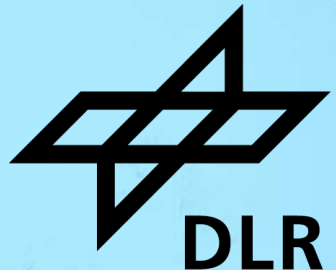


CONTRACT-BASED DESIGN IN MODEL-BASED SYSTEMS ENGINEERING

Ingo Stierand

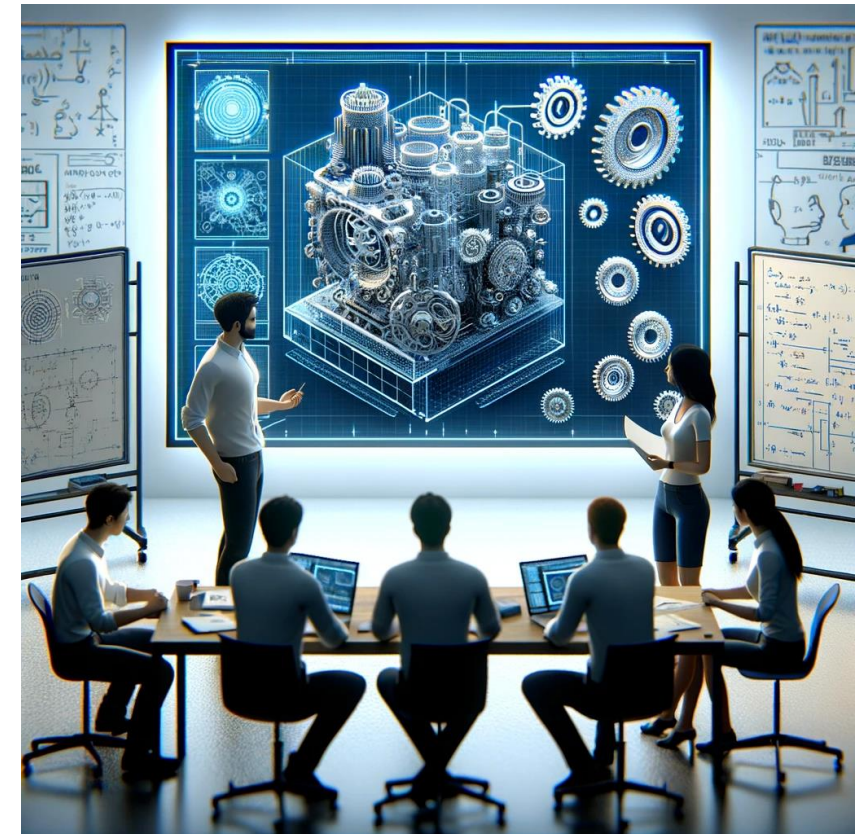
DLR – Institute Systems Engineering for Future Mobility



Why Model-Based (Systems) Engineering

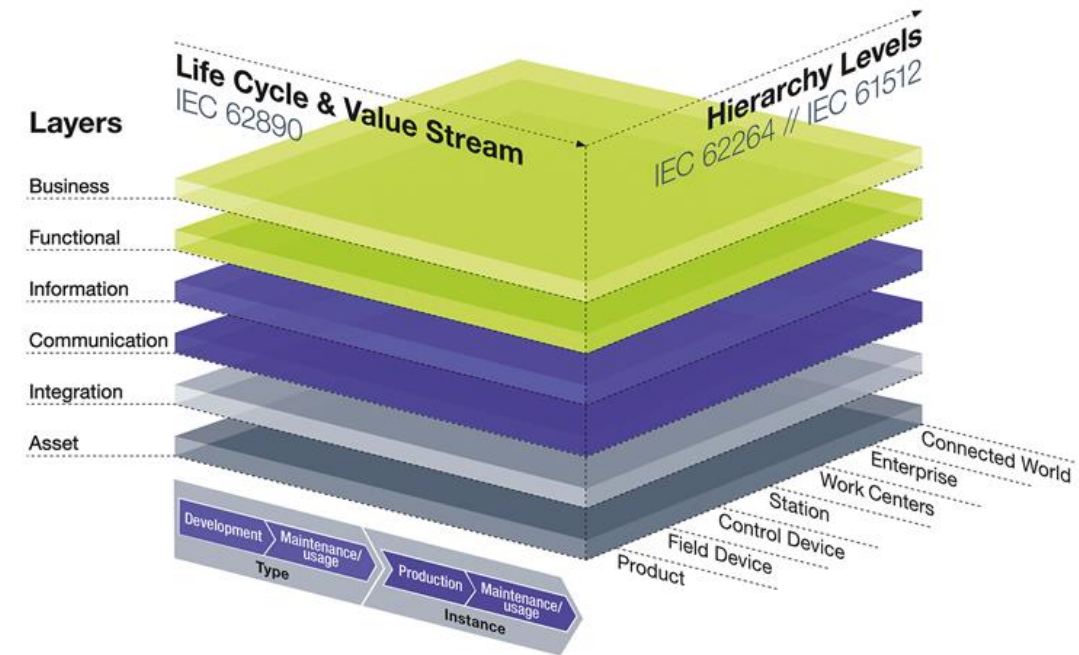
Model-Based Systems Engineering (MBSE) uses models as an integral part of the technical baseline that includes the requirements, analysis, design, implementation, and verification.

- **Complexity Management:** Comprehensive models that represent all aspects of a system help in managing complexity by providing a clear and consistent representation.
- **Collaboration:** Models as a common language provide a shared framework that everyone can understand and contribute to.
- **Adaptability:** Models can be updated more easily than traditional documents, which enables engineers to more rapidly adapt to evolving requirements.
- **Requirements Management:** MBSE enables linking system requirements directly to elements within the model. This ensures traceability and helps to see the impact of changes.
- **Risk Management:** By simulating and analyzing models, MBSE allows engineers to identify potential issues and risks early in the design process.
- **Regulatory Compliance and Safety:** Thoroughness and accuracy of MBSE models assist in ensuring that systems meet regulations and safety requirements.
- **Data Management:** MBSE models can integrate data from various sources and maintain data consistency throughout the system's lifecycle.



Modelling Frameworks

- Each application domain comes with particular demands, regulations and constraints.
- Modelling frameworks help structuring the development process:
 - Layers (perspectives, viewpoints)
 - Life cycle phases
 - Hierarchy levels
 - ...
- Many MBSE modelling languages exist to “fill in” elements in those matrixes:
 - Business Process Model and Notation (BPMN)
 - Structured Goal Notation (GSN)
 - UML diagrams
 - ...

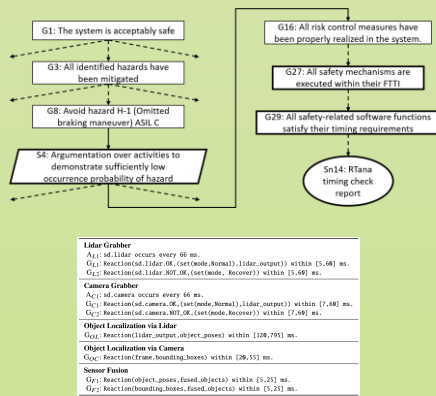


Reference Architecture Model Industry 4.0 – RAMI 4.0 (Source: DKE)

- Focus on (safety-critical) cyber-physical systems.
- Provision of semantically consistent, continuous, traceable MBSE along viewpoints and abstraction levels.

