

1 **Exit lane allocation with the lane-based signal optimization method**
2 **at isolated intersections**

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12 **Abstract**

13 In signal optimization problems, incompatible movements usually are in either of two states:
14 predecessor or successor. However, if the exit lane is well allocated, the incompatible movements
15 merging at the same destination arm can exist in parallel. The corresponding longer green
16 duration is expected to increase the capacity of intersection. This paper aims to solve the
17 exit lane allocation problem with the lane-based method by applying the three states among
18 incompatible movements at conventional signalized intersections. After introducing auxiliary
19 variables, the problem is formulated as a mixed integer programming and can be solved using a
20 standard branch-and-cut algorithm. In addition to the exit lane allocation results, this proposed
21 method can also determine the cycle length, green duration, start of green and signal sequence.
22 The results show that the proposed method can obtain a higher capacity than that without the
23 exit lane allocation. The pavement markings are further suggested for safety.

24 **Key words:** signal optimization, exit lane allocation, lane-based method, signal timing
25 plan

26 1 Introduction

27 Signal optimization is one of solutions to increase traffic capacities in urban, and the methods
28 of signal optimization thereby attract research interests. The traditional stage-based method,
29 in which a stage is a combination of non-conflict movements having the right of way at the
30 same green time, was developed decades years ago. The widely-used objective of the stage-
31 based method is to minimize the total delay at the intersection (e.g. Webster (1958)), but the
32 multiobjective problem which combines efficiency and safety is also solved in recent years (Li
33 and Sun 2018, 2019). The stage-based method determines the green time (Webster 1958; HBS
34 2001; Ceylan and Bell 2004), or green split (Al-Khalili 1985), cycle length (Webster 1958; HBS
35 2001; Ceylan and Bell 2004) or even stage sequence (Memoli et al. 2017; Tang and Friedrich
36 2018) for given stages in a signal cycle, which means that the stage composition should be
37 done before signal timing optimization. To integrate to stage composition and signal timing
38 optimization, researchers developed the group-based method (Improta and Cantarella 1984;
39 Gallivan and Heydecker 1988; Silcock 1997), in which the given lane markings can be flexibly
40 "grouped" and their signal timing can be optimized in one mathematical model. However,
41 both the stage-based method and the group-based method require lane markings as exogenous
42 inputs whereas the lane markings could not be always available.

43 For the purpose of flexibly handling the lane markings, the lane-based method is developed
44 by maximizing the reserved capacity or minimizing the cycle length or the total delay at isolated
45 intersections (Wong and Wong 2003; Wong, Wong and Tong 2006). Wong and Heydecker (2011)
46 then extend the lane-based method via relaxing the numbers of approach lane in traffic arms

47 so that the number of approaching lanes and the exit lanes can be optimized. To better
48 handle the fluctuation of traffic demands, Zhao et al. (2013), Alhajyaseen et al. (2017) and
49 Assi and Ratrout (2018) solved a dynamic lane assignment problem which can automatically
50 determine lane markings based on varied demands, but focus on approaching lanes. To explore
51 the potential application of the lane-based method in the networks, Lee, Wong and Li (2015)
52 and Lee and Wong (2017) estimate the queue length for the intersections with signal control
53 solved by the lane-based method. Meanwhile, the lane-based method is applied for the signal
54 design of unconventional intersections. Signalized roundabouts (Ma et al. 2013), displaced left
55 turn intersections (Zhao et al. 2015), special width approach (Zhao, Liu and Wang 2016) and
56 lane dynamical exclusive bus lane design at intersections (Zhao and Zhou 2018) well adjust the
57 lane-based method into different unconventional intersection designs.

58 The signal sequence determination in a one-step model is another advantage additional to
59 the lane markings. Signal sequence is affected by the conflict matrix indicating the compati-
60 bility between two movements. The compatible movements do not mutually conflict, whereas
61 the incompatible movements do. The incompatibility occurs either at intersections or in the
62 arms where two movements merge at the same destination arms. Thus, for the safety reasons,
63 the incompatible movements must not be in the same green duration to avoid the conflicts.
64 That is, a movement can only be the predecessor or successor of its incompatible movements
65 in a signal cycle, and once it is the predecessor of the incompatible movements, it is not the
66 successor (Wong and Wong 2003). It may result in the possible inefficient utilization of inter-
67 section capacities. The conflicts in the destination arms can be eliminated by appropriately
68 allocating exit lanes. In Figure 1, the incompatible movements to the same destination arm can

69 be assigned to different exit lanes so that the conflicts are eliminated. Then the green duration
70 of the related movements could increase so as their capacities. Observing this phenomenon,
71 Xie and Jiang (2016) extend the method of Wong and Heydecker (2011)'s by allocating the exit
72 lanes to turning movements. They group the incompatible movements into strictly incompat-
73 ible movements referring to the movements going to different destination arms, and potential
74 incompatible movements referring to the movements going to the same destination arm. They
75 draw the conclusion that exit lane allocation can increase the intersection capacity.

76 [Figure 1 about here.]

77 When one determines the signal sequence, the incompatible movements can usually be
78 in one of two states: predecessor or successor (e.g. Wong and Wong (2003); Xie and Jiang
79 (2016)). However, the incompatible movements merging at the same destination arm could be
80 in parallel by appropriately allocating the exit lanes. This means, in the previous research,
81 the signal sequence states, predecessor or successor, may not be consistent with the actual
82 situation due to the feasibility of "in parallel". Further, the compatible movements could also
83 be in the one of three states. Thus, the signal sequence variables, which can describe three
84 states referring to predecessor, in parallel or successor rather than the binary states referring
85 to predecessor or successor as Xie and Jiang (2016) did, can more precisely reflect the states of
86 signal sequences, especially for the incompatible movements merging at the same destination.

87 In this paper, we propose a new method to solve the exit lane allocation problem in the
88 lane-based method to maximize capacities for intersections. The proposed optimization method
89 is developed based on Wong and Wong (2003) which can determine the lane markings of
90 the approaching lanes, the green duration for each lane and for each movement, starts of

91 green for each lane and for each movement, cycle length, signal sequence, assigned flows and
 92 reserved capacity. Additional to these decision variables, the exit lane permission indicators
 93 are introduced to determine the exit lane allocation. However, the signal sequence variables
 94 with three states make the modeling linearization be a challenge. To linearize the model,
 95 we introduce two auxiliary binary variables which are explained in the next section in detail.
 96 Therefore, the proposed model is formulated as a mixed integer linear programming such that
 97 can be efficiently solved with a standard algorithm such as a branch-and-cut algorithm with
 98 ILOG CPLEX. Taking maximization of the reserved capacity as the objective function, one
 99 can observe the capacity increase by allocating the exit lanes compared with the results of
 100 Wong and Wong (2003). Considering the potential conflicts at merging movements, we further
 101 suggest the pavement markings to ensure safety.

102 **2 Problem formulation**

103 **2.1 Intersection representation**

104 An isolated intersection has N_A arms. In each arm there are L_i approaching lanes and E_i
 105 exit lanes, where i is the arm index and $i = 1, \dots, N_A$. The turning movements are belong to
 106 the direction set M and M contains the elements RT, TH, LT where RT is right turn, TH is
 107 through movement, and LT is left turn. In this paper, U-turn is not considered. A movement
 108 can be represented as $(i, j), \forall i = 1, \dots, N_A, j \in M$ which means a turning direction j in arm i .
 109 Meanwhile, a movement can also be represented as a movement from arm $i, i = 1, \dots, N_A$ to arm
 110 $i', i' = 1, \dots, N_A$ and $i \neq i'$. Thus, a relation between the turning direction and the destination

111 arm must hold: $i' = \Gamma(i, j)$. This function means that i' is the destination arm of movement
 112 (i, j) . The details of the arm index, the approaching lane index and the exit lane index can be
 113 found in Figure 2.

114 [Figure 2 about here.]

115 2.2 Decision variables

116 The proposed model can optimize lane markings for approaching lanes, exit lane allocation,
 117 signal sequences and signal timing. The decision variables are thereby relevant to these imple-
 118 mentation requirements.

119 Approaching lane permission indicators, $\delta_{i,j,k}$, indicate whether the lane marking of the
 120 movement (i, j) is drawn on the approaching lane k . For all $i = 1, \dots, N_A, j \in M, k = 1, \dots, L_i$,
 121 if $\delta_{i,j,k} = 1$, the lane marking of the movement (i, j) is permitted on the approaching lane k ;
 122 if $\delta_{i,j,k} = 0$, otherwise. Similarly, exit lane permission indicators, $\epsilon_{i,i',k'}$, indicate whether a
 123 movement from arm i to arm i' is permitted to exit the intersection via the exit lane k' . For all
 124 $i = 1, \dots, N_A, i' = 1, \dots, N_A, i' \neq i, k' = 1, \dots, E_{i'}$, if $\epsilon_{i,i',k'} = 1$, the movement is permitted on the
 125 exit lane k' ; if $\epsilon_{i,i',k'} = 0$, otherwise. Sequence indicator represents the sequence relationship
 126 between different movements. Sequence indicators between movement (i, j) and (l, m) are
 127 denoted as $\Omega_{i,j,l,m}$ where $i, l = 1, \dots, N_A, j, m \in M$:

$$\Omega_{i,j,l,m} = \begin{cases} 1, & \text{if } (i, j) \text{ is the predecessor of } (l, m) \\ 0, & \text{if } (i, j) \text{ and } (l, m) \text{ are in parallel} \\ -1, & \text{if } (i, j) \text{ is the successor of } (l, m) \end{cases} \quad (1)$$

