

High Performance Computing Workshop – Leogang, 2025/02/24

Pia Siegl, Greta Reese, Nis van Hülst, Tomohiro Hashizume, Dieter Jaksch

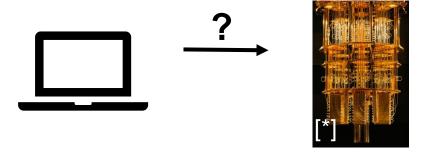


### **Computational Fluid Dynamics on Quantum Computers**



- why should we solve classical differential equations on quantum computers?
  - ➤ beneficial scaling?
  - ➤ reduce computational cost?

• how can we solve classical partial differential equations on quantum computers?



### **Example: Diffusion Equation**



$$\frac{df(x,t)}{dt} = c_d \Delta f(x,t)$$

explicit Euler time steps:

$$f^{j+1} = f^j + \Delta t \cdot c_d \Delta f^j$$

classical implementation:

$$fj+1 - fj \perp \Lambda t \cdot c \cdot \Lambda fj$$

$$f^{j} \rightarrow \begin{pmatrix} f_{0} \\ \vdots \\ f_{N} \end{pmatrix}^{j}, \Delta \rightarrow \frac{1}{\Delta x^{2}} \begin{pmatrix} -2 & 1 & \dots & 0 & 0 \\ 1 & -2 & \dots & 0 & 0 \\ & \vdots & & \ddots & & \vdots \\ 0 & 0 & \dots & -2 & 1 \\ 0 & 0 & & 1 & -2 \end{pmatrix}$$

$$\begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^{j+1} = \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^j + \frac{\Delta t \cdot c_d}{\Delta x^2} \begin{pmatrix} -2 & 1 & \dots & 0 & 0 \\ 1 & -2 & \dots & 0 & 0 \\ \vdots & \ddots & & \vdots & \\ 0 & 0 & \dots & -2 & 1 \\ 0 & 0 & \dots & 1 & -2 \end{pmatrix} \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^j$$

 $\Delta$  = Laplace Operator

 $c_d$  = diffusion constant

 $\Delta t$  = time step size

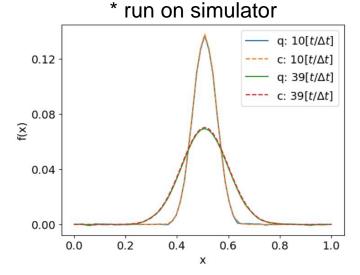
 $\Delta x$  = spatial grid size

### How to Port this on a Quantum Computer?



classical implementation:

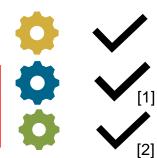
$$\begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^{j+1} = \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^j + dt \begin{pmatrix} -2 & 1 & \dots & 0 & 0 \\ 1 & -2 & \dots & 0 & 0 \\ \vdots & \ddots & & \vdots \\ 0 & 0 & \dots & -2 & 1 \\ 0 & 0 & \dots & 1 & -2 \end{pmatrix} \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^j$$



goal: port this operation on the quantum computer

### main steps:

- 1) encode vector
- 2) apply operators
- 3) compute the new time-step





<sup>[1]</sup> Termanova et al., Quantum Tensor Programming, New J. Phs. 26, 123019, 2024

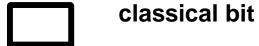
<sup>[2]</sup> M. Lubasch et al., Variational guantum algorithms for nonlinear problems, Phys. Rev. A 101, 2020

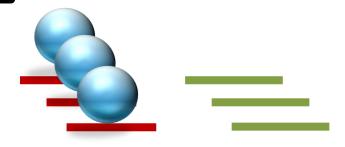
<sup>[3]</sup> P. Siegl et al., Tensor-Programmable Quantum Circuits for differential equations, https://arxiv.org/abs/2502.04425, 2025



# Introduction: Bits and Qubits



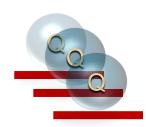




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#### quantum bit (qubit)





$$\sqrt{0.2} |0\rangle + \sqrt{0.8} |1\rangle$$

$$\begin{pmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \\ a_4 \\ a_5 \\ a_6 \\ a_7 \end{pmatrix} \begin{vmatrix} \mathbf{000} \\ |\mathbf{001} \rangle \\ |\mathbf{010} \rangle \\ |\mathbf{011} \rangle \\ |\mathbf{100} \rangle \\ |\mathbf{101} \rangle \\ |\mathbf{110} \rangle \\ |\mathbf{111} \rangle \\ |\mathbf{111} \rangle$$