Requirements Viewpoint

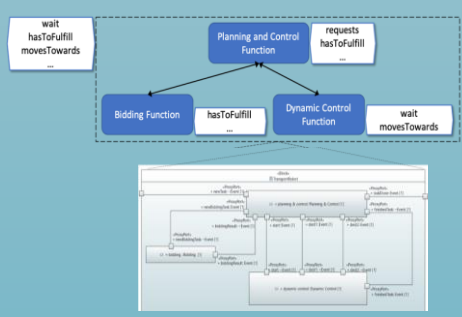
- Operational context
- Goals
- Top-level requirements



Lidar Grabber
A ₁ : lid.lidar occurs every 66 ms.
G ₁ : Reaction(sd.lidar_ok, (set(mode.Normal), lidar_output)) within (5,66) ms.
G ₂ : Reaction(sd.lidar_not_ok, (set(mode.Recovery)) within (5,66) ms.
Camera Grabber
A ₁ : sd.camera occurs every 66 ms.
G ₁ : Reaction(sd.camera_ok, (set(mode.Normal), lidar_output)) within (7,66) ms.
G ₂ : Reaction(sd.camera_not_ok, (set(mode.Recovery)) within (7,66) ms.
Object Localisation via Lidar
G ₁ : Reaction(lidar_camera_object_poses) within (109,793) ms.
Object Localisation via Camera
G ₁ : Reaction(frame_bounding_boxes) within (28,35) ms.
Scene Fusion
G ₁ : Reaction(object_poses_fused_objects) within (5,23) ms.
G ₂ : Reaction(bounding_boxes_fused_objects) within (5,23) ms.

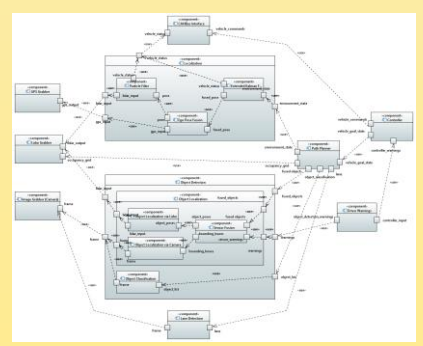
Functional Viewpoint

- System behavior, and
- Dependencies
- Decomposition of functions



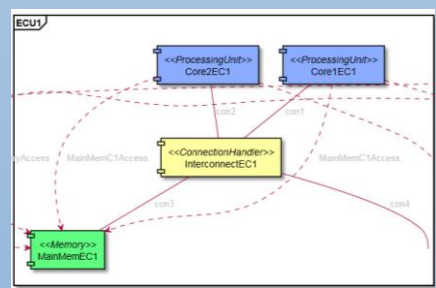
Logical Viewpoint

- Decomposition into sub-systems and components
- Assigning functions to logical units



Technical Viewpoint

- Hardware/Software Design
- Scheduling/middleware
- Optimization of resources
- Communication

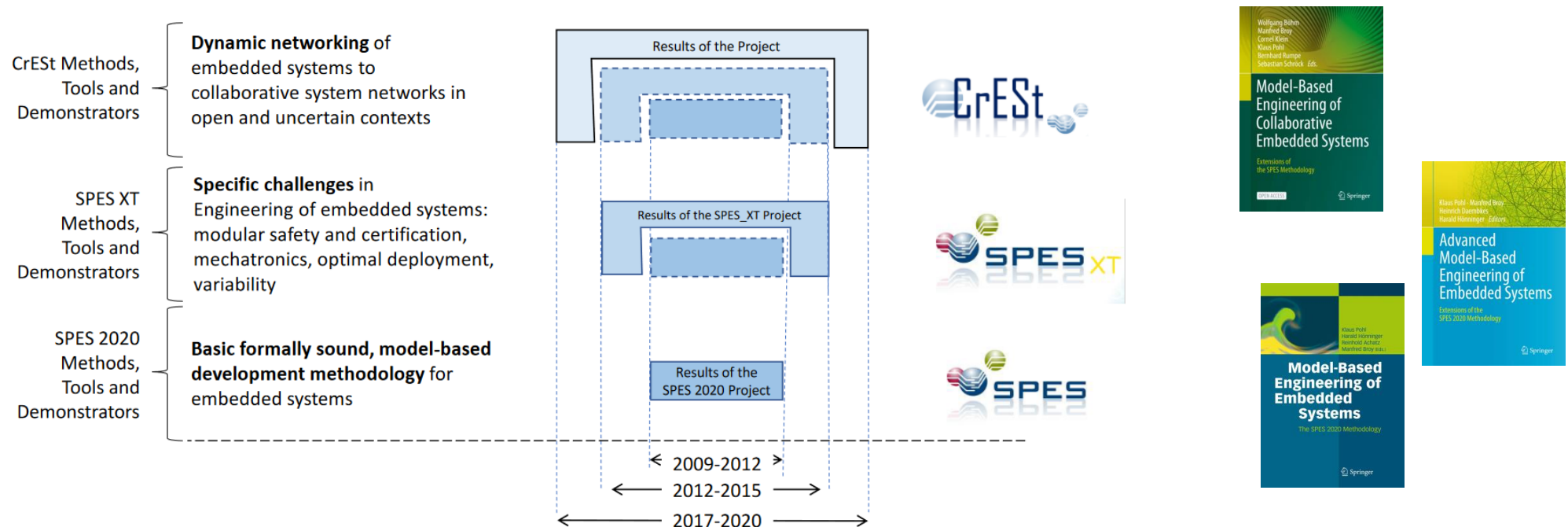


Abstraction levels

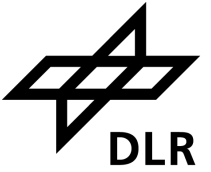
History: From SPES 2020 to CrEst

The Vision of SPES 2020

The development of software intense systems (CPS) can be accomplished through a set of **integrated modeling techniques**, their use and integration in the development is fully understood.



