## Performance of numerical algorithms for low-rank tensor operations in tensor-train / matrix-product-states format

Knowledge for Tomorrow

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## Content

- Introduction
- Tensor-train / matrix-product states for data compression

- Solving linear systems in tensor-train / matrix-product-states format
- Underlying linear algebra operations
- Conclusion





## Introduction: performance engineering

• Levels of parallelism:

- single-core: ~100 flop per cycle "on the fly"
- multi-core: e.g. 128 cores (shared memory)
- multi-node: cluster of nodes (fast network)
- Data transfers vs. arithmetic operations: Computer model:
  - 1. load *n* bytes from memory
  - 2. perform k arithmetic operations (Flop)
  - 3. store *m* bytes in memory
  - $\rightarrow$  Compute intensity:  $I_c = \frac{k}{n+m}$

#### • Expected (ideal) runtime:

• memory-bound ( $I_c < 16$ ): t = (n + m)/bandwidth

• compute-bound ( $I_c > 16$ ):





compute intensity [Flop/Byte]

## Introduction: linear algebra

#### Matrix-matrix product (GEMM):

 $\begin{array}{ccc} C &\leftarrow A & B \\ (n \times k) & (n \times m) \ (m \times k) \end{array}$ 

- 2nmk flops, 8(nk + nm + mk) data transfers
  - compute bound for  $n \approx m \approx k \gg 100$
  - memory bound for  $\min(n, m, k) \leq 100$

#### QR decomposition:

QR = A, with  $Q^TQ = I$ , and R upper triangular.

- $O(nm^2)$  flops, O(nm) data transfers
  - memory bound for  $m \lesssim 100$
  - → tall-skinny QR (TSQR)

#### Singular Value Decomposition (SVD):

 $A = USV^{T},$ with  $U^{T}U = I, V^{T}V = I$ , and  $S = \text{diag}(\sigma_{1}, ..., \sigma_{m}),$  $\sigma_{1} \ge \sigma_{2} \ge \cdots \ge \sigma_{m} \ge 0.$ 

•  $O(nm^2)$  flops (usually: runtime  $t_{SVD} \gg t_{QR} > t_{GEMM}$ )

• Truncated SVD 
$$\rightarrow$$
 best rank-r approximation:  
 $\|A - \hat{B}\|_F = \min_{\operatorname{rank}(B) \leq r} \|A - B\|_F$   
with  $\hat{B} = U \operatorname{diag}(\sigma_1, \dots, \sigma_r, 0, \dots, 0) V^T$ 



## **Tensor network notation**

- Tensor = multi-dimensional array:  $X \in \mathbb{R}^{n_1 \times n_2 \times \cdots \times n_d}$
- Contraction: sum over product of entries

   → generalization of matrix-matrix product
   example (contraction of 3<sup>rd</sup> dim. of X and 2<sup>nd</sup> of Y)

$$Z_{i_1,i_2,i_4,j_1,j_3} = \sum_{i_3} X_{i_1,i_2,\mathbf{i_3},i_4} Y_{j_1,\mathbf{i_3},j_3}$$

- In 2d: all contractions with transpose+GEMM:  $AB, A^{T}B, AB^{T}, A^{T}B^{T}$
- > 2d: too many combinations!
   → tensor network notation

#### Tensor network notation (from physics!):



Matrix with orthogonal columns:





#### **Tensor network notation (cont.)**

#### **Contractions:**



**Decompositions:** 

ranks /

bond-dimensions

 $n_5$ 

n<sub>4</sub>

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## Data compression with TT-SVD: motivation

- Input: large, dense, high-dimensional data
- Example: Cardio MRT videos:
  - Multiple patients
  - Multiple views/cuts
  - Slices in z direction
  - Finer resolution in x/y direction
  - Videos over time
- Data looks very "similar" in all directions
- Goals:
  - 1. Compress data in tensor-train format
  - 2. Data analysis with compressed data





## **Truncated SVD in 2d**

- Interpret picture as matrix
- Approximate matrix by lower rank:



(	(	
	Data size	
	uncompressed:	160 kB
	compressed:	1.6 <i>r</i> kB

#### rank 1 approximation





## **Truncated SVD over time**

- We can do the same for a movie over time...
- Interpret movie as matrix (time x pixels)
- Approximate matrix by lower rank:







## Truncate in 2d and over time $\rightarrow$ tensor-train / MPS

- We can combine both ideas...
- First truncate over time, then in x/y direction:







# $\begin{array}{c|c} r & r \\ t & y \\ t & y \\ \end{array}$

## **Tensor-train decomposition / MPS for data compression**





## **Performance of the tensor-train SVD**

**Setup:** given dense  $X \in \mathbb{R}^{n^d}$ , desired rank  $r_{max}$ 

#### Core idea:

 neat combination of (tall-skinny) matrix operations (TSQR, SVD, fused GEMM+reshape)