128 To linearize the model, we introduce two auxiliary binary variables $x_{i,j,l,m}, y_{i,j,l,m}, \forall i, l =$
129 $1, \dots, N_A, j, m \in M$. If $\Omega_{i,j,l,m} \geq 1, x_{i,j,l,m} = 0$; $\Omega_{i,j,l,m} < 1, x_{i,j,l,m} = 1$. If $\Omega_{i,j,l,m} \geq 0, y_{i,j,l,m} = 1$;
130 $\Omega_{i,j,l,m} < 0, y_{i,j,l,m} = 0$. Thus, if $x_{i,j,l,m} = 0, y_{i,j,l,m} = 1$, the movement (i, j) is the predecessor of
131 movement (l, m) ; if $x_{i,j,l,m} = 1, y_{i,j,l,m} = 1$, the movement (i, j) and the movement (l, m) are in
132 parallel; if $x_{i,j,l,m} = 1, y_{i,j,l,m} = 0$, the movement (i, j) is the successor of the movement (l, m) .
133 Summarily, $\forall i, l = 1, \dots, N_A, j, m \in M$, the following relations hold for the sequence indicators
134 and the auxiliary variables:

$$0 \leq 1 - \Omega_{i,j,l,m} \leq Hx_{i,j,l,m}, \quad (2)$$

$$-1 - H(1 - x_{i,j,l,m}) \leq \Omega_{i,j,l,m} \leq H(1 - x_{i,j,l,m}), \quad (3)$$

$$0 \leq 1 + \Omega_{i,j,l,m} \leq Hy_{i,j,l,m}, \quad (4)$$

$$-H(1 - y_{i,j,l,m}) \leq \Omega_{i,j,l,m} \leq H(1 - y_{i,j,l,m}) + 1, \quad (5)$$

$$x_{i,j,l,m} + y_{i,j,l,m} \geq 1, \quad (6)$$

135 where H is an arbitrary large positive constant.

136 Assigned flow $q_{i,j,k}$ is the number of vehicles in the movement (i, j) turning via lane k .
137 The signal timing decision variables include cycle length, green duration and starts of green.
138 The cycle length ξ is formulated as the reciprocal of the actual cycle length for the purpose
139 of linearization. Thus, the actual cycle length is obtained with $1/\xi$. The green duration of a
140 movement $\phi_{i,j}$, the start of green of a movement $\theta_{i,j}$, the green duration of a lane $\Phi_{i,K}$ and the
141 start of green of a lane $\Theta_{i,k}$ are the fraction of the actual cycle length. Hence, the actual green
142 duration of a movement, the actual start of green of a movement, the actual green duration of
143 a lane and the actual start of green of a lane are $\phi_{i,j}/\xi, \theta_{i,j}/\xi, \Phi_{i,k}/\xi$ and $\Theta_{i,k}/\xi$, respectively.

144 For simplification, in this paper, the reciprocal of the actual cycle length and the fractions
 145 are directly called cycle length, green duration and starts of green. Reserved capacity in this
 146 paper is a common flow multiplier which indicates whether the intersection is overloaded or has
 147 reserved capacity. The original definition can be found in Allsop (1972).

148 Decision variables and their domains are summarized in Table 1. In Table 1, c_{min}, c_{max}
 149 and g_{min} are the minimum cycle length, maximum cycle length and minimum green duration,
 150 respectively.

151 [Table 1 about here.]

152 Optionally, if the signal timing of the pedestrian movements needs to be optimized, the
 153 relevant decision variables are initialized. They are green duration for pedestrian ($\phi_{i,0}, \forall i =$
 154 $1, \dots, N_A$), start of green for pedestrian ($\theta_{i,0}, \forall i = 1, \dots, N_A$), signal sequence for pedestrian
 155 ($\Omega_{i,j,i,0}, \Omega_{i,0,i,j}, \Omega_{l,m,i,0}, \Omega_{i,0,l,m}, \forall i, l = 1, \dots, N_A, j, m \in M, (i, j) \neq (l, m), i = \Gamma(l, m)$) and the
 156 relative auxiliary variables which are similar as the decision variables and the auxiliary variables
 157 between movements.

158 2.3 Objective function

159 The objective is to maximize the reserved capacity, μ , because the goal of this paper is to
 160 gain capacity by allocating the exit lanes. If $\mu > 1$, the intersection has reserved capacity
 161 with $100(\mu - 1)$ percent; if $\mu < 1$, the intersection is overloaded with $100(1 - \mu)$ percent. The
 162 objective function is

$$\max \mu, \tag{7}$$

163 subject to the relations (2) - (6) and the constraints (8) - (23). If pedestrian movements are
 164 considered, constraints (24) - (31) need to be added as well.

165 **2.4 Constraints**

166 The constraints for the vehicle movements adjusted to our model are explained in detail in
 167 constraints (8) - (16) whereas the constraints from the original lane-based method developed
 168 by Wong and Wong (2003) are briefly summarized in constraints (17) - (23), followed by the
 169 constraints about pedestrian movements.

170 **2.4.1 Minimum and maximum number of permitted lanes**

171 Each movement (i, j) must occupy at least one approaching lane. Meanwhile, the number of
 172 permitted lanes for movement (i, j) , is less than or equal to the number of exit lanes of the
 173 movement; otherwise, vehicles on that movement merging into fewer lanes may cause safety
 174 problems (Wong and Wong 2003).

175 For all $i' = \Gamma(i, j)$, the following constraint holds:

$$1 \leq \sum_{k=1}^{L_i} \delta_{i,j,k} \leq \sum_{k'=1}^{E_{i'}} \epsilon_{i,i',k'}, \forall i, i' = 1, \dots, N_A, i' \neq i, j \in M. \quad (8)$$

176 **2.4.2 Conflict elimination on adjacent lanes**

177 The movements on adjacent approaching lanes and exit lanes may conflict with each other.
 178 The conflicts must be eliminated for safety reasons (Wong and Wong 2003). Figure 3 is an
 179 example of conflicts generated on the adjacent approaching lanes. If a right turn is permitted
 180 on lane $k = 2$, the through movement conflict with the left turn on lane $k = 1$ and should not

181 be allowed; if a through movement is permitted on lane $k = 2$, the left turn on the lane $k = 1$
 182 will conflict with the through movement.

183 [Figure 3 about here.]

184 To eliminate conflicts in the approaching lanes, the constraint (9) holds. We denote M' as
 185 the subset of M . If $j = RT$, $M' = \{TH, LT\}$; if $j = TH$, $M' = \{RT\}$; if $j = LT$, $M' = \emptyset$.

$$\delta_{i,j,k+1} - 1 \leq \delta_{i,m,k} \leq 1 - \delta_{i,j,k+1}, \forall i = 1, \dots, N_A, j \in M, k = 1, \dots, L_i - 1, m \in M'. \quad (9)$$

186 The conflicts on the adjacent exit lanes must be also eliminated. The conflicts on the
 187 adjacent exit lanes only occur when two conflicted movements have the same destination arm
 188 and have the signal sequence in parallel. Figure 4 is an example of the conflicts on adjacent
 189 exit lanes. If the exit lane $k' = 2$ is permitted for a through movement, the exit lane $k' = 1$ is
 190 not allowed for a right turn; if the exit lane $k' = 2$ is permitted for a left turn, the exit lane
 191 $k = 1$ is not allowed for a though movement.

192 [Figure 4 about here.]

193 Before handling the exit lanes, we denote the conflict matrix as Ψ . $\psi_{i,j,l,m} \in \Psi, \forall i, l =$
 194 $1, \dots, N_A, j, m \in M$. If $\psi_{i,j,l,m} = 1$, it means a conflict exists between movement (i, j) and
 195 movement (l, m) ; if $\psi_{i,j,l,m} = 0$, otherwise. If $\psi_{i,j,l,m} = 1$, for the movements have the same des-
 196 tination arm which means $i' = \Gamma(i, j) = \Gamma(l, m), \forall i' = 1, \dots, N_A$, and have the signal sequence in
 197 parallel, the constraint (10) holds. According to the definition of the auxiliary binary variables
 198 $x_{i,j,l,m}$ and $y_{i,j,l,m}$, when $x_{i,j,l,m} + y_{i,j,l,m} = 2$, movement (i, j) and (l, m) are in parallel. Hence,
 199 if $x_{i,j,l,m} + y_{i,j,l,m} = 2$ and the exit lane k' is allocated for movement (i, j) , then the exit lane

200 k'' cannot be allocated for movement (l, m) ; if the two movements are not in parallel, the exit
 201 lane allocation of lane k' does not influence the exit lane allocation of lane k'' .

$$\begin{aligned} \epsilon_{i,\Gamma(i,j),k'} + x_{i,j,l,m} + y_{i,j,l,m} - 3 &\leq \epsilon_{l,\Gamma(l,m),k''} \leq 3 - (\epsilon_{i,\Gamma(i,j),k'} + x_{i,j,l,m} + y_{i,j,l,m}), \\ \forall i, l = 1, \dots, N_A, i \neq l, j \in M, m \in M', \Gamma_{i,j} = \Gamma_{l,m}, k' \text{ and } k'' = 1, \dots, E_i, k'' > k'. \end{aligned} \quad (10)$$

202 2.4.3 Order of signal displays

203 The conflicts among movements mainly influence the order of signal displays. If two conflict
 204 movements have different destination arms, they can be either predecessor or successor of each
 205 other. If two conflict movements have the same destination arm, they could be either in parallel
 206 or not. If two movements do not conflict with each other, they could be the predecessor or
 207 successor or in parallel. Conflict matrix records the conflicts among movements and contributes
 208 to the constraint construction.