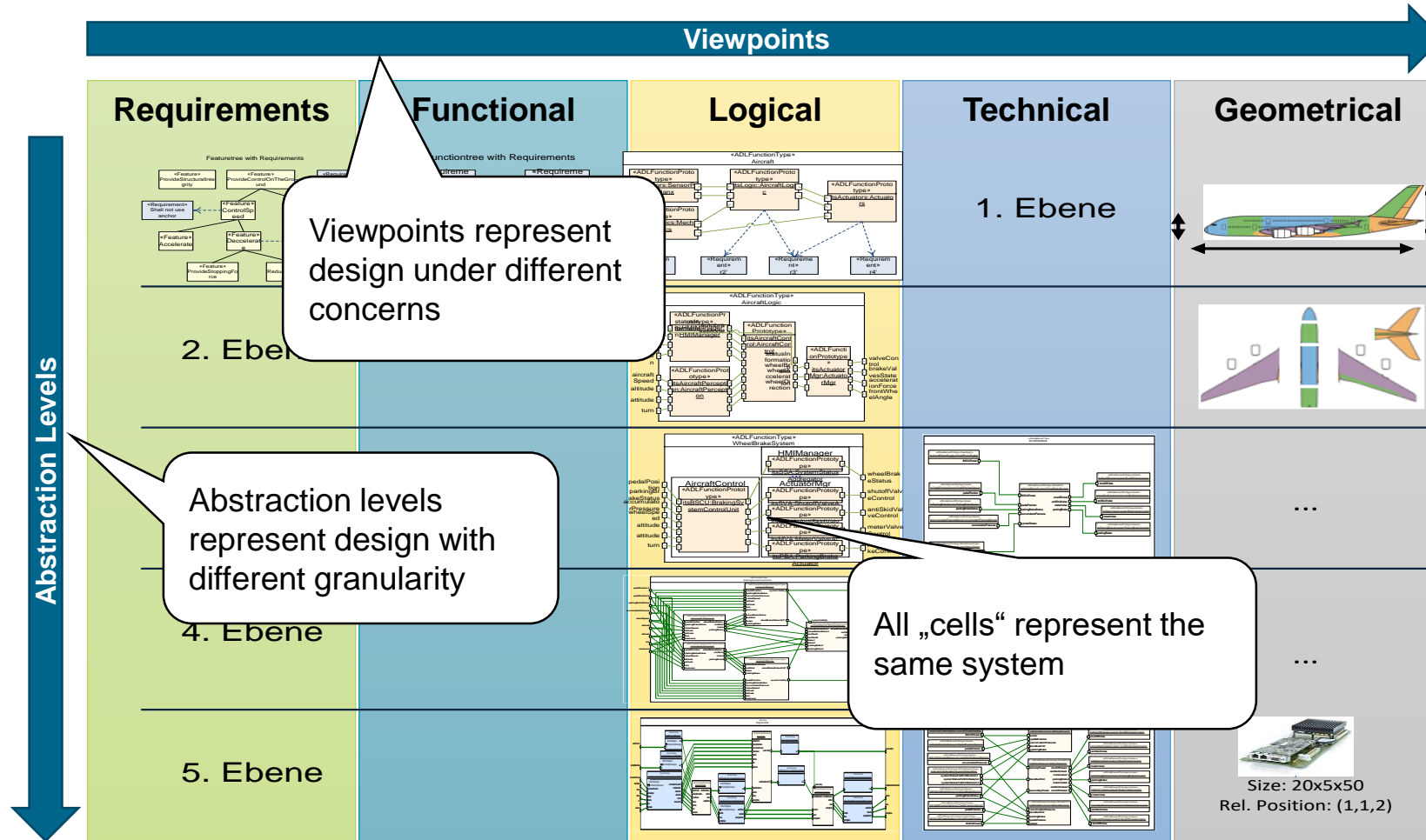
Agenda



- **Compositional Semantic Framework**
- A "Meta Theory" of Contract Based Design
- Example