$$a_i \in \mathbb{C}, \ \sum_{i=0}^N {a_i}^2 = 1$$
 normalized vector



### Encoding Field into a Quantum Computer



recap: 
$$f^{j+1} = f^j + \Delta t \cdot c_d \Delta f^j$$

$$\rightarrow \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^{j+1} = \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^j + dt \begin{pmatrix} -2 & 1 & \dots & 0 & 0 \\ 1 & -2 & \dots & 0 & 0 \\ \vdots & \ddots & & \vdots \\ 0 & 0 & \dots & -2 & 1 \\ 0 & 0 & \dots & 1 & -2 \end{pmatrix} \begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^j$$



field: unnormalized vector  $\begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}$ 



state: normalized vector  $\begin{pmatrix} a_0 \\ \vdots \\ a_N \end{pmatrix}$ 

$$\begin{pmatrix} a_0 \\ \vdots \\ a_N \end{pmatrix}$$

#### connecting both worlds

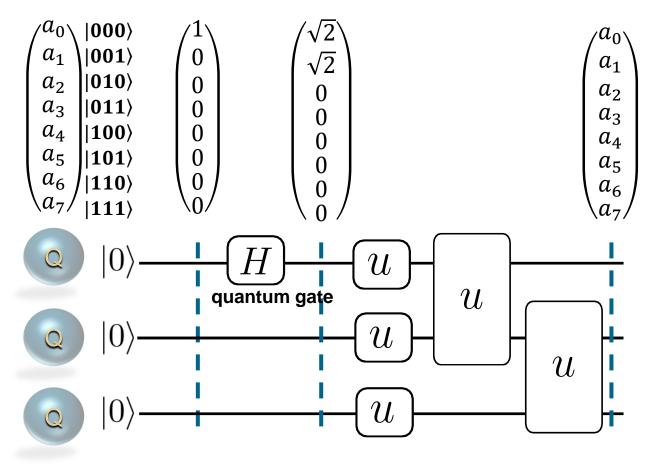


field: 
$$\begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix} = \theta_0 \begin{pmatrix} a_0 \\ \vdots \\ a_N \end{pmatrix}$$



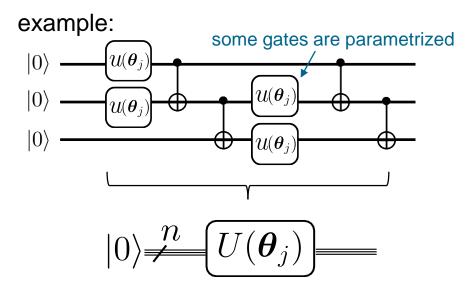
### Encoding Field into a Quantum Computer





quantum circuit

#### here: encode field with classically parametrized ansatz



classical parameters:  $\theta_i$ 



### **Porting Operators on Quantum Computers**



goal: matrix vector multiplication on quantum computer

$$|0\rangle \stackrel{n}{=} U(\boldsymbol{\theta}_j)$$

but: gates are always unitary
(norm-conserving + reversible)



classically

operator: non-unitary



quantum computer

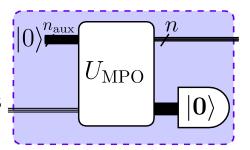
gates: always unitary

#### connecting both worlds



#### probabilistic application of $\Delta$ :

- success probability  $\alpha_{succ}$
- using extra qubits + measurements
- tensor networks





## Computing the New Time Step



recap: explicit Euler time stepping  $f^{j+1} = f^j + \Delta t \cdot c_d \Delta f^j = (I_n + c \Delta) f^j$ 

 $c = \Delta t \cdot c_d$ 

$$\begin{vmatrix} a_0 \\ \vdots \\ a_N \end{vmatrix}^j \qquad \begin{vmatrix} a_0 \\ \vdots \\ a_N \end{vmatrix}^{j+1} \qquad \widehat{0}$$
 quantum measurements 
$$|0\rangle \stackrel{n}{=} U(\boldsymbol{\theta}_j) \qquad \widehat{0} \qquad |0\rangle \text{ or } |1\rangle?$$