209 Although in a cycle a signal could appear multiple times so that the signal is both the
 210 predecessor and successor of another, this paper limits this case for simplification as this is not
 211 that common in signal planning. Therefore, no matter whether the movements conflict with
 212 each other, constraint(11) holds, saying that if one movement is the predecessor of another,
 213 another movement can only be the successor of the one; or if one movement is in parallel with
 214 another, another movement is also in parallel with the one. Similar constraint can also be found
 215 in Wong and Wong (2003).

$$\Omega_{i,j,l,m} + \Omega_{l,m,i,j} = 0, \forall i, l = 1, \dots, N_A, i \neq l, j \text{ and } m \in M. \quad (11)$$

216 If $\psi_{i,j,l,m} = 1$, for the movements have different destination arms, i.e. $\Gamma(i, j) \neq \Gamma(l, m)$, they

217 cannot be in parallel. Hence,

$$x_{i,j,l,m} + y_{i,j,l,m} = 1; \forall i, l = 1, \dots, N_A, i \neq l, j \text{ and } m \in M. \quad (12)$$

218 If $\psi_{i,j,l,m} = 1$ and the movement (i, j) and the movement (l, m) , which have the same
 219 destination arm, are in parallel (i.e. $x_{i,j,l,m} = 1, y_{i,j,l,m} = 1$), the exit lane k' can only be assigned
 220 for either the movement (i, j) or the movement (l, m) to avoid conflicts; if the movement (i, j)
 221 and the movement (l, m) are not in parallel (i.e. $x_{i,j,l,m} + y_{i,j,l,m} = 1$), they cannot conflict with
 222 each other, so it does not matter that the exit lane k' is assigned for which movements.

$$\epsilon_{i,i',k'} + \epsilon_{l,i',k'} \leq 3 - (x_{i,j,l,m} + y_{i,j,l,m}), \quad (13)$$

$$\forall i, l = 1, \dots, N_A, i \neq l, i' = \Gamma(i, j) = \Gamma(l, m), j, m \in M, k' = 1, \dots, E_{i'}.$$

223 2.4.4 Identical signal sequence on shared approaching lanes

224 When two movements share the same lane, the signal sequence between the two movements
 225 and the other movements must be the same to avoid internal conflicts on the lanes. Thus, if
 226 movement (i, j) and (i, j') are permitted on the approaching lane k , the values of their signal
 227 sequence indicator must be the same.

$$\delta_{i,j,k} + \delta_{i,j',k} - 2 \leq \Omega_{i,j,l,m} - \Omega_{i,j',l,m} \leq 2 - (\delta_{i,j,k} + \delta_{i,j',k}), \quad (14)$$

$$\forall i, l = 1, \dots, N_A, i \neq l, j, j', m \in M, j' > j, k = 1, \dots, L_i.$$

228 2.4.5 Clearance time

229 If two movements are predecessor/successor of each other, there is at least a clearance time
 230 in-between the green duration of the movements due to potential safety problems. Thus, if the
 231 movement (i, j) is the predecessor of the movement (l, m) , then $x_{i,j,l,m} = 0, y_{i,j,l,m} = 1$, and

232 the start of green of the movement (l, m) must be later than the sum of the start of green and
 233 the green duration of the movement (i, j) and the clearance time (See constraint(15)); if the
 234 movement (i, j) is the successor of the movement (l, m) , then $x_{i,j,l,m} = 1, y_{i,j,l,m} = 0$, and the
 235 constraint(16) holds.

$$\theta_{i,j} + \phi_{i,j} + \omega_{i,j,l,m}\xi \leq \theta_{l,m} + x_{i,j,l,m}, \forall i, l = 1, \dots, N_A, i \neq l, j, m \in M. \quad (15)$$

$$\theta_{l,m} + \phi_{l,m} + \omega_{l,m,i,j}\xi \leq \theta_{i,j} + y_{i,j,l,m}, \forall i, l = 1, \dots, N_A, i \neq l, j, m \in M. \quad (16)$$

236 2.4.6 Constraints from the original lane-based method

237 This section includes the constraints from Wong and Wong (2003)' model. Considering better
 238 readability of this paper, we summarize the constraints below and apply a brief explanation.
 239 More details can be found in Wong and Wong (2003).

$$\sum_{j \in M} \delta_{i,j,k} \geq 1, \forall i = 1, \dots, N_A; k = 1, \dots, L_i, \quad (17)$$

$$-H(1 - \delta_{i,j,k}) \leq \Phi_{i,k} - \phi_{i,j} \leq H(1 - \delta_{i,j,k}), \forall i = 1, \dots, N_A, j \in M, k = 1, \dots, L_i, \quad (18)$$

$$-H(1 - \delta_{i,j,k}) \leq \Theta_{i,k} - \theta_{i,j} \leq H(1 - \delta_{i,j,k}), \forall i = 1, \dots, N_A, j \in M, k = 1, \dots, L_i, \quad (19)$$

$$\mu Q_{i,j} = \sum_{k=1}^{L_i} q_{i,j,k}, \forall i = 1, \dots, N_A, j \in M, \quad (20)$$

$$q_{i,j,k} \leq H\delta_{i,j,k}, \forall i = 1, \dots, N_A, j \in M, k = 1, \dots, L_i, \quad (21)$$

$$-H(2 - \delta_{i,j,k} - \delta_{i,j,k+1}) \leq v_{i,k} - v_{i,k+1} \leq H(2 - \delta_{i,j,k} - \delta_{i,j,k+1}), \quad (22)$$

$$\forall i = 1, \dots, N_A, j \in M, k = 1, \dots, L_i - 1,$$

$$u_{i,k} = \frac{v_{i,k}}{\Phi_{i,k} + e\xi} \leq u_{max,i,k}, \forall i = 1, \dots, N_A, k = 1, \dots, L_i, \quad (23)$$

240 where H is an arbitrary large positive constant, $v_{i,k} = \sum_{j \in M} \frac{q_{i,j,k}}{s_j}$ is flow factor of lane k in
241 arm i , s_j is the saturation flow of movements on exclusive lanes, e is the difference between
242 actual green time and effective green time and predefined as 1 s, and $u_{max,i,k}$ is the maximum
243 acceptable degree of saturation.

244 Due to the completeness principle of signal timing plan design, all movements should be
245 included in the signal cycle, so constraint (17) holds. Signal timing is the most important
246 issue to be solved. When more than one movements share one lane, the signal settings of these
247 movements are identical to avoid internal conflict on the lane (constraint (18) and (19)). Traffic
248 flows must be treated as well which refers to constraint (20) - (23). The maximum amount of
249 traffic increase, which confirms the reasonable performance of the intersection, is the product
250 of reserved capacity μ and demands $Q_{i,j}$. The maximum amount is equal to the sum of traffic
251 flows of movement (i, j) being assigned to all lanes on arm i (constraint (20)). The assigned
252 flow $q_{i,j,k}$ must be 0 if movement (i, j) is not permitted on lane k (constraint (21)). If two
253 movements share the same lane and two adjacent lanes are permitted, the degree of saturation
254 on both lanes is identical, resulting in equal flow factors, for signal settings of these adjacent
255 lanes are the same (constraint (22)). The degree of saturation should be no more than the
256 maximum acceptable degree of saturation (constraint (23)).

257 **2.4.7 Pedestrian movement**

258 If the decision variables of pedestrian movements need to be determined, the constraints for the
259 signal sequence and the clearance time will be handled. However, firstly all relevant decision
260 variables must be in their domain.

$$0 \leq \theta_{i,0} \leq 1, \forall i = 1, \dots, N_A, \quad (24)$$

$$g_{min,0}\xi \leq \phi_{i,0} \leq 1, \forall i = 1, \dots, N_A, \quad (25)$$

261 where $g_{min,0}$ is the minimum green duration for the pedestrian movement.

262 For each pedestrian movement, it can conflict with both the movements starting from an
 263 arm and the movements ending at the same arm. To distinguish the two cases, we use movement
 264 (i, j) as the movement starting at arm i and movement (l, m) as the movement ending at arm
 265 i . Similar as the vehicle movements, the auxiliary binary variables, $x_{i,j,i,0}, y_{i,j,i,0}, x_{l,m,i,0}$ and
 266 $y_{l,m,i,0}$ are introduced. $\forall i = 1, \dots, N_A, j \in M$, if $\Omega_{i,j,i,0} \geq 1, x_{i,j,i,0} = 0; \Omega_{i,j,i,0} < 1, x_{i,j,i,0} = 1$; if
 267 $\Omega_{i,j,i,0} \geq 0, y_{i,j,i,0} = 1; \Omega_{i,j,i,0} < 0, y_{i,j,i,0} = 0. \forall l = 1, \dots, N_A, m \in M, \Gamma(l, m) = i$, if $\Omega_{l,m,i,0} \geq 1,$
 268 $x_{l,m,i,0} = 0; \Omega_{l,m,i,0} < 1, x_{l,m,i,0} = 1$; if $\Omega_{l,m,i,0} \geq 0, y_{l,m,i,0} = 1; \Omega_{l,m,i,0} < 0, y_{l,m,i,0} = 0$. Then
 269 $x_{i,j,i,0}, y_{i,j,i,0}, x_{l,m,i,0}$ and $y_{l,m,i,0}$ follow the relations similar as the relations (2) - (6).

270 For the movement (i, j) starting from arm i ,

$$\Omega_{i,j,i,0} + \Omega_{i,0,i,j} = 0, \forall i = 1, \dots, N_A, j \in M, \quad (26)$$

$$\theta_{i,j} + \phi_{i,j} + \omega_{i,j,i,0}\xi \leq \theta_{i,0} + x_{i,j,i,0}, \forall i = 1, \dots, N_A, j \in M, \quad (27)$$

$$\theta_{i,0} + \phi_{i,0} + \omega_{i,0,i,j}\xi \leq \theta_{i,j} + y_{i,j,i,0}, \forall i = 1, \dots, N_A, j \in M. \quad (28)$$

271 For the movement (l, m) merging at arm i ,

$$\Omega_{l,m,i,0} + \Omega_{i,0,l,m} = 0, \forall i, l = 1, \dots, N_A, m \in M, i = \Gamma(l, m), \quad (29)$$

$$\theta_{l,m} + \phi_{l,m} + \omega_{l,m,i,0}\xi \leq \theta_{i,0} + x_{l,m,i,0}, \forall i, l = 1, \dots, N_A, m \in M, i = \Gamma(l, m), \quad (30)$$

$$\theta_{i,0} + \phi_{i,0} + \omega_{i,0,l,m}\xi \leq \theta_{l,m} + y_{l,m,i,0}, \forall i, l = 1, \dots, N_A, m \in M, i = \Gamma(l, m), \quad (31)$$

272 where $\omega_{i,j,i,0}, \omega_{i,0,i,j}, \omega_{l,m,i,0}$ and $\omega_{i,0,l,m}$ are the clearance time.