#### **Results:**

- Memory bound for small r<sub>max</sub> (transfer volume ~2.2n<sup>d</sup> values)
- Compute bound for  $r_{\text{max}} \gtrsim 100$ (operations  $\leq 12n^d r_{\text{max}}$  Flops)
- Common implementations are too slow (factor  $\geq 50$ )

Röhrig-Zöllner et.al.: Performance of the low-rank TT-SVD for large dense tensors on modern multi-core CPUs, SISC, 2022 (doi: 10.1137/21m1395545)



TT-SVD of a 2<sup>27</sup> tensor on a 14-core Intel Skylake Gold 6132



## **TT-SVD** for a real-world data set

Cardio-MRT data (subset):  $19 \times 19 \times 300 \times 200 \times 200$ patients slices frames x-dir. y-dir. (complete data set: more patients × different views)

Data size	
Uncompressed:	8.7 GB (uint16)
	17.3 GB (float)
Compressed:	124 MB
$\epsilon = 0.002$	



TensorTrain\_float
with dimensions = [19, 19, 5, 5, 3, 2, 2, 2, 2, 2, 5,
5, 2, 2, 2, 5, 5]
ranks = [19, 361, 711, 1336, 1565, 1267,
861, 943, 1177, 1382, 693, 187, 100, 50, 25, 5]



## **Example for data analysis step in TT format: extract distances**

- Given: vectors  $(w_1, ..., w_n) =: W, w_i \in \mathbb{R}^m$
- Calculate the pair-wise distances of all vectors:  $d_{ij}^{2} = \left\|w_{i} - w_{j}\right\|_{2}^{2} = \left\|w_{i}\right\|_{2}^{2} - 2\langle w_{i}, w_{j} \rangle + \left\|w_{j}\right\|_{2}^{2}$
- Or their cosine similarities:  $s_{ij} = \frac{\langle w_i, w_j \rangle}{\|w_i\|_2 \|w_j\|_2}$
- Both require  $O(mn^2)$  operations
- Alternative approach:
  - Compress in TT format: O(mnr) operations
  - Rearrange sub-tensors
    - → approximation for cosine similarities





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## Linear algebra in tensor-train / matrix-product-states format



## Linear solvers in TT format: motivation

Solve large-scale linear systems infeasible to store solution as a "dense" vector (e.g., operator dimension:  $\mathbb{R}^{100^{10} \times 100^{10}}$ )

#### Possible applications:

- High-dim. / parameterized / stochastic PDEs
- Large-scale optimal control problems
- Problems from quantum physics: often Hermitian eigenvalue problems (closely related to linear systems but not discussed here!)
- Data science: e.g., compressed sensing

#### **Problem definition**

Given:

- Low-rank linear operator:  $\mathcal{A}_{TT} \in \mathbb{R}^{n^d} \to \mathbb{R}^{n^d}$
- Low-rank right-hand side:  $B_{TT} \in \mathbb{R}^{n^d}$
- Desired tolerance:  $\epsilon_{tol}$

#### Calculate:

- Iterative solution  $X_{TT}^*$  with  $\|\mathcal{A}_{TT}X_{TT}^* - B_{TT}\|_* \le \epsilon_{tol}$
- (choice of norm  $\|\cdot\|_*$  depends on solution method)

Focus here on  $n \gg 2$ , e.g. dimensions  $50^{10}$  (performance characteristics differ for  $2^{N}$ !)



## Common algorithms: TT-GMRES, TT-MALS, TT-AMEn

#### Methods:

- "Global": TT-GMRES GMRES applied in TT format with additional truncations
- "Local" projection onto 2 sub-tensors: TT-MALS (like DMRG but formulated for generic linear systems)
- "Local" projection onto 1 sub-tensor: TT-AMEn (subspace "enriched" with projected residuals)



(simple PDE,  $20^6 \le n^d \le 100^{14}$ , r < 100,  $\epsilon_{tol} = 10^{-8}$ ):

- TT-GMRES about 100x slower than TT-MALS
- TT-MALS about 100x slower than TT-AMen



## Preconditioning

Common practice for **sparse solvers**...

**Required properties:** 

- 1.  $\operatorname{cond}(PA) \ll \operatorname{cond}(A)$
- 2. cheap  $y \leftarrow Px$  (apply precond. to vector)

Different variants:

- left preconditioning: PAx = Pb
- right preconditioning: APy = b, x = Py
- two-sided precond.:  $P_LAP_Ry = P_Lb$ ,  $x = P_Ly$

Need a few more constraints for **TT solvers**!