### state read out is expensive!

[2] we do:

create a variational ansatz  $U^{\dagger}(oldsymbol{ heta}_{j+1})$ 



compare with

$$|0\rangle \stackrel{n}{=} U(\boldsymbol{\theta}_j)$$

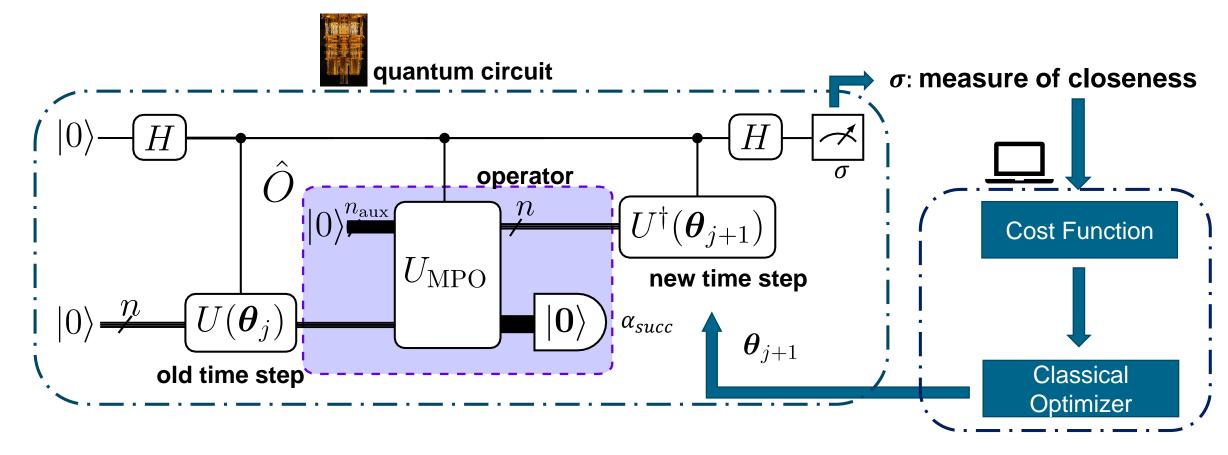




combining tensor network based operator encoding \_\_\_\_with variational time stepping \_\_\_\_









## **Necessary Correction**





**problem:** solution has wrong norm field:  $\begin{pmatrix} f_0 \\ \vdots \\ f_N \end{pmatrix}^J = \theta_0 \begin{pmatrix} a_0 \\ \vdots \\ a_N \end{pmatrix}^j \rightarrow \theta_0 = \theta_0^J$ 

#### reason:

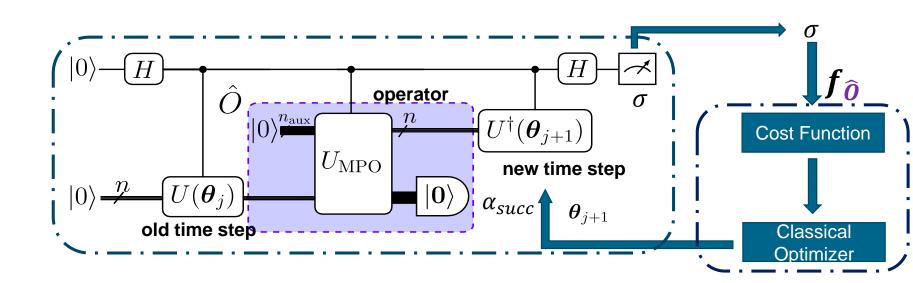
$$\widehat{O}_{classic}\begin{pmatrix} n_0 \\ \dots \\ n_N \end{pmatrix} = \begin{pmatrix} \dots \\ \dots \\ \dots \end{pmatrix}, \qquad \widehat{O}_{quantum}\begin{pmatrix} n_0 \\ \dots \\ n_N \end{pmatrix} = \begin{pmatrix} \dots \\ \dots \\ \dots \end{pmatrix},$$

$$\widehat{O}_{quantum} \begin{pmatrix} n_0 \\ \cdots \\ n_N \end{pmatrix} = \begin{pmatrix} \cdots \\ \cdots \\ \cdots \end{pmatrix},$$

$$\begin{pmatrix} \cdots \\ \cdots \\ \cdots \end{pmatrix} = f_{\widehat{o}} \begin{pmatrix} \cdots \\ \cdots \\ \cdots \end{pmatrix},$$