273 **3 Numerical examples**

274 **3.1 Numerical configurations**

275 The layout of the studied intersection can be found in Figure 2 whereas the number of ap-
276 proaching lanes and exit lanes vary according to Table 2, which summarizes the details of the
277 number of approaching lanes and exit lanes for each intersection.

278 [Table 2 about here.]

279 Additional to the configurations of studied intersections, traffic demand, conflict matrix,
280 saturation flows for each movement and the values of bounds for signal timing have to be given.
281 The traffic demand and the conflict matrix can be seen in Table 3 and Table 4, respectively.
282 In this section, only protected left turns are used. If permitted left turns could be present, the
283 conflicts between left turns and the opposing through movements and the conflicts between left
284 turns and the opposing right turns should be marked as 0 in Table 4. However, the optimal result
285 could show that the "permitted left turns" are not in the same green duration as their opposing
286 through movement because compatible movements could be in different green duration, and
287 then the left turns are actually protected. According to HCM (2000), the saturation flow of
288 through movement is assigned as 1900 veh/h; the saturation flow of right turn is 1615 veh/h;
289 the saturation flow of left turn is 1805 veh/h. The cycle length is within the range of 60 s
290 and 90 s. The green duration for all movements must be no less than 5 s. The clearance
291 time between conflicted movements is 5 s. The maximum acceptable degrees of saturation,
292 $u_{max,i,k}, \forall i = 1, \dots, N_A, k = 1, \dots, L_i$, are assigned as 90%.

293 For the purpose of observation of the capacity increase, we compare our proposed model
294 with a reference model which is developed by Wong and Wong (2003) whereas the signal timing
295 for pedestrian movements is excluded as the focus of this paper is the exit lane allocation for
296 vehicles. The reference model and the proposed model are implemented in Java integrated with
297 ILOG CPLEX 12.8 which is a professional solver for the linear programming. The PC, whose
298 CPU is Intel Core i7 with 2.7GHz and memory is 16.0GB, performs the numerical example.

299 [Table 3 about here.]

300 [Table 4 about here.]

301 **3.2 Overall optimization results**

302 Table 5 summarizes the overall optimization results for both the reference model and the
303 proposed model so that their optimization results can be better compared. The number of
304 variables and the number of constraints represent the problem sizes of the models. The problem
305 size of the proposed model is larger than the problem size in the reference model, but the
306 proposed model can still be efficiently solved. The optimal reserved capacities increase as the
307 number of approaching lanes and the number of exit lanes increase. With the positive values of
308 μ (see Column 8), all intersection have reserved capacity which can be found in Column 9. The
309 capacity increase due to the exit lane allocation is in the final column of Table 5. Compared to
310 the reference model, the capacity increase for intersection 1, 2, 3 and 4 is 1.14%, 9.23%, 9.23%
311 and 13.23%, respectively. This means that appropriately allocating the exit lanes can increase
312 the capacity, and as the number of approaching lanes and the number of exit lanes increase,
313 the capacity increase goes larger.

3.3 Signal timing plan and lane allocation

Signal timing plan reflects the results of cycle length, green duration, and starts of green and signal sequences. Lane allocation reflects the results of approaching lane permission and exit lane allocation. Signal timing plan and lane allocation are demonstrated in this section together because the link between the signal timing and lane allocation can be clearly built. By maximizing the reserved capacity, the cycle length for all intersections is 90 s. The signal timing plan for each intersection is shown in Figure 5(a), Figure 6(a), Figure 7(a) and Figure 8(a), respectively. In each signal timing plan, there are green duration which incompatible movements have the right-of-way at the same time. The green duration is numbered at the bottom of the signal timing plans. Therefore, Figure 5(b), Figure 6(b), Figure 7(b) and Figure 8(b) display how the incompatibility is eliminated for each intersections, respectively. The lane markings, i.e. the approaching lane allocation, can also be found in these figures.

At intersection 1, the conflicts of incompatible movements in four intervals of green duration are eliminated (see Figure 5(b)). In the first green duration, the through movement from arm 2 goes to the exit lane 2 and 3 in arm 4, and the left turn from arm 3 goes to the exit lane 1, so they do not conflict. The clearance time between these two movements is then not necessary, which means longer green duration for the relevant movements. Similarly, the elimination of conflicts can be found in the second, the third and the fourth interval of green duration. At intersection 2, there three intervals of green duration for the incompatible movements (see Figure 6(b)). As the exit lane of the right turn from arm 4 can be well allocated, this movement has a full-green

335 duration. Four intervals of green duration for the incompatible movements at intersection 3
336 can be seen in Figure 7(b). Due to the appropriate exit lane allocation, at this intersection, the
337 right turns from arm 2, arm 3 and arm 4 have a full-green duration as well, and their clearance
338 time is not necessary. Intersection 4 also has four green duration for incompatible movements
339 (see Figure 7(b)). The movements from three arms could be in the green duration at the same
340 time at both intersection 3 and 4 as their number of exit lanes increases.

341 [Figure 5 about here.]

342 [Figure 6 about here.]

343 [Figure 7 about here.]

344 [Figure 8 about here.]

345 **3.4 Traffic performance measurement**

346 In this section, the assigned flows for each lane are recorded and then lane flows, lane saturation
347 flow, lane capacities and degree of saturation are calculated and shown in Table 6. With these
348 traffic performance measures, we can better evaluate the results. Column 3-6 display the results
349 of assigned flows. Then for each lane, the lane flows can be obtained by summing up the assigned
350 flow by lane (see Column 7). Meanwhile, the lane saturation flow can be calculated according
351 to the assigned flow and the lane flow. The value of lane saturation flow is shown in Column
352 8. The capacity of the lane in Column 9 is the product of saturation flow and lane green split.
353 Degree of saturation for each lane in Column 10 is then the lane flow divided the capacity.
354 When movements share the same lane, their degrees of saturation are the same and it can be

355 found from the values of the degrees of saturation. As the number of approaching lane increases
356 from intersection 1 to intersection 4, the degree of saturation obviously decreases. However, as
357 intersection 2 and 3 has the same optimal reserved capacity, they have the same value of the
358 performance measures. Right turns benefit from the exit lane allocation, and their degrees of
359 saturation are relatively small. For example, at intersection 2 & 3, the degree of saturation of
360 the right turn in arm 2 (from arm 2 to arm 1) is 0.2594, and the degree of saturation of the
361 right turn in arm 4 (from arm 4 to arm 3) is 0.1238. They are much fewer than the degrees
362 of saturation for the rest movements. That means, they gain capacities due to the exit lane
363 allocation.

364 [Table 6 about here.]

365 4 Discussion

366 We propose a method to solve the exit lane allocation problem as a mixed integer linear pro-
367 gramming in the lane-based method. Capacity increases after the exit lanes are suitably al-
368 located for incompatible movements ending at the same destination arm. We apply exit lane
369 indicators and precise sequence states among incompatible movements: predecessor, in parallel
370 and successor, so that extend the lane-based signal optimization method. We would like to
371 discuss the reasons why exit lane allocation can increase capacities and the safety aspect of exit
372 lane allocation.

373 The incompatible movements require clearance time in-between for safety reasons, but the
374 well allocated exit lanes avoid the conflicts and so as the clearance time. The green duration
375 for related incompatible movements thereby becomes longer, resulting in larger capacities for

376 these movements. It is interesting to notice that the number of exit lanes influence the capacity
377 increase. As the number of exit lanes increases, the capacity increases. Large number of exit
378 lanes allows more approaching lanes for a movement, including the incompatible movements
379 ending at the same destination arm. A large number of approaching lanes for a movement
380 obviously increases the capacity. On another hand, the total green duration of incompatible
381 movements increases from intersection 1 to 4. Longer green duration is another factor to
382 increase the capacity.

383 To ensure the safety of the exit lanes of incompatible movements ending at the same desti-
384 nation, the design of pavement markings at intersections needs to be considered. The pavement
385 markings can guide drivers to correctly follow a lane to avoid conflicts from the vehicles of their
386 incompatible movements, when they have the right-of-way in the overlapping green duration.
387 Figure 9 shows the pavement markings of all exit lanes of incompatible movements being cor-
388 rectly guided. The standard of pavement marking design in Germany can be found in FGSV
389 (1993). Although our model can get the exit lane allocation for all movements, the exit lane
390 allocation for compatible movements and the incompatible movements in different green dura-
391 tion is less important because they do not mutually conflict. Thus, only the pavement markings
392 of the incompatible movements having overlapping green duration are kept in Figure 9 to re-
393 duce the complexity of pavement markings at intersections. However, as the number of lanes
394 increases, the pavement markings still become complicated. We should pay attention on the
395 potential accidents caused by overloaded information of pavement markings.

396 [Figure 9 about here.]

397 **5 Conclusions**

398 A lane-based signal optimization method integrating the exit lane allocation at isolated inter-
399 sections is proposed. The proposed method can determine lane markings, exit lane allocation,
400 signal sequences, green duration, starts of green, cycle length, assigned flows and reserved
401 capacity. The precise sequence states referring to predecessor, successor and in parallel are
402 applied. With the linear objective of maximizing the reserved capacity and constraints, this
403 problem is formulated as a mixed integer linear programming, which can be efficiently solved
404 with standard branch-and-cut algorithm. By applying this method, the exit lane allocation
405 for incompatible movements ending at the same destination arm becomes feasible, which is a
406 significant extension on the original lane-based method. Numerical example shows that the
407 appropriate exit lane allocation can gain capacity at intersections. Pavement marking design
408 suggestions also contribute to avoiding accidents in the practical application.