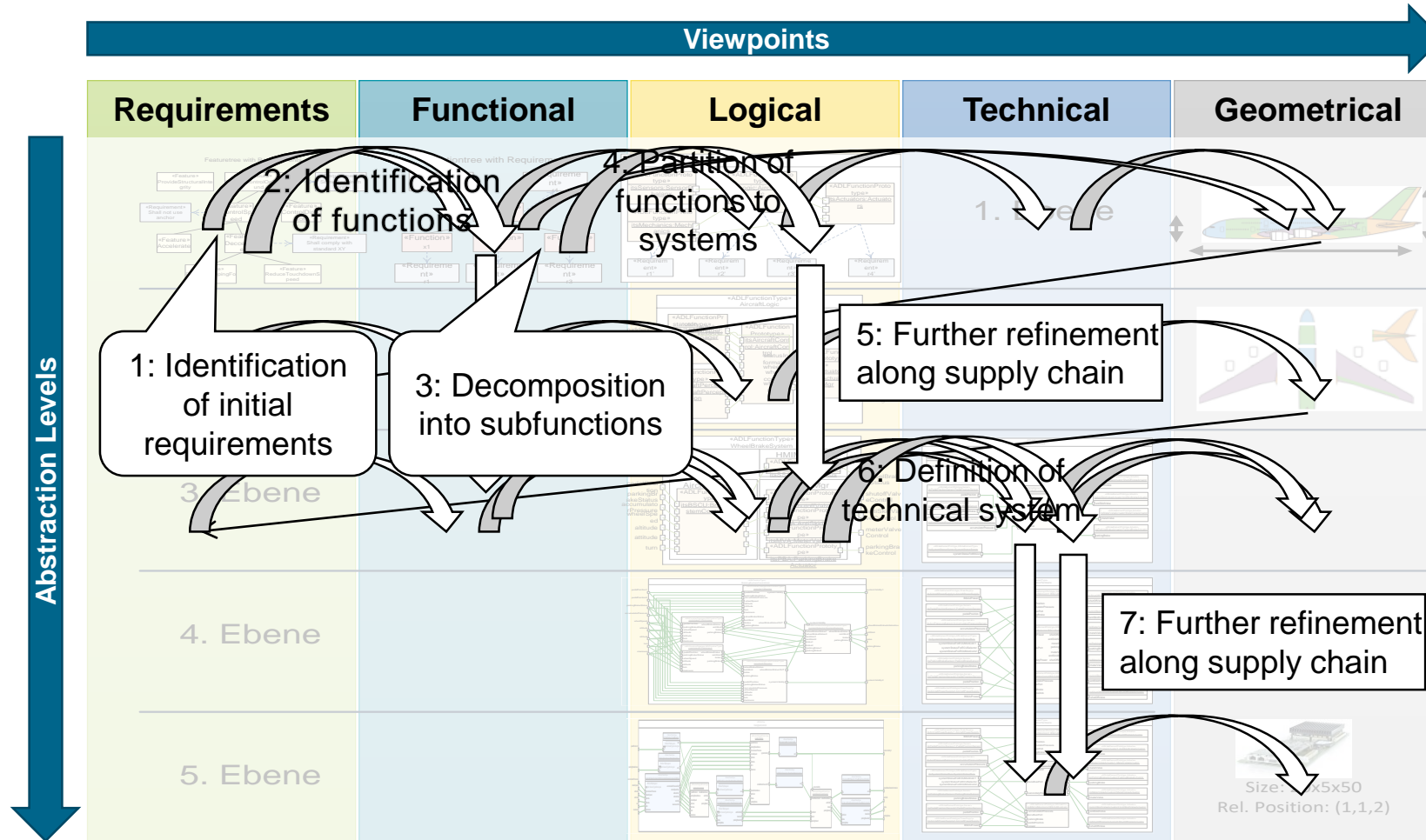
Compositional Semantic Framework

Abstraction Levels and Viewpoints



Compositional Semantic Framework

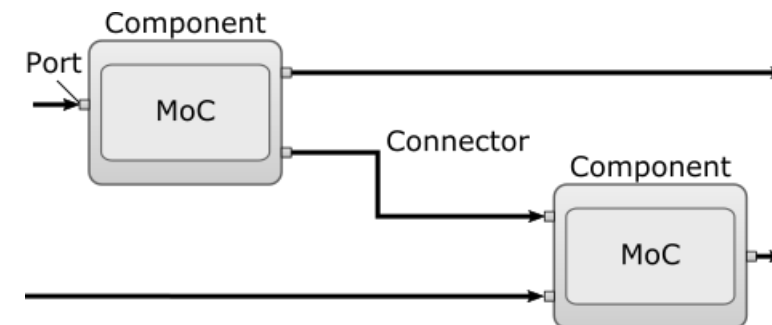
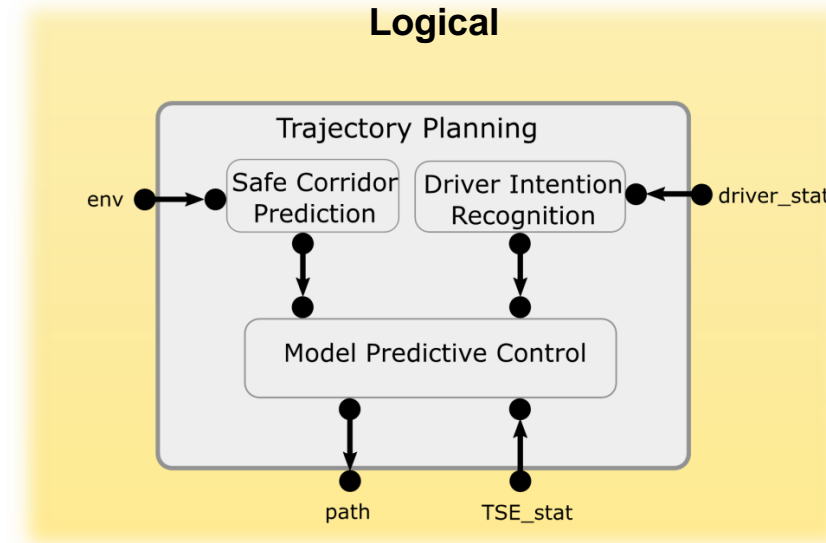
Design Steps (Examples)



Compositional Semantic Framework

Components and Hierarchy

- Components
 - Basic design entity to structure models
 - Well defined interfaces
 - Can be reused in different design contexts
- Typed Ports
 - Define syntactical interface to adjacent components / environment
- Hierarchy and Composition
 - Allows deeply nested component hierarchies
 - Supports top-down and bottom-up design
- Connectors
 - Component interaction via port connections
 - Simple and complex connectors



Compositional Semantic Framework

Component Behaviour (Semantic Domain)

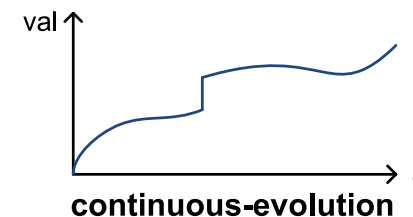
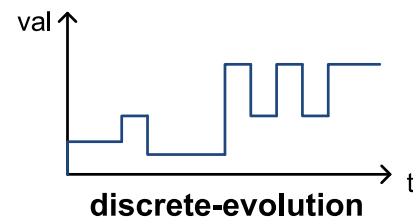
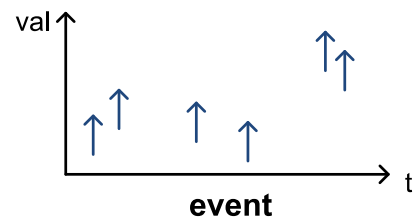


- Behaviour is visible at component ports

- Common dense time domain: $\mathbb{T} = \mathbb{R}^{\geq 0}$ (for example)
- Port types and value domains: D_p for port p
- Behaviour in term of signals: $s_p: \mathbb{T} \rightarrow V_p \cup \{\perp\}$
 - \perp means “absent value”

- Allows for specification of different “types” of behaviour:

- **Discrete event** (absent values except for discrete set of time points)
- **Discrete evolution** (changes only at discrete time points)
- **Continuous evolution**



Compositional Semantic Framework

Contracts – Assume / Guarantee Reasoning



- **Assumptions (green):**
 - Specify necessary conditions of environment and surrounding components for component to work properly
- **Guarantees (blue):**
 - Specify required behaviour that must be guaranteed by implementation if used in context compliant to assumptions
- Split in assumptions and guarantees allows **compositional** reasoning schemes
 - Implementations can be developed independently
 - Reduces verification complexity
 - Detect integration issues early in design process before developing implementations

Assumption A:

driver_stat occurs every 1s.

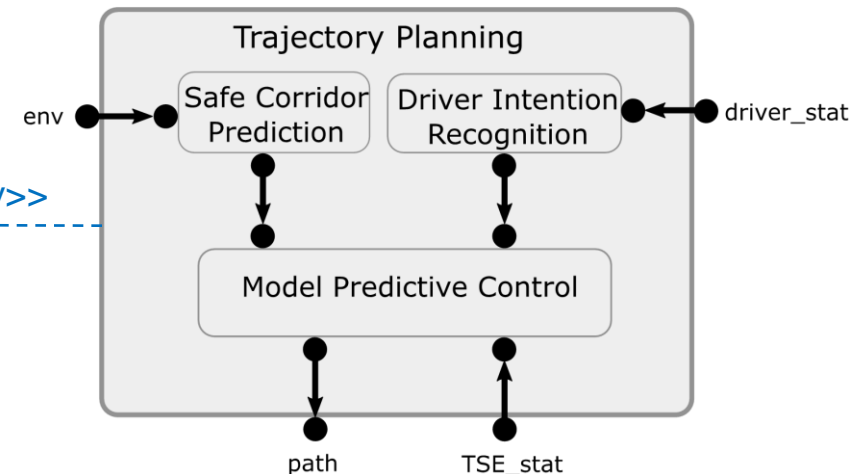
TSE_stat occurs after every path.

Guarantee G:

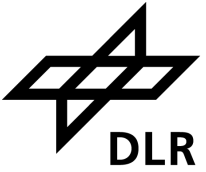
Reaction(TSE_stat.reached,path)
within [150,1500] ms.

Reaction(TSE_stat.paa,path) within
[150,200] ms.

<<satisfy>>



Agenda



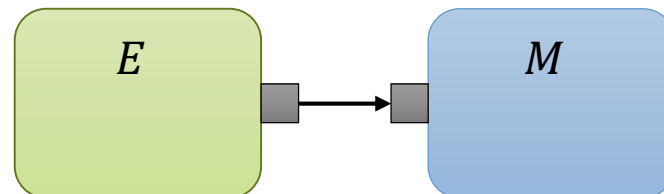
- Compositional Semantic Framework
- **A "Meta Theory" of Contract Based Design**
- Example

Contract Based Design

Components



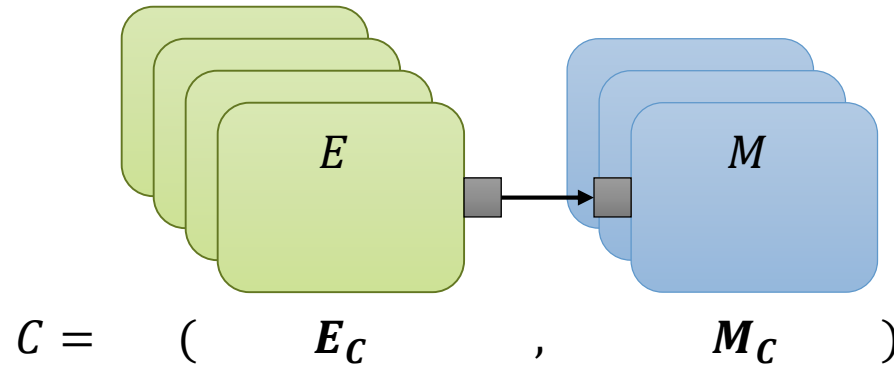
- Contract Based Design implies existence of Components
 - A *component* is the basic design entity.
 - Components can be *composed*: $M_1 \times M_2$.
 - Not all components can be composed
 - Components are *composable* if $M_1 \times M_2$ is well-defined.
 - An *environment* of M is a component E such that $E \times M$ is composable.