can be computed from  $\alpha_{succ}$ 







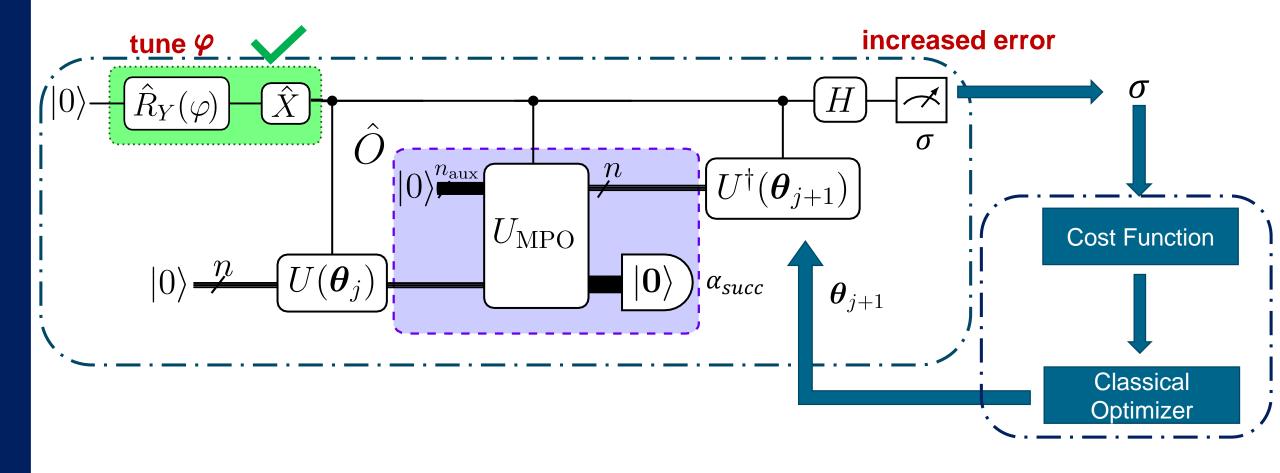
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problem: extra measurements

 $|\mathbf{0}\rangle$ 

increase measurement error of







### **Example: Euler Equation**



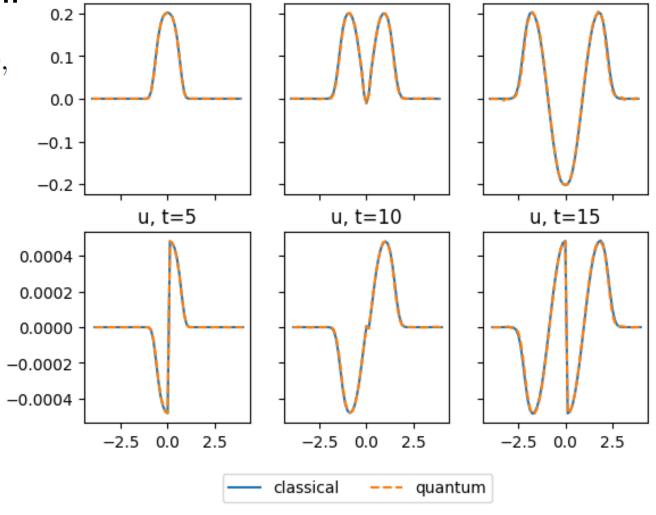
p, t = 15

#### 1D linear incompressible Euler equation

$$\begin{split} \frac{\partial p}{\partial t} &= -\bar{\rho}c^2 \left(\frac{\partial u}{\partial x}\right) + f(x,t) - \gamma(x)p, \\ \frac{\partial u}{\partial t} &= -\frac{1}{\bar{\rho}}\frac{\partial p}{\partial x} - \gamma(x)u, \end{split}$$
 source sponge



> run on quantum simulator



p, t=10

p, t=5

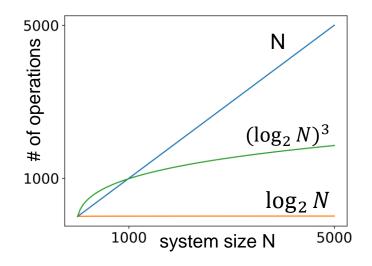
### **Scaling with System Size**



advantage expected for large scale simulations

### classically

matrix – vector multiplication:  $N^2$  sparse matrix-vector multiplication: >N

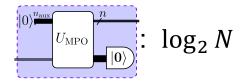




quantum computer → expected:

$$\overline{(U(oldsymbol{ heta}_j)}$$
: poly( $\log_2 N$ )

→ dependence on amount of randomness expected



tensor network algorithms can help to estimate scaling



how does the training scale with  $\theta_i$ ?

- $\rightarrow$  expected: # $\theta_j$  increases poly( $\log_2 N$ )
- → trainability can be assured for larger circuit sizes





- very easy incorporation of various operators with little additional cost:
  - ➤ different boundary conditions, higher order accuracies
  - flexible incorporation of non-derivative operators.
- possible reduction of computational cost

- big advantage identified: in-the loop quantification of quality of solution.
  - > currently working on quantifying this and increasing the accuracy of this.

### **Summary**



- why should we solve partial differential equations on quantum computers?
  - > promising scaling: advantages for large scale simulations?
- can we solve classical partial differential equations on quantum computers?
  - > implementation of broad range of differential equations possible
  - > in-the loop quantification of quality of solution possible













