\frac{1}{2}}\sigma\mathcal{U}_1^T\sigma\mathcal{R}\mathcal{A}\mathcal{S}^T\mathcal{V}_2\mathcal{V}_2^T\mathcal{S}\mathcal{A}^T\sigma\mathcal{R}^T\sigma\mathcal{U}_1\sigma\tilde{\Sigma}^{-\frac{1}{2}} \\ &= \sigma\mathcal{L}\mathcal{A}\mathcal{T}\tilde{\Sigma}\mathcal{T}^T\mathcal{A}^T\sigma\mathcal{L}^T. \end{aligned} \quad (24)$$

Now, by using the definitions of $\hat{\mathcal{A}}$ and $\hat{\mathcal{B}}$ in (11), and taking into account (21) and (24), (20) becomes

$$\sigma\tilde{\Sigma} = \hat{\mathcal{A}}\tilde{\Sigma}\hat{\mathcal{A}}^T + \hat{\mathcal{B}}\hat{\mathcal{B}}^T. \quad (25)$$

In a similar way we can show that

$$\tilde{\Sigma} = \hat{\mathcal{A}}^T\sigma\tilde{\Sigma}\hat{\mathcal{A}} + \hat{\mathcal{C}}^T\hat{\mathcal{C}} \quad (26)$$

and thus $(\hat{\mathcal{A}}, \hat{\mathcal{B}}, \hat{\mathcal{C}})$ is balanced. Since $\tilde{\Sigma} > 0$, from (25) and (26) it follows that $(\hat{\mathcal{A}}, \hat{\mathcal{B}}, \hat{\mathcal{C}})$ is completely reachable and completely observable, and thus also minimal [4]. Moreover, the minimal system is also asymptotically stable.

To complete the proof it remains to show that the two realizations $(\hat{\mathcal{A}}, \hat{\mathcal{B}}, \hat{\mathcal{C}})$ and $(\mathcal{A}, \mathcal{B}, \mathcal{C})$ achieve the same input-output operator. To this end, we define the error system $(\mathcal{A}_e, \mathcal{B}_e, \mathcal{C}_e)$ by

$$A_{k,e} = \begin{bmatrix} A_k & 0 \\ 0 & \hat{A}_k \end{bmatrix}, \quad B_{k,e} = \begin{bmatrix} B_k \\ \hat{B}_k \end{bmatrix}, \quad C_{k,e} = \begin{bmatrix} C_k & -\hat{C}_k \end{bmatrix}.$$

The reachability Gramian \mathcal{P}_e of the error system satisfies the periodic Lyapunov equation

$$\sigma\mathcal{P}_e = \mathcal{A}_e\mathcal{P}_e\mathcal{A}_e^T + \mathcal{B}_e\mathcal{B}_e^T.$$

If we partition $P_{k,e}$ in the form

$$P_{k,e} = \begin{bmatrix} P_{k,11} & P_{k,12} \\ P_{k,12}^T & P_{k,22} \end{bmatrix},$$

then we obtain immediately that $\mathcal{P}_{11} = \mathcal{S}^T\mathcal{S}$, $\mathcal{P}_{22} = \tilde{\Sigma}$, and \mathcal{P}_{12} satisfies the periodic Sylvester equation

$$\sigma\mathcal{P}_{12} = \mathcal{A}\mathcal{P}_{12}\hat{\mathcal{A}}^T + \mathcal{B}\hat{\mathcal{B}}^T. \quad (27)$$

This equation has a unique solution guaranteed by the asymptotic stability of both the non-minimal and minimal realizations.