Contract Based Design

Contracts

- A contract $C = (E_C, M_C)$ specifies two sets of components
 - We say, a component $M \in M_C$ is an *implementation* of C : $M \models^M C \Leftrightarrow M \in M_C$
 - We say, a component $E \in E_C$ is an *environment* of C : $E \models^E C \Leftrightarrow E \in E_C$
 - Each $E \models^E C$ and $M \models^M C$ must be composable



- We say, C is *consistent* iff it has at least one implementation: $M_C \neq \emptyset$
- We say, C is *compatible* iff it has at least one environment: $E_C \neq \emptyset$

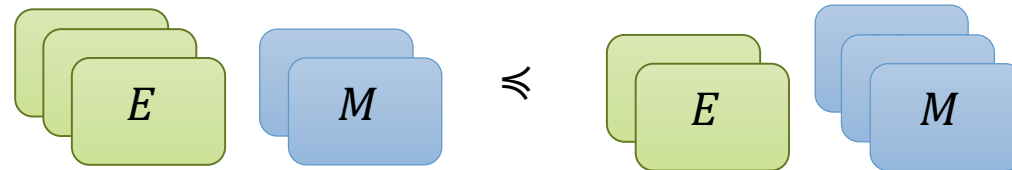
Contract Based Design

Refinement, Composition



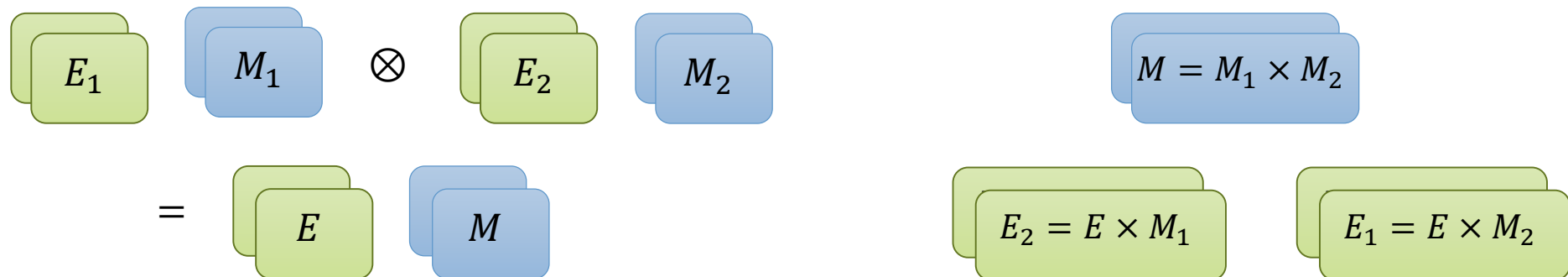
Refinement

Contract C' refines C , $C' \preceq C$, iff $E_{C'} \supseteq E_C$ and $M_{C'} \subseteq M_C$



Composition

$$C_1 \otimes C_2 = \min \left\{ C \mid \begin{array}{l} \forall M_1 \models^M C_1 \\ \forall M_2 \models^M C_2 \\ \forall E \models^E C \end{array} \Rightarrow \begin{array}{l} M_1 \times M_2 \models^M C \\ E \times M_1 \models^E C_2 \\ E \times M_2 \models^E C_1 \end{array} \right\}$$



Contract Based Design

Important Properties



■ Refinement

- Let be $C' \preceq C$

- Any implementation of C' is an implementation of C

$$M \models^M C' \implies M \models^M C$$

- Any environment of C is an environment of C'

$$E \models^E C \implies E \models^E C'$$

■ Independent Implementability

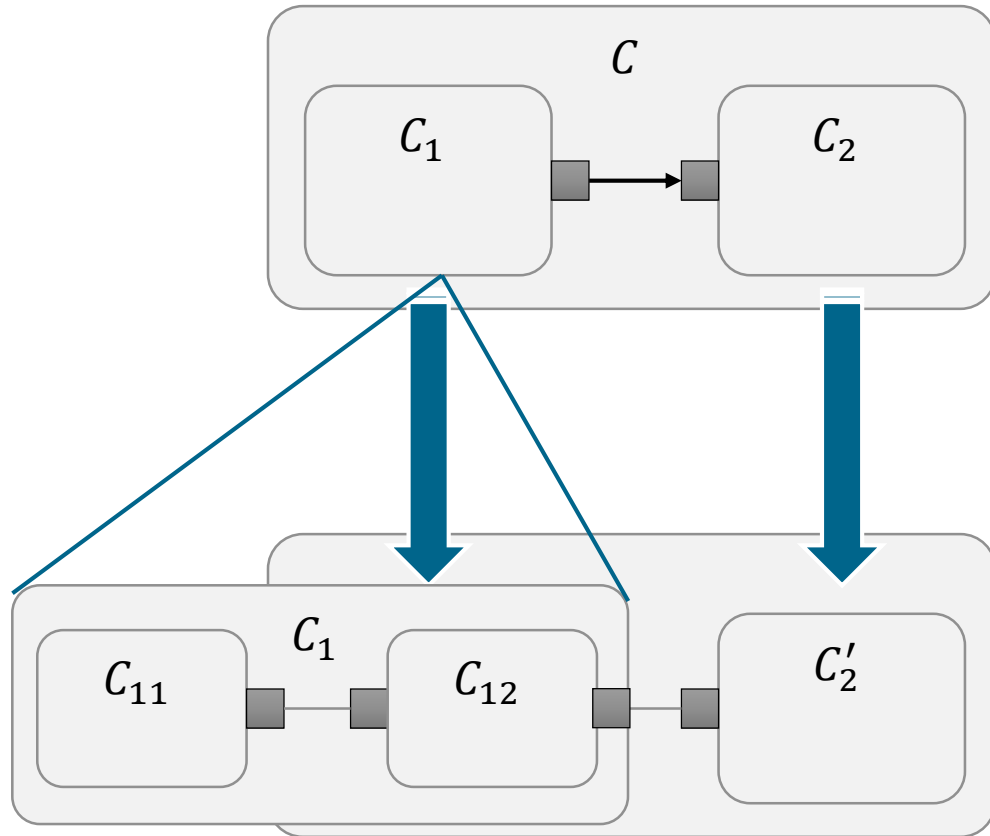
- Let be $C'_1 \preceq C_1$ and $C'_2 \preceq C_2$

- The composition of C'_1 and C'_2 refines the composition of C_1 and C_2

$$C'_1 \otimes C'_2 \preceq C_1 \otimes C_2$$

Contract Based Design

Important Properties Applied



$$C_1 \otimes C_2 \preceq C \quad \text{Virtual Integration Test}$$

$$\left. \begin{array}{l} C'_1 \preceq C_1 \\ C'_2 \preceq C_2 \end{array} \right\} \Rightarrow C'_1 \otimes C'_2 \preceq C_1 \otimes C_2 \preceq C$$

$$(C_{11} \otimes C_{12}) \otimes C'_2 \preceq C_1 \otimes C_2$$

Contract Based Design

Virtual Integration Test



Recall

- $C' \preceq C$ iff $E_{C'} \supseteq E_C$ and $M_{C'} \subseteq M_C$

$$C_1 \otimes C_2 = \min \left\{ C \mid \begin{array}{l} \left[\begin{array}{l} \forall M_1 \models^M C_1 \\ \forall M_2 \models^M C_2 \\ \forall E \models^E C \end{array} \right] \Rightarrow \left[\begin{array}{l} M_1 \times M_2 \models^M C \\ E \times M_1 \models^E C_2 \\ E \times M_2 \models^E C_1 \end{array} \right] \end{array} \right\}$$