Desired properties ("global" preconditioner):

- 1.  $\operatorname{cond}(\mathcal{PA}_{TT}) \ll \operatorname{cond}(\mathcal{A}_{TT})$ (fewer total (inner) iterations)
- 2. rank( $\mathcal{PA}_{TT}$ )  $\approx$  rank( $\mathcal{A}_{TT}$ ) (complexity is cubic in the rank)
- 3. "make the operator more symmetric" (better convergence / possibly smaller  $r_{max}$ )
- 4. "preserve problem structure"



## Suggestion: simple rank-1 preconditioner

Idea:

- approximate TT-operator with rank 1:  $\tilde{\mathcal{A}}_{TT} \approx \mathcal{A}_{TT}$  with rank $(\tilde{\mathcal{A}}_{TT}) = 1$
- Rank-1 inverse:  $(\tilde{A}_1 \otimes \tilde{A}_2 \otimes \cdots \otimes \tilde{A}_d)^{-1} = \tilde{A}_1^{-1} \otimes \tilde{A}_2^{-1} \otimes \cdots \otimes \tilde{A}_d^{-1}$

Two-sided preconditioned operator (for symm. problems  $\mathcal{L}_{TT}^T = \mathcal{R}_{TT}$ ):

 $\mathcal{L}_{TT}\mathcal{A}_{TT}\mathcal{R}_{TT}$ 

using SVDs  $\tilde{A}_k = U_k S_k V_k^T$ :

$$L_k = S_k^{-1/2} U_k^T$$
,  $R_k = V_k S_k^{-1/2}$ 

Generic, fast and works well in my test cases.

Replace by problem-specific preconditioner if possible!

 $\rightarrow$  Combines properties 1, 2, 3 but not 4.



## **TT-MALS** projection

#### Idea:

- Vary only  $(X_k, X_{k+1})$ (keeping  $X_1, \dots, X_{k-1}, X_{k+2}, \dots, X_d$  fixed)
- Minimize energy:
  - $J(X_{TT}) \coloneqq 0.5 \langle X_{TT}, \mathcal{A}_{TT} X_{TT} \rangle \langle X_{TT}, B_{TT} \rangle$
- Sweep over dimensions  $(k \leftarrow k \pm 1)$

#### **Properties:**

•  $\mathcal{V}^T \mathcal{V} = I$  with  $\mathcal{V} y_{TT} = X_{TT}$ 

For sym. pos. def. operator  $A_{TT}$ :

- Minimizes  $||X_{TT} X_{TT}^*||_{\mathcal{A}_{TT}}$
- $\operatorname{cond}(\mathcal{V}^T \mathcal{A}_{TT} \mathcal{V}) \leq \operatorname{cond}(\mathcal{A}_{TT})$

#### Resulting "local" problem





## Idea for non-symmetric projection

Sym. projection sub-optimal for non-sym. operator!  $\rightarrow$  use  $\mathcal{W}^T \mathcal{A}_{TT} \mathcal{V}$  with  $\mathcal{W} \neq \mathcal{V}$ 

Idea:

• Try to build  $\mathcal{W}$  to span directions of  $\mathcal{A}_{TT}\mathcal{V}$ 

#### **Properties:**

- $\mathcal{W}^T \mathcal{W} = I$  with  $\mathcal{A}_{TT} \mathcal{V} \approx \mathcal{W} M$
- Solution  $\mathcal{W}^T \mathcal{A}_{TT} \mathcal{V}_{TT} = \mathcal{W}^T B_{TT}$  approximates:  $\min_{\mathcal{Y}_{TT}} \|\mathcal{A}_{TT} \mathcal{V}_{TT} - B_{TT}\|_F$
- $W_l, W_r$  chosen to make  $\mathcal{W}^T \mathcal{A}_{TT} \mathcal{V}$  more normal

#### Derivation of the "local" operator





## **Results with non-symmetric projection**

Setup:

- Simple PDE (dimensions 20<sup>10</sup>)
- TT-MALS with inner TT-GMRES
- Varying asymmetry (conv. to diff. ratio)

Observation:

 Alternative projection beneficial for strongly non-symmetric problems





## **TT-AMEn** performance

**Setup:** simple PDE (conv.-diff. ratio = 10,  $\epsilon_{tol} = 10^{-8}$ )

#### Core ideas:

- Simple preconditioner
- Improved orthogonalization and SVD steps (with TSQR)
- Faster contractions (dim. reordering + padding)

#### **Results:**

- Not 1 dominating part in the algorithm
   → needs combination of improvements!
- Significant speedup (factor  $\sim$ 5)

#### **Remark:**

• Tweaked ttpy version with fast underlying BLAS (MKL)



TT-AMEn for a 50<sup>10</sup> problem on a 64-core AMD EPYC 7773X



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## **Underlying linear algebra operations: motivation**

High speedups: 50x (TT-SVD), 5x (TT-AMEn) (on the same hardware!)