After substituting $\hat{\mathcal{B}}$ in (11) into (27) we obtain

$$\sigma\mathcal{P}_{12} = \mathcal{A}\mathcal{P}_{12}\hat{\mathcal{A}}^T + \mathcal{B}\mathcal{B}^T\sigma\mathcal{L}^T.$$

Further, we substitute $\mathcal{B}\mathcal{B}^T = \sigma\mathcal{S}^T\sigma\mathcal{S} - \mathcal{A}\mathcal{S}^T\mathcal{S}\mathcal{A}^T$ and $\mathcal{L} = \tilde{\Sigma}^{-\frac{1}{2}}\mathcal{U}_1^T\mathcal{R}$ in the above equation and rearrange the terms to get

$$\begin{aligned} \sigma(\mathcal{P}_{12} - \mathcal{S}^T\mathcal{S}\mathcal{R}^T\mathcal{U}_1\tilde{\Sigma}^{-\frac{1}{2}}) &= \\ \mathcal{A}\mathcal{P}_{12}\hat{\mathcal{A}}^T - \mathcal{A}\mathcal{S}^T\mathcal{S}\mathcal{A}^T\sigma\mathcal{R}^T\sigma\mathcal{U}_1\sigma\tilde{\Sigma}^{-\frac{1}{2}}. \end{aligned} \quad (28)$$

Using again the identity $\mathcal{V}_1\mathcal{V}_1^T + \mathcal{V}_2\mathcal{V}_2^T = I$ and (23), we become for the right-hand side of (28)

$$\begin{aligned} \mathcal{A}\mathcal{P}_{12}\hat{\mathcal{A}}^T - \mathcal{A}\mathcal{S}^T(\mathcal{V}_1\mathcal{V}_1^T + \mathcal{V}_2\mathcal{V}_2^T)\mathcal{S}\mathcal{A}^T\sigma\mathcal{R}^T\sigma\mathcal{U}_1\sigma\tilde{\Sigma}^{-\frac{1}{2}} &= \\ \mathcal{A}(\mathcal{P}_{12} - \mathcal{S}^T\mathcal{V}_1\tilde{\Sigma}^{\frac{1}{2}})\hat{\mathcal{A}}^T, \end{aligned}$$

where we also used the expression of $\hat{\mathcal{A}}$ in (11). Finally, observing that

$$\mathcal{S}\mathcal{R}^T\mathcal{U}_1 = \mathcal{V}\Sigma\mathcal{U}^T\mathcal{U}_1 = \mathcal{V}_1\tilde{\Sigma},$$

we get from (28)

$$\sigma(\mathcal{P}_{12} - \mathcal{S}^T\mathcal{V}_1\tilde{\Sigma}^{\frac{1}{2}}) = \mathcal{A}(\mathcal{P}_{12} - \mathcal{S}^T\mathcal{V}_1\tilde{\Sigma}^{\frac{1}{2}})\hat{\mathcal{A}}^T,$$

which has the unique solution

$$\mathcal{P}_{12} = \mathcal{S}^T\mathcal{V}_1\tilde{\Sigma}^{\frac{1}{2}}.$$

Thus, the reachability Gramian is given by

$$P_{k,e} = \begin{bmatrix} S_k^T S_k & S_k^T V_{k,1} \tilde{\Sigma}_k^{\frac{1}{2}} \\ \tilde{\Sigma}_k^{\frac{1}{2}} V_{k,1}^T S_k & \tilde{\Sigma}_k \end{bmatrix}.$$

Similarly, we find that the observability Gramian is

$$Q_{k,e} = \begin{bmatrix} R_k^T R_k & -R_k^T U_{k,1} \tilde{\Sigma}_k^{\frac{1}{2}} \\ -\tilde{\Sigma}_k^{\frac{1}{2}} U_{k,1}^T R_k & \tilde{\Sigma}_k \end{bmatrix}.$$

A direct check shows that $\mathcal{P}_e\mathcal{Q}_e = 0$ and therefore all Hankel singular values are zero. Using operator theory arguments [4], it results that the Hankel-operator associated to the error system is zero and thus the corresponding input-output (Toeplitz) operator is also zero. Thus the two systems realize the same input-output operator. \square