Virtual Integration Test

- In top-down design processes, often C is given as well as C_1 and C_2 .
 - Calculating $C_1 \otimes C_2$ and checking $C_1 \otimes C_2 \preceq C$ is not necessary.
- For any contract C with

$$\begin{array}{l} \left[\begin{array}{l} \forall M_1 \models^M C_1 \\ \forall M_2 \models^M C_2 \\ \forall E \models^E C \end{array} \right] \Rightarrow \left[\begin{array}{l} M_1 \times M_2 \models^M C \\ E \times M_1 \models^E C_2 \\ E \times M_2 \models^E C_1 \end{array} \right] \end{array} \quad \text{Virtual Integration Test}$$

holds that $C_1 \otimes C_2 \preceq C$.

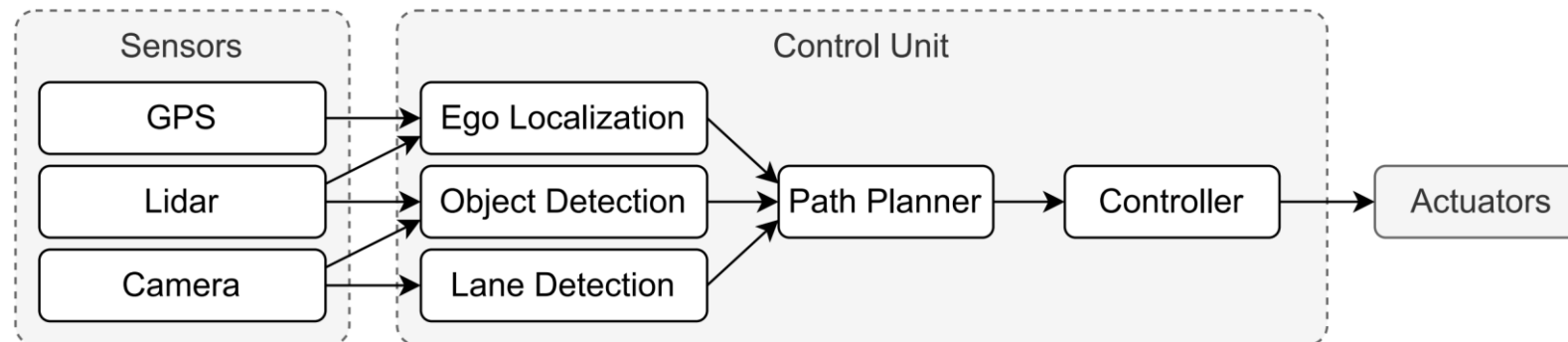
Agenda



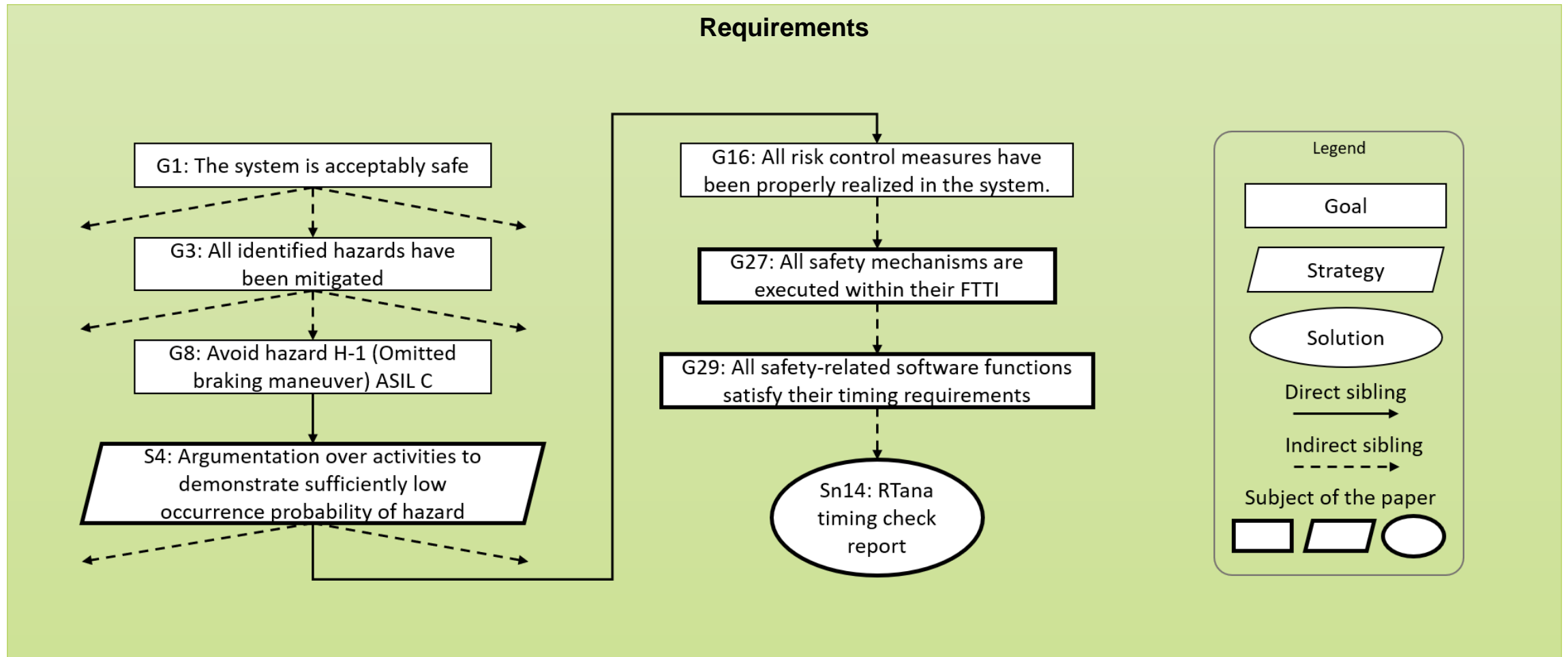
- Compositional Semantic Framework
- A "Meta Theory" of Contract Based Design
- **Example**

Example – Overview

- (Part of) a highly automated vehicle that shall perform two main functions:
 - i. following a predefined route,
 - ii. avoiding collisions with traffic participants / obstacles.
- Sensors, Control Unit, Actuators:
 - Ego Localization (localize the ego vehicle),
 - Object Detection (detect, localize, and classify traffic participants and obstacles),
 - Lane Detection (detect lane boundaries),
 - Path Planner (plan maneuvers),
 - Controller (and calculate actuator commands).