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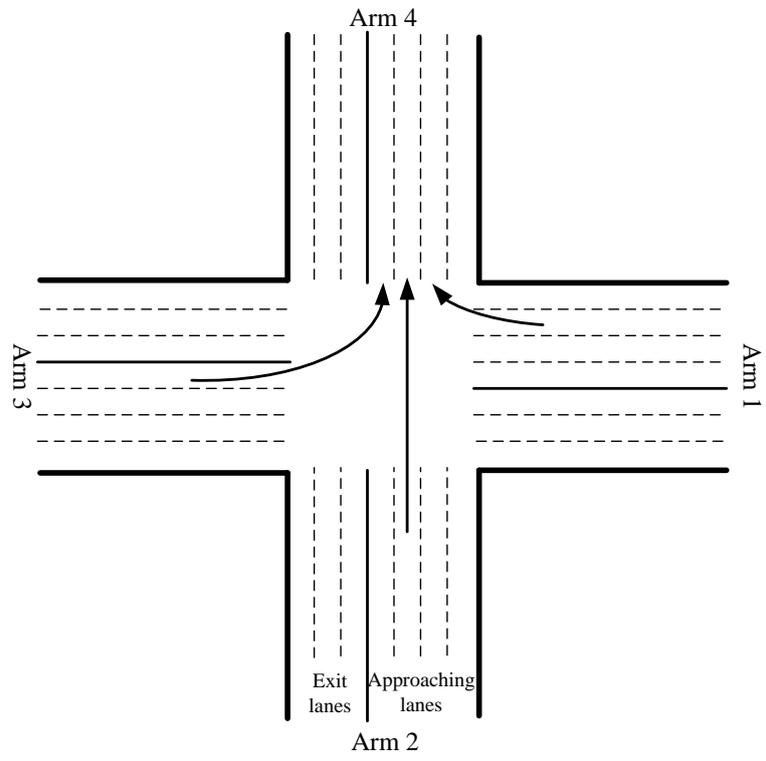


Figure 1: Appropriate exit lane allocation can eliminate the conflicts in destination arms