Fair comparisons:

- No comparison vs. "unoptimized" code!  $\rightarrow$  all implementations call BLAS/LAPACK
- Use the same (multi-threaded) BLAS/LAPACK library (MKL with workaround for AMD)
- Except for some specialized operations...
  - Q-less tall-skinny QR ("TSQR" that only returns R)
  - Fused tall-skinny GEMM+reshape
  - Fused axpy+dot

So what did I change?

- Improve mapping of high-level operations to "building blocks"
- Exploit specialized operations (less generic/accurate than BLAS/LAPACK for all inputs)
- Improve data layout (BLAS/LAPACK must work with what it gets...)





## Building blocks of one TT-SVD "step"

Standard implementation (large SVD for each step): Given tall-skinny  $X \in \mathbb{R}^{n \times k}$ , calculate:

$$\begin{split} X &= USV^{T}, \\ Q &\leftarrow V_{:,1:r}, \\ B &\leftarrow U_{:,1:r}S_{1:r,1:r}, \\ X' &\leftarrow \text{reshape}(B, \dots) \end{split}$$

Actual problem: calculate *Y*, *Q* with  $Q^T Q = I$ :  $||X - BQ^T||_F \le \tau$ ,  $X' \leftarrow \operatorname{reshape}(B, ...)$ 

Optimized implementation:

X = QR,  $R = USV^{T},$   $Q \leftarrow V_{:,1:r},$  $X' \leftarrow \text{reshape}(XQ, ...)$ 

Underlying operations:

• SVD  $(n \times k)$ 

• copy  $(k \times r)$ 

• *r* axpy (*n*)

• reshape  $(n \times r)$ 

Underlying operations:

- Q-less TSQR  $(n \times k)$
- SVD  $(k \times k)$
- copy  $(k \times r)$
- tall-skinny GEMM+reshape  $(n \times k \cdot k \times r)$

## Linear solver building blocks: QRs and SVDs

Optimizations assume tallskinny / very rectangular matrices!

#### Orthogonalization

Given  $X = X_1 X_2$ , calculate:  $X_1 = QB$ ,  $X'_1 \leftarrow Q$ ,  $X'_2 \leftarrow BX_2$ 

Standard: pivoted Householder QR

Optimized with Q-less TSQR:

$$\begin{array}{l} X_1 = QR, \\ X_1' \leftarrow X_1 R^{-1}, \\ X_2' \leftarrow RX_2 \end{array}$$

but  $X'_1$  inaccurate for  $cond(R) \gg 1$ 

#### Truncation

Given 
$$X = X_1 X_2$$
 with  $X_2^T X_2 = I$ , calculate:  
 $\|X_1 - QB\|_F < \tau$ ,  
 $X'_1 = Q$ ,  
 $X'_2 = BX_2$ 

Standard: truncated SVD

Optimized with Q-less TSQR:  $X_1 = QR,$   $R \approx USV^T,$   $X'_1 = X_1VS^{-1},$   $X'_2 = SV^TX_2$ Again less orthogonal  $X'_1$  but product  $X_1X_2$  ok!



## **QRs+SVDs** in **TT/MPS** addition+truncation

High-level operation:  $Z_{TT} \approx X_{TT} + Y_{TT}$ Setup: dim. 50<sup>10</sup>,  $r_X = 50$ ,  $r_Y = 1, ..., 700$ 

#### Background:

- Combines QR- / SVD-steps
- Additional optimization for previously orthog.  $X_{TT}$  or  $Y_{TT}$  $\rightarrow$  reuse orthogonal columns



TT-AXPY+TRUNC for 50<sup>10</sup> TTs on a 64-core AMD EPYC 7773X



#### Linear solver building blocks: contractions

Multiply TT operator (MPO) with dense array

- Easily leads to complicated array accesses
- Freedom in memory-layout and padding! (operator prepared once and applied often)

#### **Optimizations:**

- Reorder and combine dimensions
   (big 1<sup>st</sup> dim./ e.g. ∑<sub>i,j</sub> A<sub>:,i,j</sub> B<sub>:,:,i,j</sub> instead of ∑<sub>i,j</sub> A<sub>:,i,j</sub> B<sub>i,:,i,j</sub>)
- Pad 1<sup>st</sup> dim. (introduces zeros the dense array!) (to avoid cache thrashing)





## Conclusion

#### **TT-SVD** compression of large dense data:

- Common implementations are about >50x too slow
- → Allows extracting interesting quantities for data analysis

#### Linear solvers in TT / MPS format:

- Obtained ~5x speedup over the standard implementation
- Presented some ideas on numerical aspects
  - Generic rank-1 preconditioner
  - Non-symmetric projection
- → Allow solving really high-dimensional linear systems!

Generic optimizations\* for building blocks of tensor-network algorithms.

\*mostly for very non-square matrix operations

#### Ideas for future work?