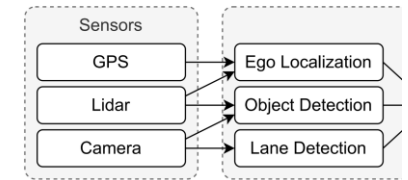
Safety Case (Excerpt)



HARA – Safety Goals – Causes – Safety Concept



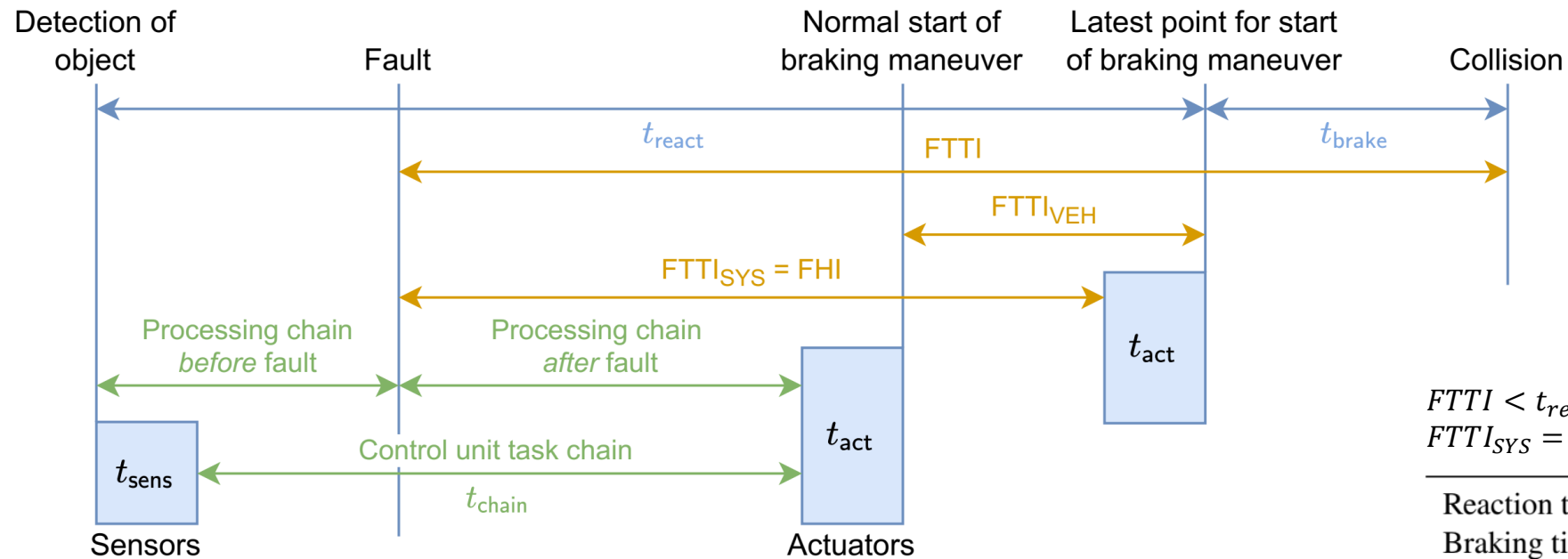
- HARA: Identification of hazards and corresponding risks
 - Here: **H-1: Omitted braking maneuver**
(a necessary braking maneuver of the ego vehicle is not performed (in time))
- Two failure modes are considered:
 1. One of *Lidar* or *Camera* fails permanently.
 2. One or both of *Lidar* and *Camera* fail for a limited amount of time.



■ Safety Requirements:

ID	Description
SG-1	The system shall prevent omitting required braking maneuvers.
SR-1-1a	The system shall use lidar and camera for object detection.
SR-1-1b	The system shall ensure that objects are detected if lidar or camera fails.
SR-1-1.1	The system shall identify sensor failures.
SR-1-1.1.1	The system shall identify when the lidar sensor has failed.
SR-1-1.1.2	The system shall identify when the camera has failed.
SR-1-1.2	The system shall mitigate sensor failures.
SR-1-1.2.1	The system shall detect objects using lidar.
SR-1-1.2.2	The system shall detect objects using camera.
SR-1-1.2.3	The system shall fuse the objects detected by lidar and camera.
SR-1-1.3	The system shall use only information from working sensors.
SR-1-2	...

Fault Tolerant Time Interval (FTTI)



$$FTTI < t_{react} + t_{brake} \approx 4.57s$$

$$FTTI_{SYS} = FHI < t_{react} + t_{act} + t_{sense} \approx 2.43s$$

Reaction time	t_{react}	2.57 s
Braking time	t_{brake}	2.0 s
Sensor processing time	t_{sense}	50 ms
Actuator time	t_{act}	100 ms
Control unit WCRT	t_{chain}	< 800 ms

Remark: We consider a *safety mechanism implemented with emergency operation* [ISO26262:2018]

Mapping Requirements to Technical Architecture

Timing Contracts



A: sd occurs every 66 ms.
 G1: Reaction(sd.OK,fused_objects) within [100,1006] ms in mode Normal.
 G2: Reaction(sd.NOT_OK,set(mode,Recover)) within [5,66] ms in mode Normal.
 G3: Reaction(sd.OK,fused_objects) within [100,940] ms in mode Recover.

$\Upsilon ?$ (VIT)

Lidar Grabber

A_{L1} : sd.lidar occurs every 66 ms.
 G_{L1} : Reaction(sd.lidar.OK,(set(mode,Normal),lidar_output)) within [5,60] ms.
 G_{L2} : Reaction(sd.lidar.NOT_OK,(set(mode,Recover))) within [5,60] ms.

Camera Grabber

A_{C1} : sd.camera occurs every 66 ms.
 G_{C1} : Reaction(sd.camera.OK,(set(mode,Normal),lidar_output)) within [7,60] ms.
 G_{C2} : Reaction(sd.camera.NOT_OK,(set(mode,Recover))) within [7,60] ms.

Object Localization via Lidar

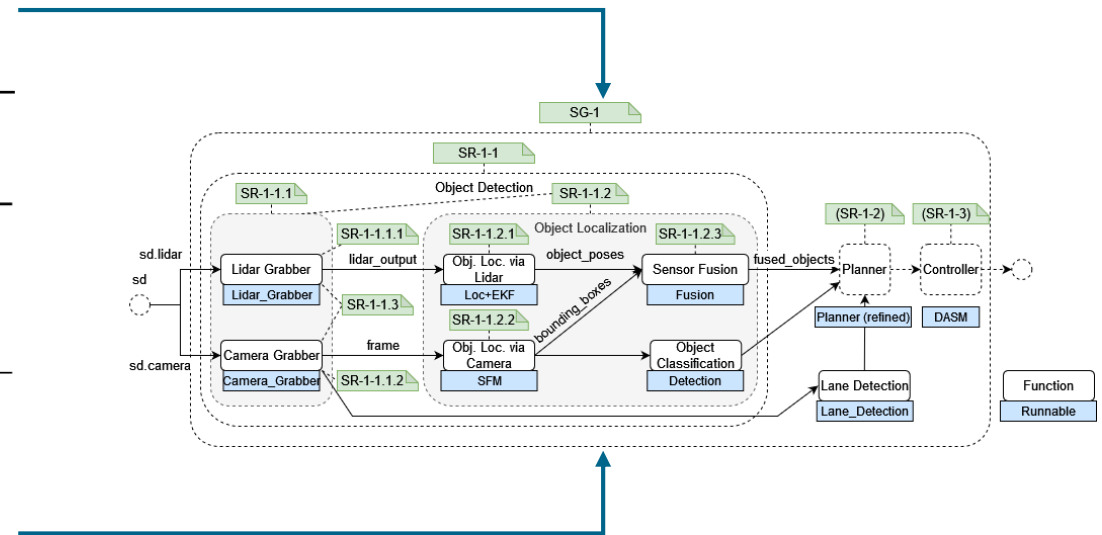
G_{OL} : Reaction(lidar_output,object_poses) within [120,795] ms.