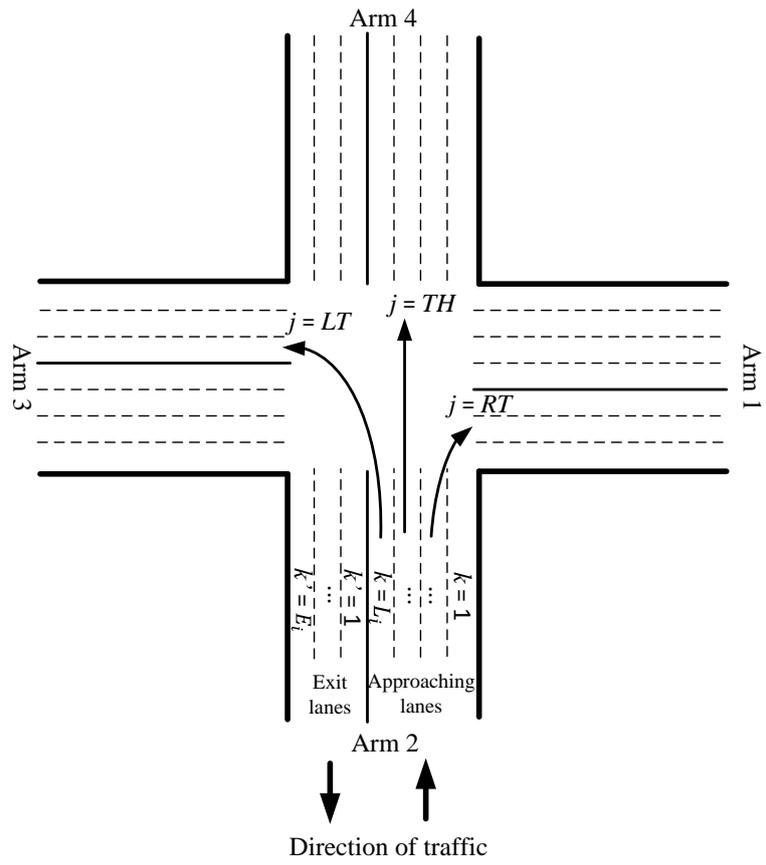


Figure 2: Intersection representation

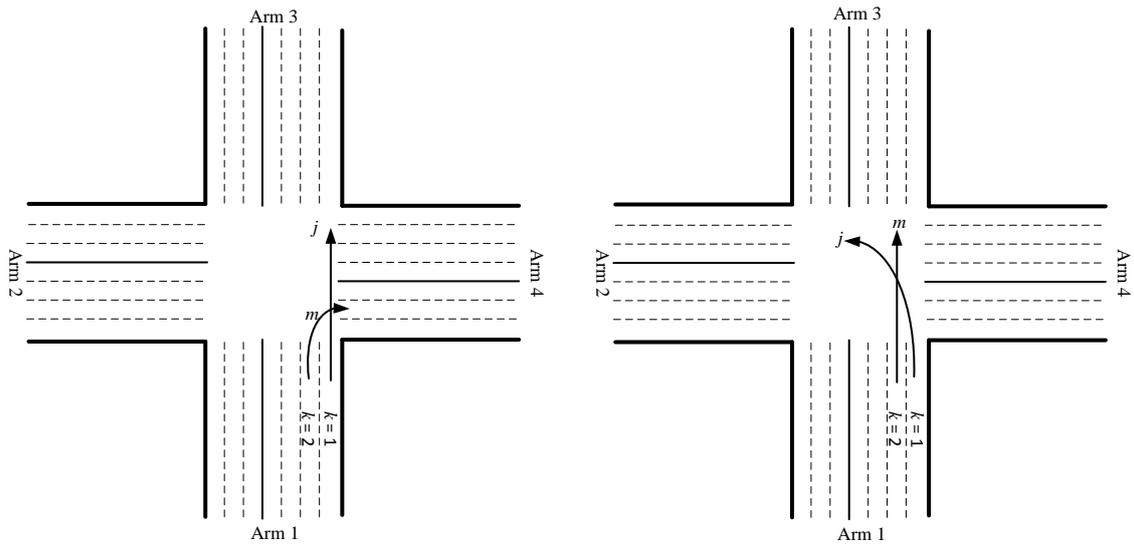


Figure 3: Examples of conflict on adjacent approaching lanes

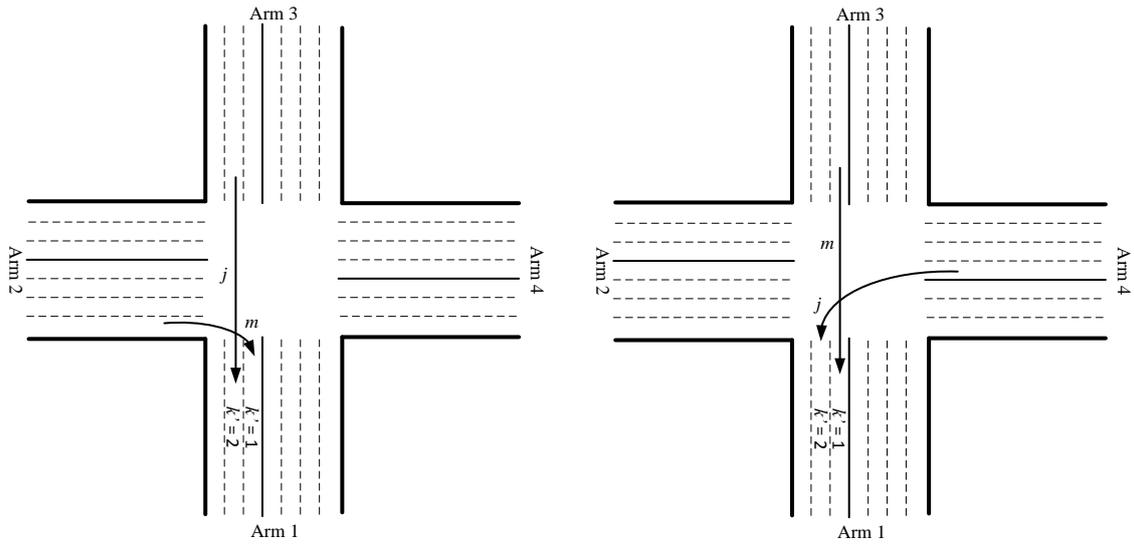


Figure 4: Examples of conflict on adjacent exit lanes

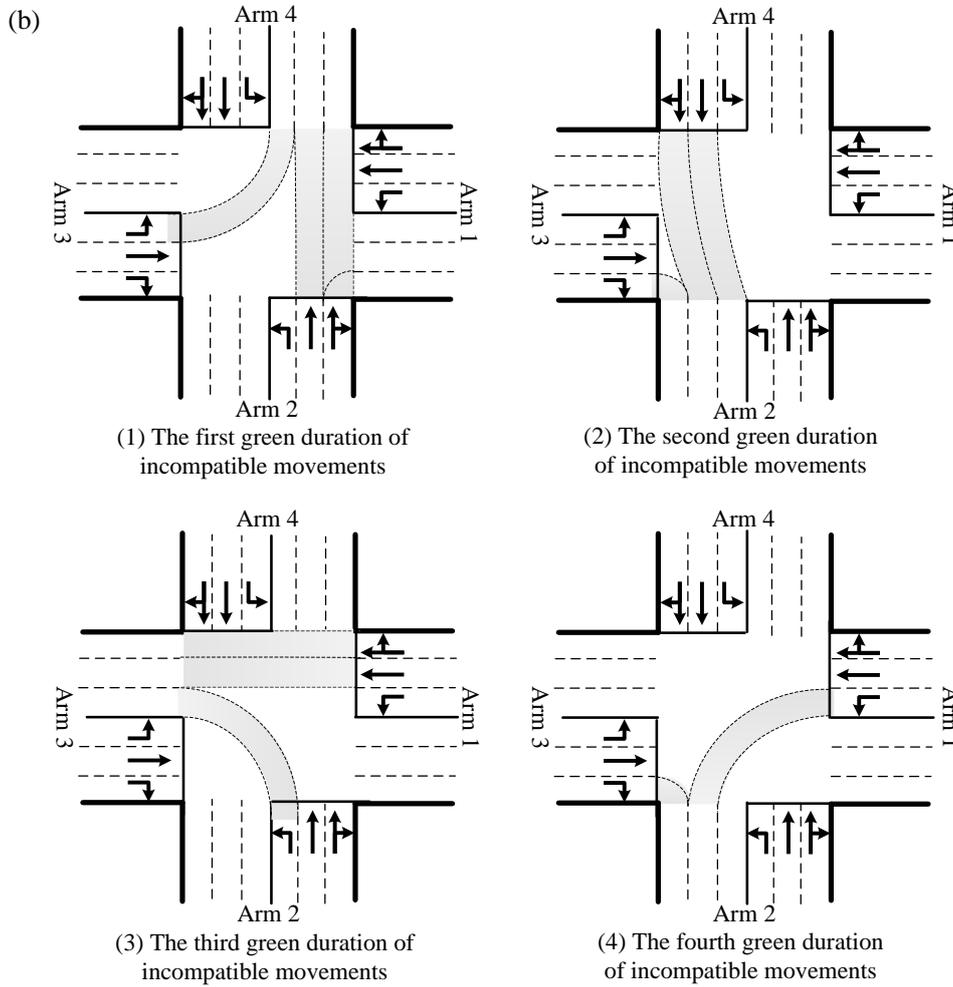
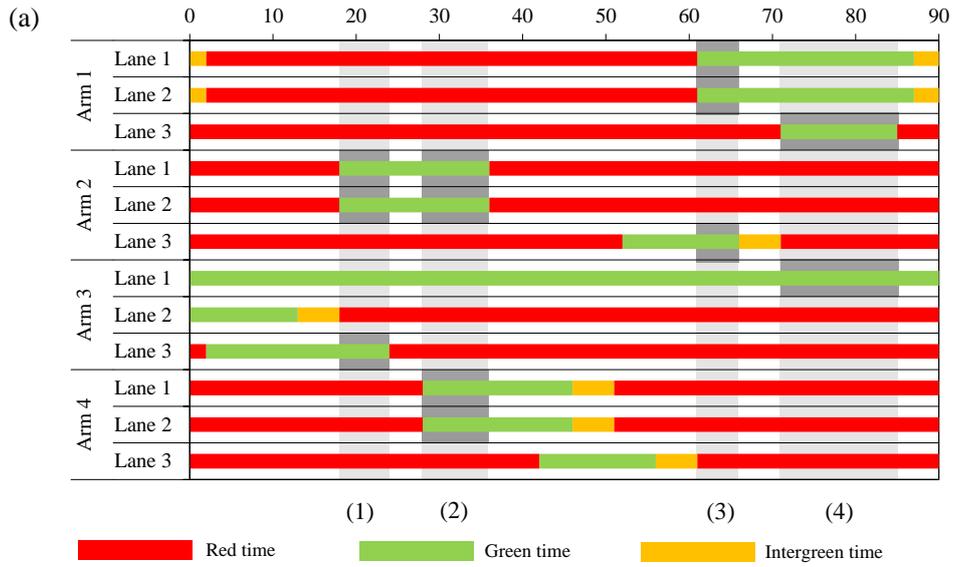


Figure 5: (a) Signal timing plan of intersection 1; (b) Exit lane allocation to eliminate conflicts at intersection 1

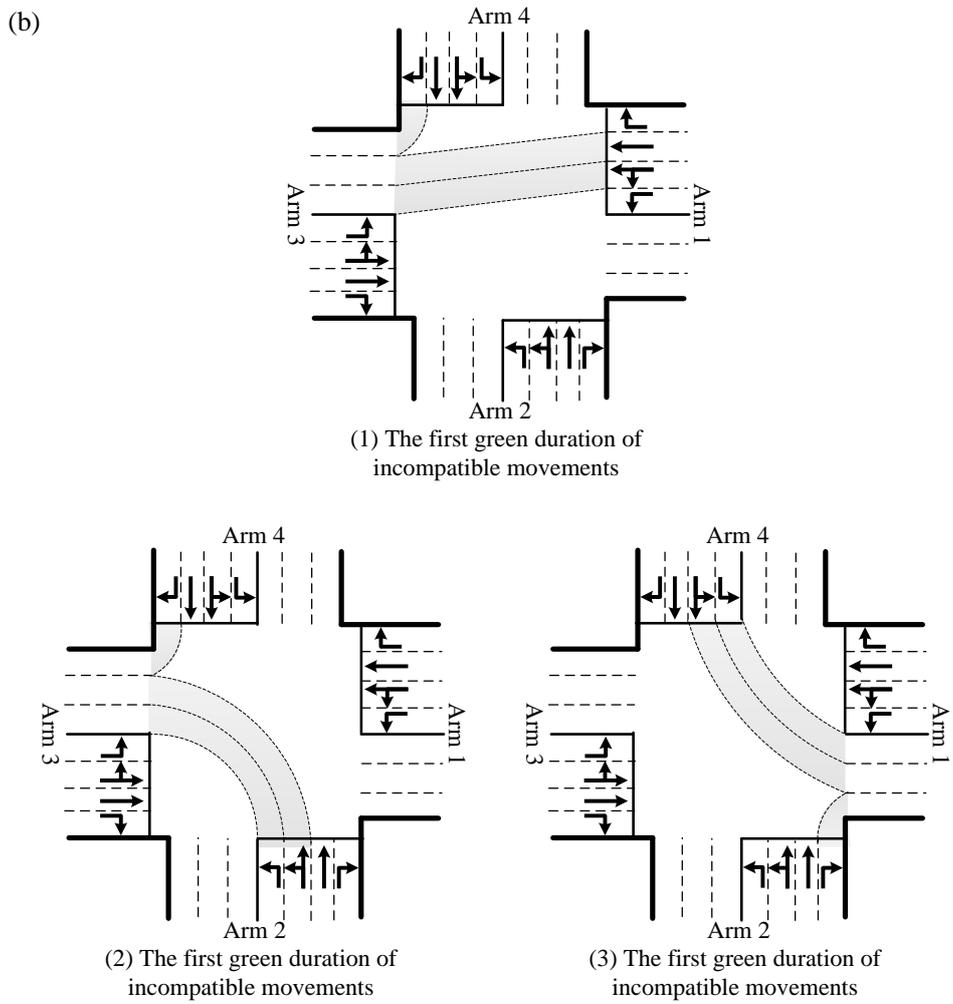
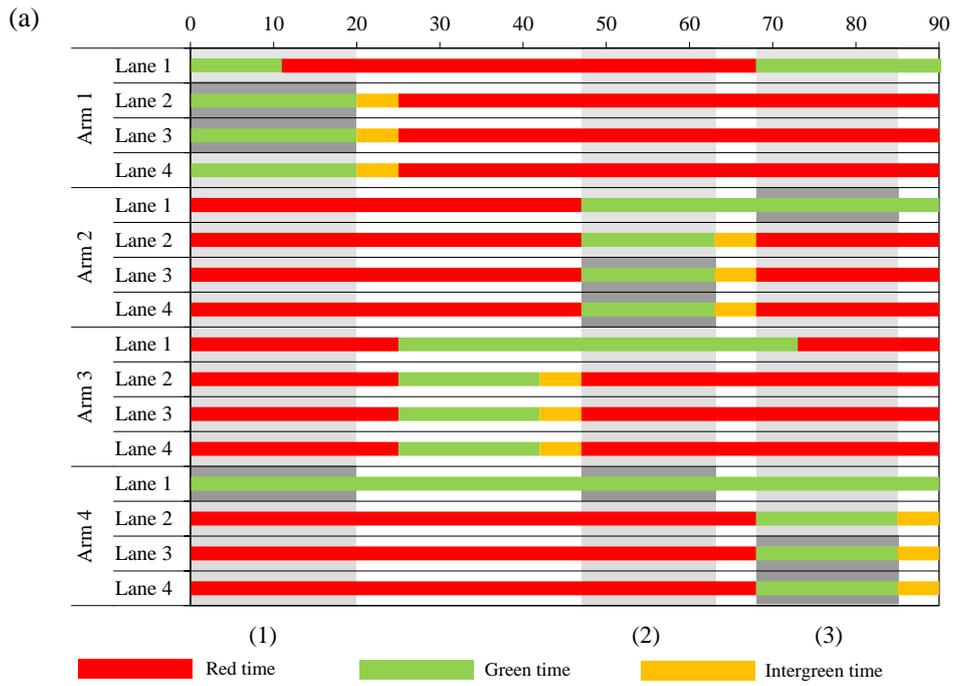


Figure 6: (a) Signal timing plan of intersection 2; (b) Exit lane allocation to eliminate conflicts at intersection 2