Object Localization via Camera

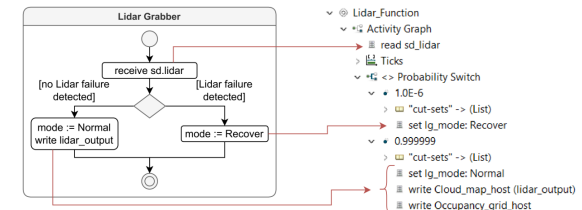
G_{OC} : Reaction(frame,bounding_boxes) within [20,55] ms.

Sensor Fusion

G_{F1} : Reaction(object_poses,fused_objects) within [5,25] ms.
 G_{F2} : Reaction(bounding_boxes,fused_objects) within [5,25] ms.



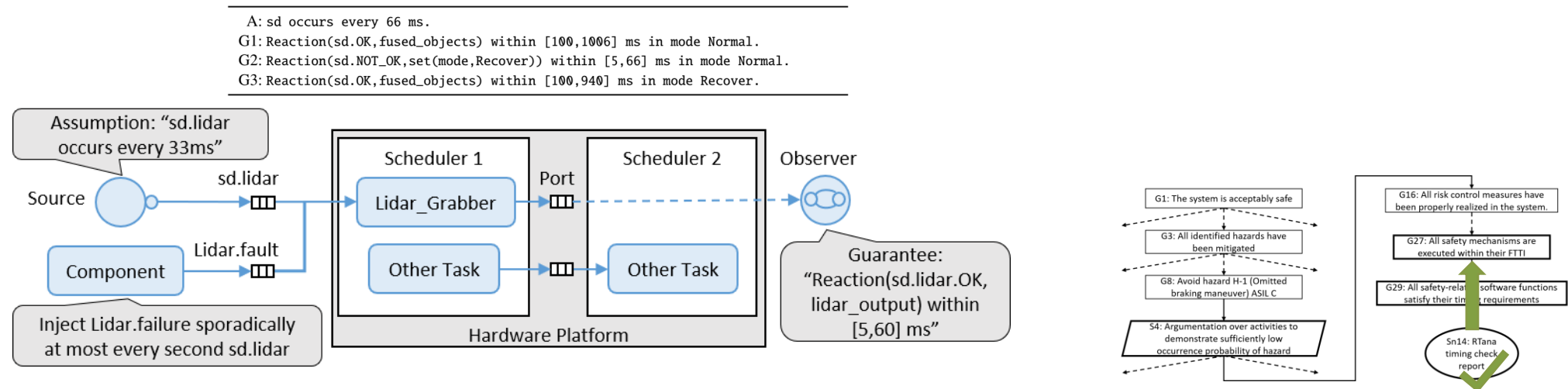
$\equiv ?$



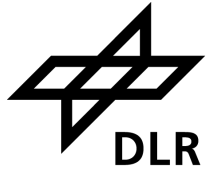
Analysis

Approach and Structure of the Rtana_{2sim} Model

- APP4MC model is translated into RTana_{2sim} model.
- Assumptions are translated into event sources.
- Guarantees are translated into observer automata.
- Additional components provide for fault injection.



Analysis Example



rtana2sim

File

Input

```
1 // Copyright (C) 2022 by DLR e. V.
2
3 const FAULT_NOT_PRESENT=1;
4 const FAULT_PRESENT=2;
5 const BLACK=1;
6 const RED=2;
7 const GrabberMode_Normal=1;
8 const GrabberMode_Recover=2;
9 ports dummy[0]:=1 OS_Overhead_trigg OS_Overhead_fin[0]:=1 Lidar_Grabber_trigg Lidar_Grabbe
10
11 signal periodic_100ms_signal -> OS_Overhead_trigg;
12 signal periodic_33ms_signal -> Lidar_Grabber_trigg PRE_SFM_gpu_POST_trigg;
13 signal periodic_5ms_signal -> DASM_trigg;
14 signal periodic_10ms_signal -> CANbus_polling_trigg;
15 signal periodic_15ms_signal -> EKF_trigg Planner_trigg;
16 signal SFM_stim_signal -> SFM_trigg;
17 signal periodic_400ms_signal -> PRE_Localization_gpu_POST_trigg;
18 signal Localization_stim_signal -> Localization_trigg;
19 signal periodic_66ms_signal -> PRE_Lane_detection_gpu_POST_trigg;
20 signal Lane_detection_stim_signal -> Lane_detection_trigg;
21 signal periodic_200ms_signal -> PRE_Detection_gpu_POST_trigg;
22 signal detection_stim_signal -> Detection_trigg;
23 source periodic_100ms [100,100] periodic_100ms_signal!1;
24 source periodic_33ms [33,33] periodic_33ms_signal!1;
25 source periodic_5ms [5,5] periodic_5ms_signal!1;
```

Log

```
Start RTana parser...
Started in Analysis Mode
Analysis: started.
Analysis: performing random unroll strategy...
Analysis: clearing state space...
Analysis: unrolling finished.
Unrolling finished with 1.002 states and queue size 433 (0.01 secs)
Analysis: checking properties...
Analysis: STS is not closed.
Analysis: property (C2_G_1.mode != BAD && C2_G_2.mode != BAD) is satisfied.
Analysis: property C2_G_2.mode != BAD is satisfied.
```

Go

mode Analysis

strategy Random

max. states [#] 1000

max. runtime [s] 60

Transitions

Ports

- dummy
- OS_Overhead_trigg
- OS_Overhead_fin
- Lidar_Grabber_trigg
- Lidar_Grabber_fin
- Lidar.transient_failure
- Lidar.permanent_failure
- Camera.permanent_failure
- Camera.transient_failure
- sd_lidar
- lg_mode
- lidar_output
- lidar_input
- sd_camera
- cg_mode
- Image_host
- frame
- Image_lane_lines_host
- DASM_trigg
- DASM_fin
- speed_objective
- steer_objective
- CANbus_polling_trigg
- CANbus_polling_fin
- Vehicle_status_host

Sources/Components

- periodic_100ms (----)
- periodic_33ms (----)
- periodic_5ms (----)
- periodic_10ms (----)
- periodic_15ms (----)
- periodic_400ms (----)
- periodic_66ms (----)
- periodic_200ms (----)
- C2_A1 (----)
- C2_A2b (----)

Core5 (Type: sched_fpp)

- DASM (----)
- OS_Overhead (----)

Core4 (Type: sched_fpp)

Lidar_Grabber

PRE_Detection_gpu_POST

CANbus_polling

EKF

Planner

PRE_SFM_gpu_POST

PRE_Lane_detection_gpu_POST

Lane_detection

PRE_Localization_gpu_POST

Localization

SFM

Show Error Trace

Conclusion



- Model-based systems engineering helps in solving many challenges in engineering processes.
- The SPES modelling framework aims at supporting engineering of (safety-critical) CPS.
- Contract-based design provides formal design and engineering support:
 - Correctness of key design steps becomes verifiable.
 - Enables tool support in verification tasks.
- Industrial example demonstrates applicability.