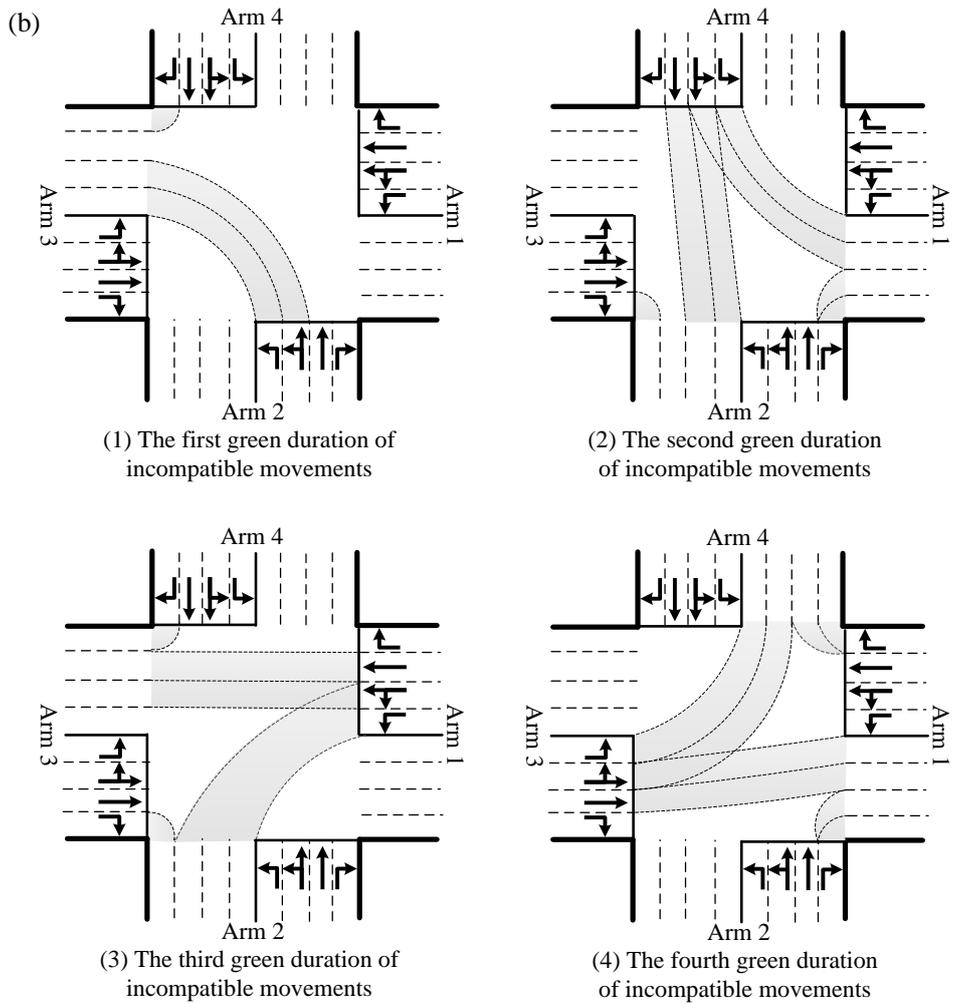
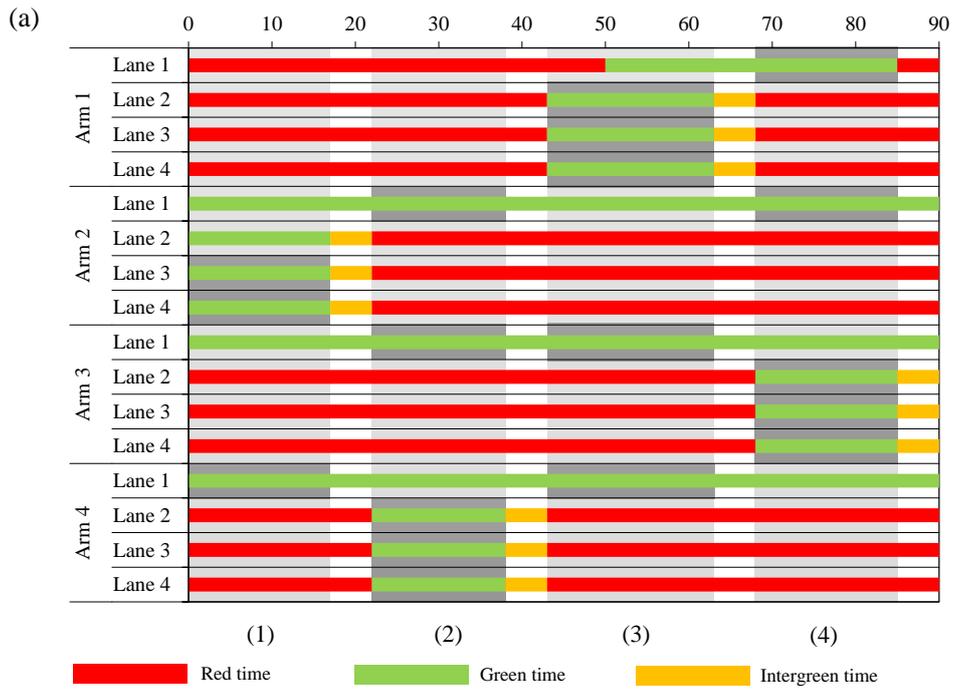


Figure 7: (a) Signal timing plan of intersection 3; (b) Exit lane allocation to eliminate conflicts at intersection 3

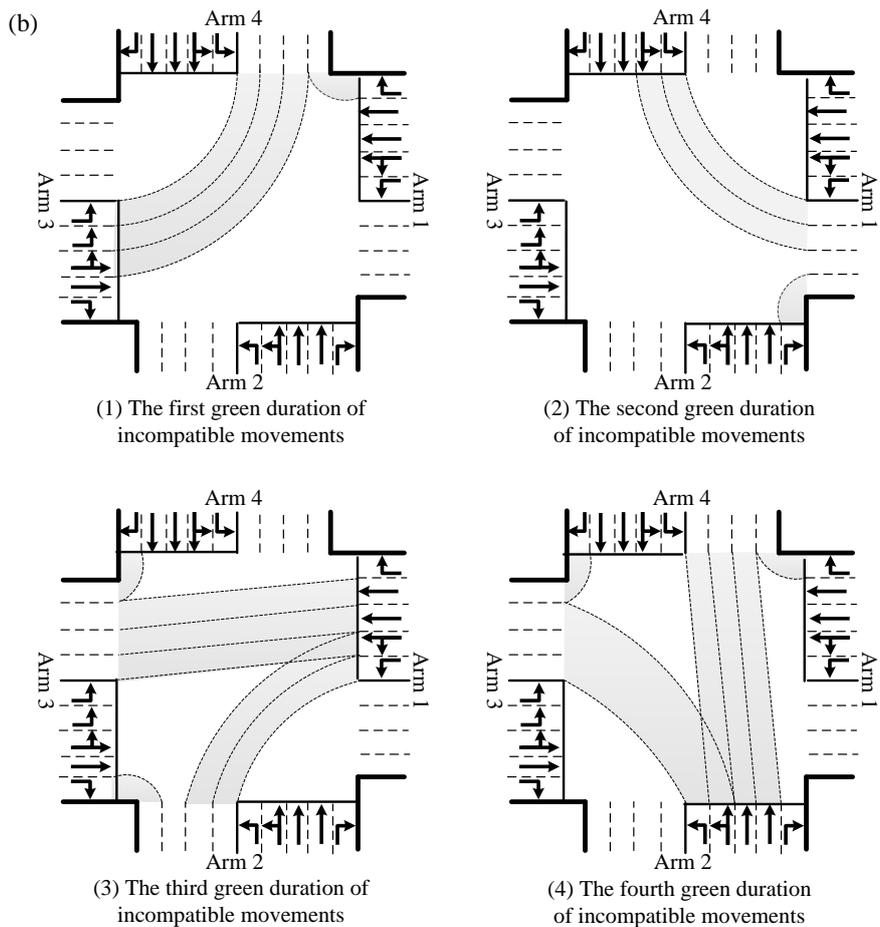
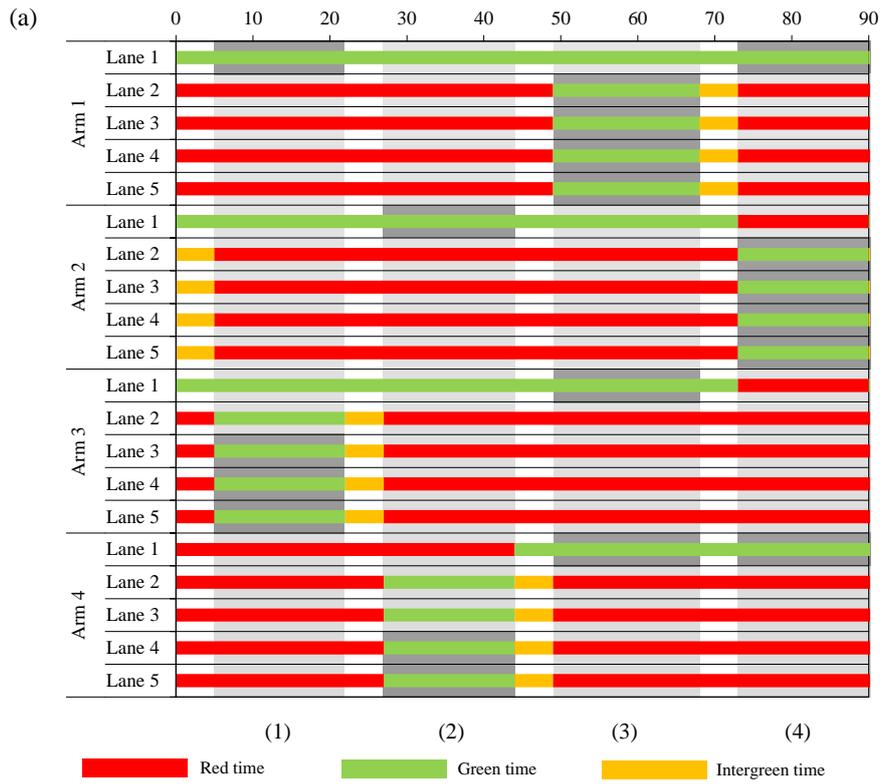


Figure 8: (a) Signal timing plan of intersection 4; (b) Exit lane allocation to eliminate conflicts at intersection 4

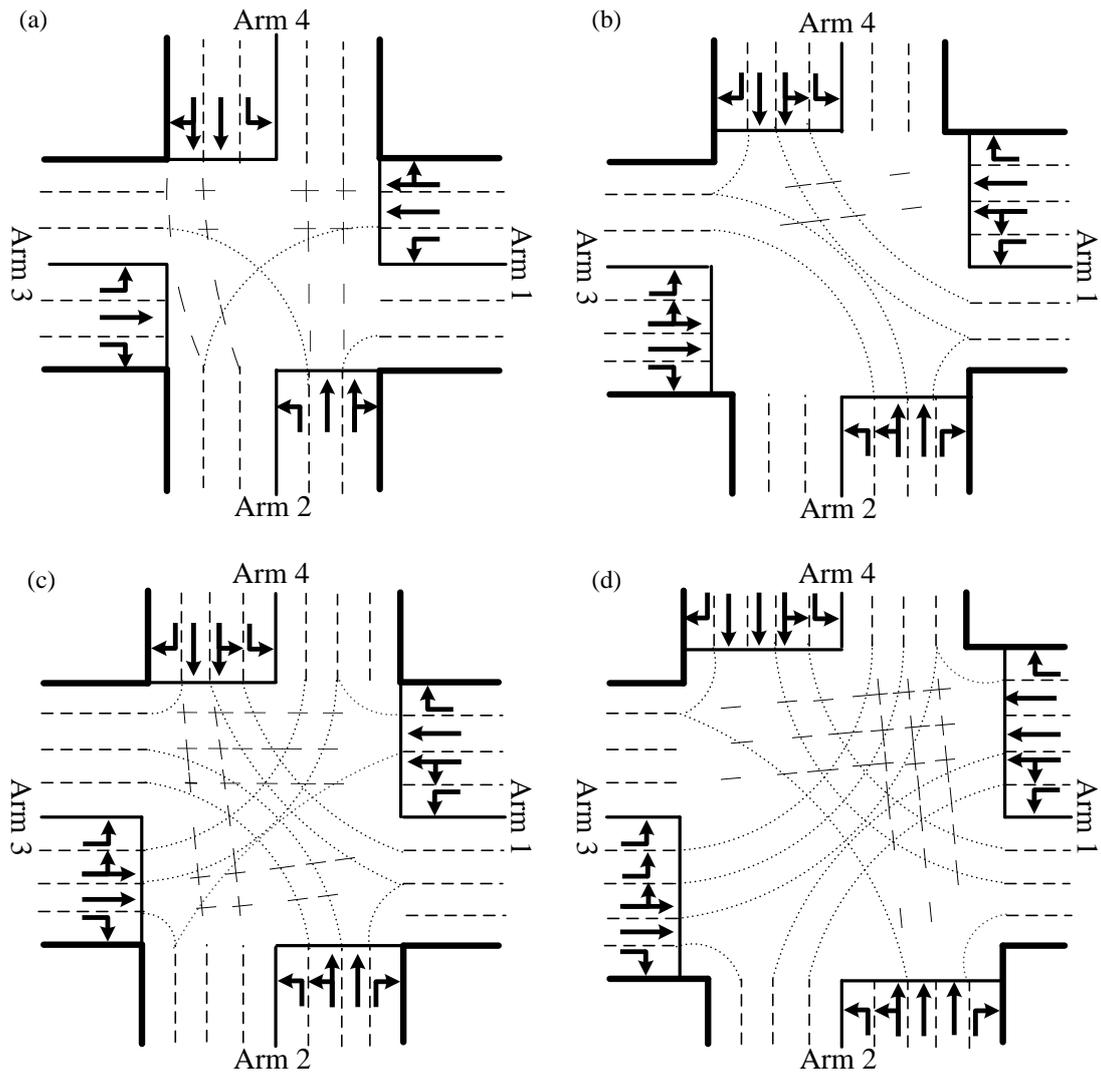


Figure 9: Pavement markings of exit lanes of incompatible movements: (a) Pavement markings of intersection 1; (b) Pavement markings of intersection 2; (c) Pavement markings of intersection 3; (d) Pavement markings of intersection 4

Table 1: List of decision variables

Decision variables	Notation	Domain
Approaching lane permission indicator	$\delta_{i,j,k}$	$\delta_{i,j,k} \in \{0, 1\}$
Exit lane permission indicator	$\epsilon_{i,i',k}$	$\epsilon_{i,j,k} \in \{0, 1\}$
Sequence indicator	$\Omega_{i,j,l,m}$	$\Omega_{i,j,l,m} \in \{-1, 0, 1\}$
Assigned flow of a movement on a lane (veh/h)	$q_{i,j,k}$	$q_{i,j,k} \in [0, \infty)$
Reciprocal of cycle length	ξ	$\xi \in [1/c_{max}, 1/c_{min}]$
Green duration for lane	$\Phi_{i,k}$	$\Phi_{i,K} \in [g_{min}\xi, 1]$
Start of green for lane	$\Theta_{i,k}$	$\Theta_{i,k} \in [0, 1]$
Green duration for movement	$\phi_{i,j}$	$\phi_{i,j} \in [g_{min}\xi, 1]$
Start of green for movement	$\theta_{i,j}$	$\theta_{i,j} \in [0, 1]$
Reserved capacity	μ	$\mu \in [0, \infty)$

Table 2: Studied intersection configuration

Intersection	Number of approaching lanes	Number of exit lanes
1	3	3
2	4	3
3	4	4
4	5	4

Table 3: Traffic demand for studied intersections (veh/h)

From arm	To arm			
	1	2	3	4
1	-	300	400	200
2	200	-	300	200
3	400	200	-	300
4	200	300	200	-

Table 4: Conflict matrix for studied intersections

Arm, i	Direction, j	1			2			3			4		
		RT	TH	LT									
1	RT	0	0	0	0	1	0	0	0	1	0	0	0
	TH	0	0	0	0	1	1	0	0	1	1	1	1
	LT	0	0	0	0	1	1	1	1	0	0	1	1
2	RT	0	0	0	0	0	0	0	1	0	0	0	1
	TH	1	1	1	0	0	0	0	1	1	0	0	1
	LT	0	1	1	0	0	0	0	1	1	1	1	0
3	RT	0	0	1	0	0	0	0	0	0	0	1	0
	TH	0	0	1	1	1	1	0	0	0	0	1	1
	LT	1	1	0	0	1	1	0	0	0	0	1	1
4	RT	0	1	0	0	0	1	0	0	0	0	0	0
	TH	0	1	1	0	0	1	1	1	1	0	0	0
	LT	0	1	1	1	1	0	0	1	1	0	0	0

Table 5: Summary of overall optimization results

Intersection	Model	Number of variables			Number of Constraints	Computing time (s)	μ	Reserved capacity	Capacity increase
		Binary	Integer	Continuous					
1	Reference	92	0	86	898	2.48	1.3502	35.02%	-
	Proposed	336	144	86	2708	3.49	1.3656	36.56%	1.14%
2	Reference	104	0	106	1406	10.73	1.7983	79.83%	-
	Proposed	348	144	106	3112	9.61	1.9643	96.43%	9.23%
3	Reference	104	0	106	1406	7.22	1.7983	79.83%	-
	Proposed	360	144	106	3208	9.31	1.9643	96.43%	9.23%
4	Reference	116	0	126	2026	29.16	2.3129	131.29%	-
	Proposed	372	144	126	3612	16.63	2.6190	161.90%	13.23%

Table 6: Traffic performance measurement results

From arm	Lane	To arm (Assigned flow, veh/h)				Lane flow (veh/h)	Saturation flow (veh/h)	Capacity (veh/h)	Degree of saturation
		1	2	3	4				
Intersection 1 ($\mu = 1.3656$, cycle length = 90 s):									
1	1			24	300	324	1633	471	0.6879
	2				376	376	1900	548	0.6861
	3		200			200	1805	280	0.7143
2	1	200			32	232	1649	329	0.7052
	2				268	268	1900	380	0.7053
	3			200		200	1805	280	0.7143
3	1		400			400	1615	1615	0.2477
	2	200				200	1900	274	0.7299
	3				300	300	1805	441	0.6803
4	1		32	200		232	1649	329	0.7052
	2		268			268	1900	380	0.7053
	3	200				200	1805	280	0.7143
Intersection 2&3 ($\mu = 1.9643$, cycle length = 90 s):									
1	1				300	300	1615	628	0.4777
	2				204	204	1900	422	0.4834
	3		7	196		203	1897	421	0.4822
	4		193			193	1805	401	0.4813
2	1	200				200	1615	771	0.2594
	2				170	170	1900	358	0.4749
	3			38	130	168	1878	354	0.4746
	4				162	162	1805	340	0.4765
3	1		400			400	1615	861	0.4646
	2	172				172	1900	358	0.4804
	3	28			137	165	1820	343	0.4810
	4				163	163	1805	340	0.4794
4	1				200	200	1615	1615	0.1238
	2		170			170	1900	358	0.4749
	3	38	130			168	1878	354	0.4746
	4	162				162	1805	340	0.4765
Intersection 4 ($\mu = 2.6190$, cycle length = 90 s):									
1	1				300	300	1615	1525	0.1967
	2				153	153	1900	422	0.3626
	3				153	153	1900	422	0.3626
	4		55	95		150	1864	414	0.3623
	5		145			145	1805	401	0.3616
2	1	200				200	1615	1309	0.1528
	2				128	128	1900	358	0.3575
	3				128	128	1900	358	0.3575
	4			79	45	124	1838	347	0.3573

Continued on next page

Table 6 Continued: Traffic performance measurement results

From arm	Lane	To arm (Assigned flow, veh/h)				Lane flow (veh/h)	Saturation flow (veh/h)	Capacity (veh/h)	Degree of saturation
		1	2	3	4				
3	5			121		121	1805	340	0.3559
	1		400			400	1615	1309	0.3056
	2	129				129	1900	358	0.3603
	3	71			55	126	1857	350	0.3600
	4				122	122	1805	340	0.3588
4	5				122	122	1805	340	0.3588
	1			200		200	1615	843	0.2372
	2		128			128	1900	358	0.3575
	3		128			128	1900	358	0.3575
	4	79	45			124	1838	347	0.3573
	5	121				121	1805	340	0.3559